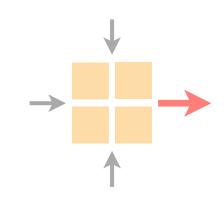
Communication Networks

Spring 2022





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nsg.ee.ethz.ch

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Materials inspired from Scott Shenker & Jennifer Rexford

Last week on Communication Networks

TCP: Reliable, in-order delivery

TCP provides a connection-oriented, reliable, bytestream transport service

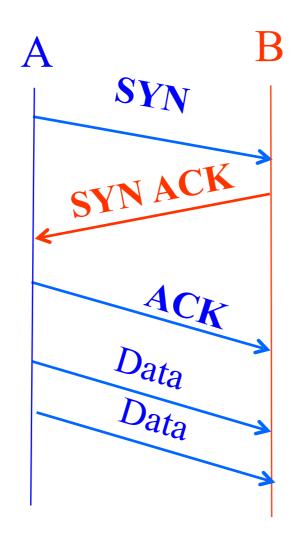
What UDP provides, plus:

- Retransmission of lost and corrupted packets
- Flow control (to not overflow receiver)
- Congestion control (to not overload network)
- "Connection" set-up & tear-down

TCP Header

Source port			Destination port		
Sequence number					
Acknowledgment					
HdrLen	0	Flags	Advertised window		
Checksum			Urgent pointer		
Options (variable)					
Data					

Establishing a TCP Connection



Each host tells its ISN to the other host.

Three-way handshake to establish connection

- Host A sends a SYN (open; "synchronize sequence numbers")
- Host B returns a SYN acknowledgment (SYN ACK)
- Host A sends an ACK to acknowledge the SYN ACK

This week on Communication Networks

Congestion Control

DNS

Introduction to 2nd project



ethz.ch ⇒
129.132.19.216

reliable transport starts *today*!

Congestion Control

DNS

Introduction to 2nd project



Congestion Control

DNS

Introduction to 2nd project

ethz.ch ⇒
129.132.19.216

Internet has one global system for

addressing hostsby design

naming hostsby "accident", an afterthought

Internet has one global system for

naming hosts

DNS

by "accident", an afterthought

Using Internet services can be divided into four logical steps

step 1	A person has name of entity she wants to access	www.ethz.ch
step 2	She invokes an application to perform the task	Chrome
step 3	The application invokes DNS to resolve the name into an IP address	129.132.19.216
step 4	The application invokes transport protocol to establish an app-to-app connection	

The DNS system is a distributed database which enables to resolve a name into an IP address



In practice, names can be mapped to more than one IP



In practice, IPs can be mapped by more than one name

name DNS	IP address		
www.ethz.ch	129.132.19.216		
www.vanbever.eu	82.130.102.71		
www.route-aggregation.net	82.130.102.71		

82.130.102.71

comm-net.ethz.ch

How does one resolve a name into an IP?

initially

all host to address mappings were in a file called hosts.txt

in /etc/hosts

problem

scalability in terms of query load & speed management

consistency

availability

When you need... more flexibility, you add... a layer of indirection

When you need... more scalability,

you add... a hierarchical structure

To scale, DNS adopt three intertwined hierarchies

naming structure hierarchy of addresses

https://www.ee.ethz.ch/de/departement/

management hierarchy of authority

over names

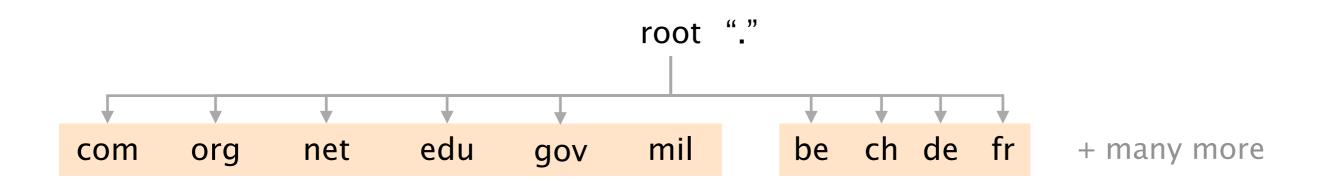
infrastructure hierarchy of DNS servers

naming structure

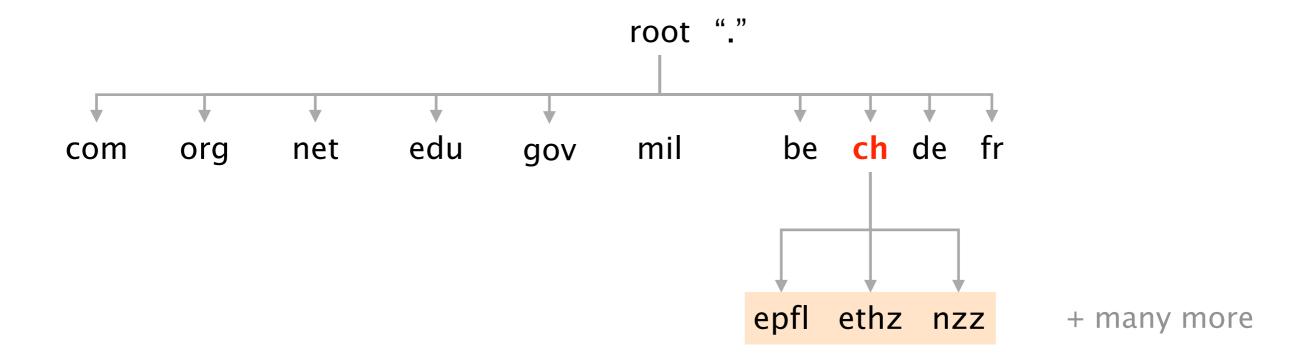
hierarchy of addresses

https://www.ee.ethz.ch/de/departement/

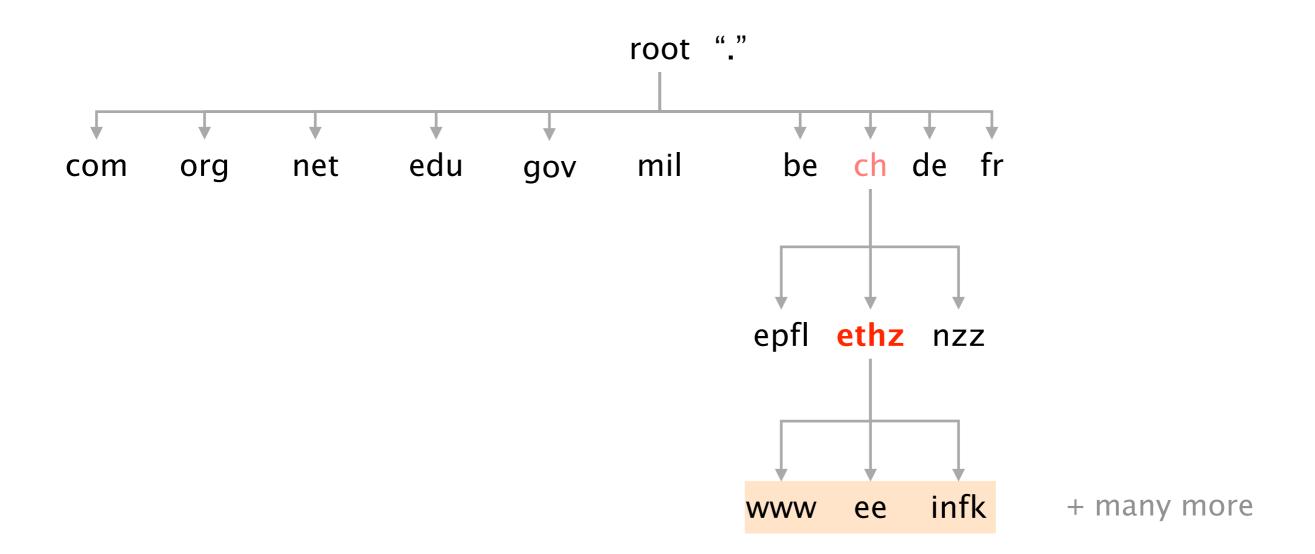
Top Level Domain (TLDs) sit at the top



Domains are subtrees



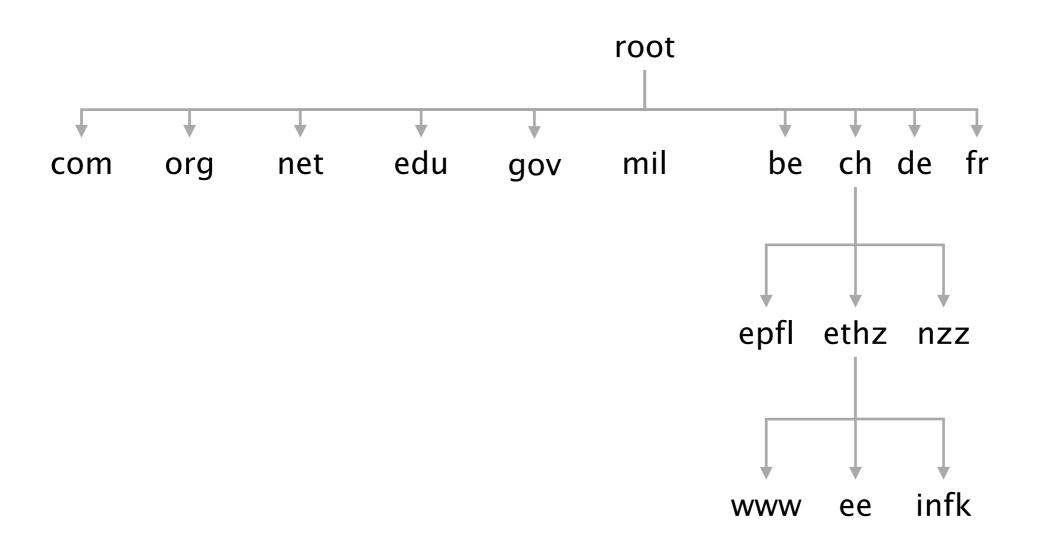
A name, *e.g.* ee.ethz.ch, represents a leaf-to-root path in the hierarchy

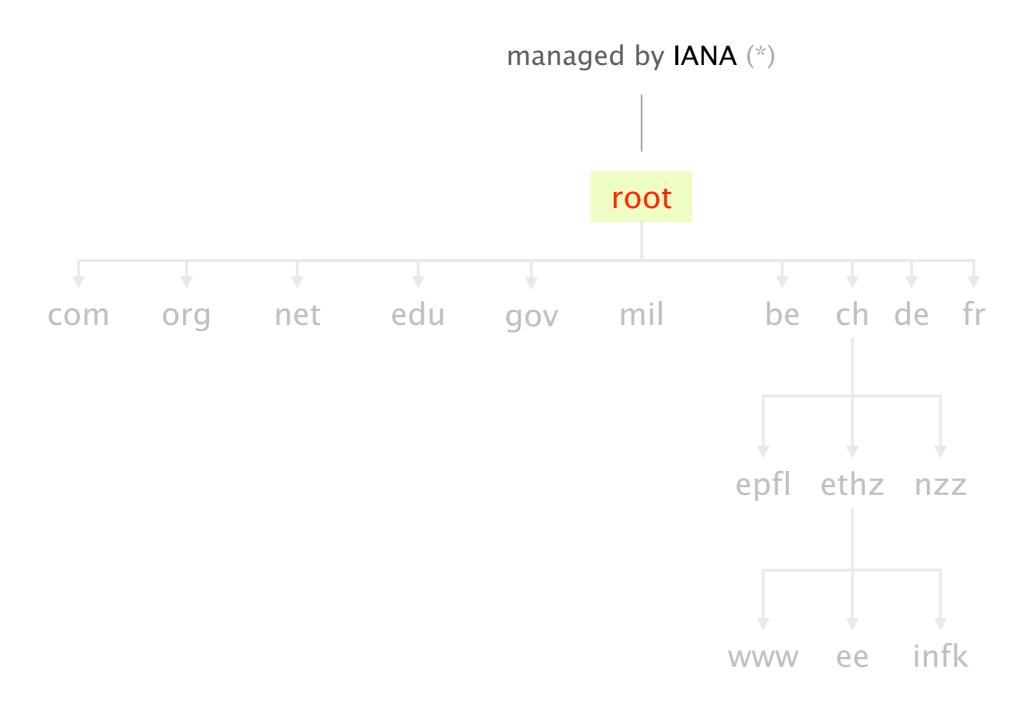


management

hierarchy of authority over names

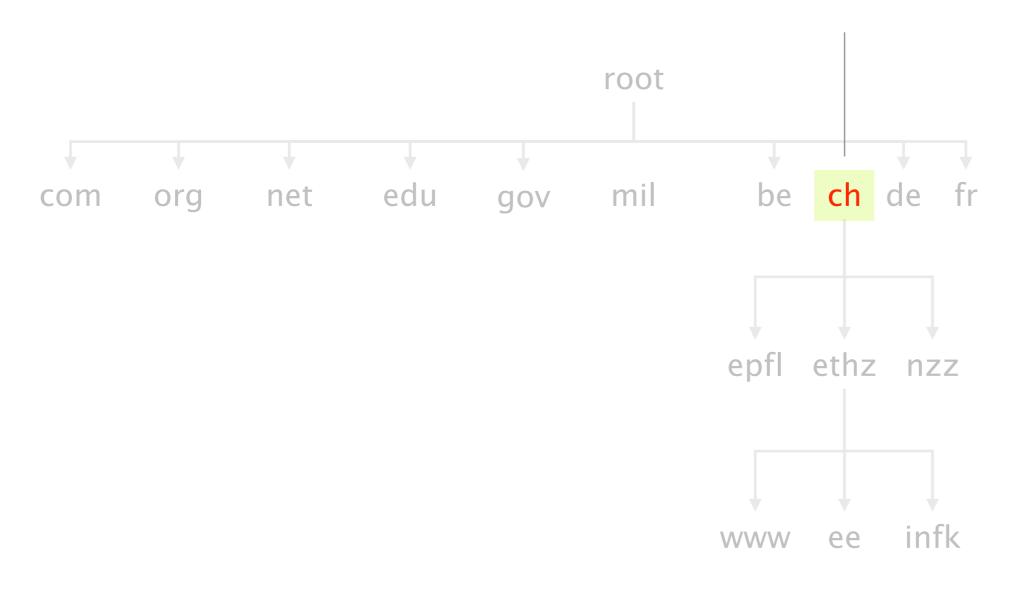
The DNS system is hierarchically administered



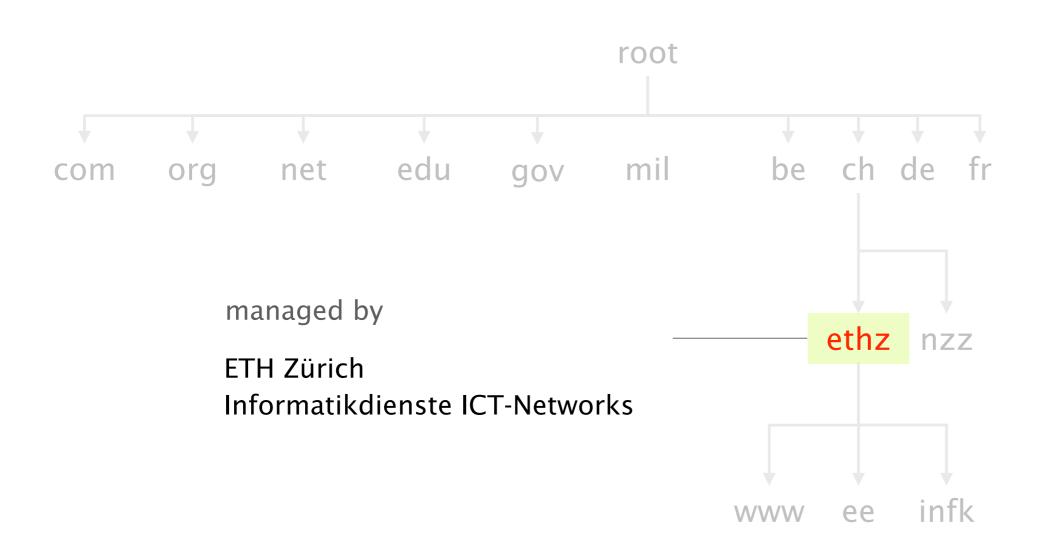


(*) see http://www.iana.org/domains/root/db

managed by The Swiss Education & Research Network (*)



(*) see https://www.switch.ch/about/id/

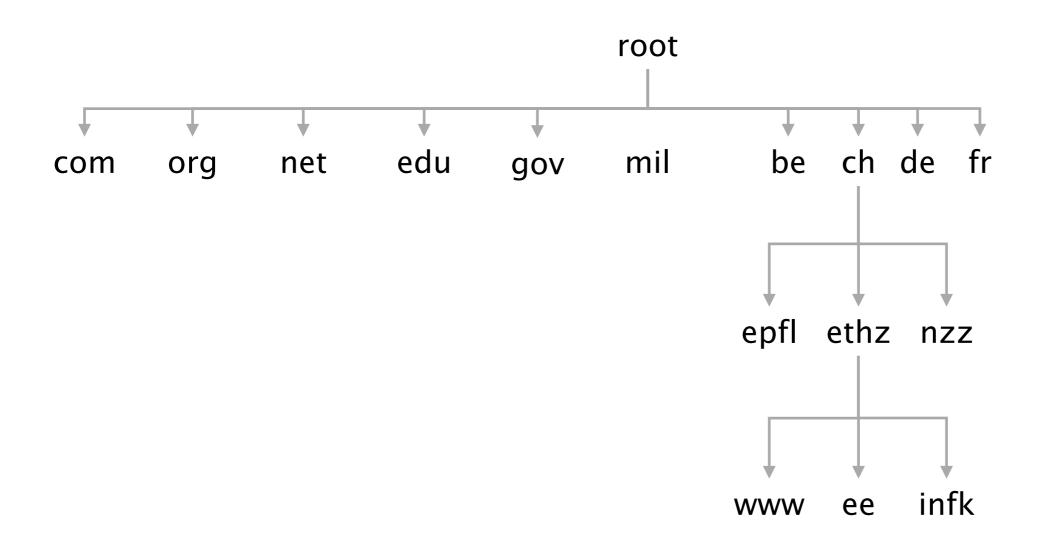


Hierarchical administration means that name collision is trivially avoided

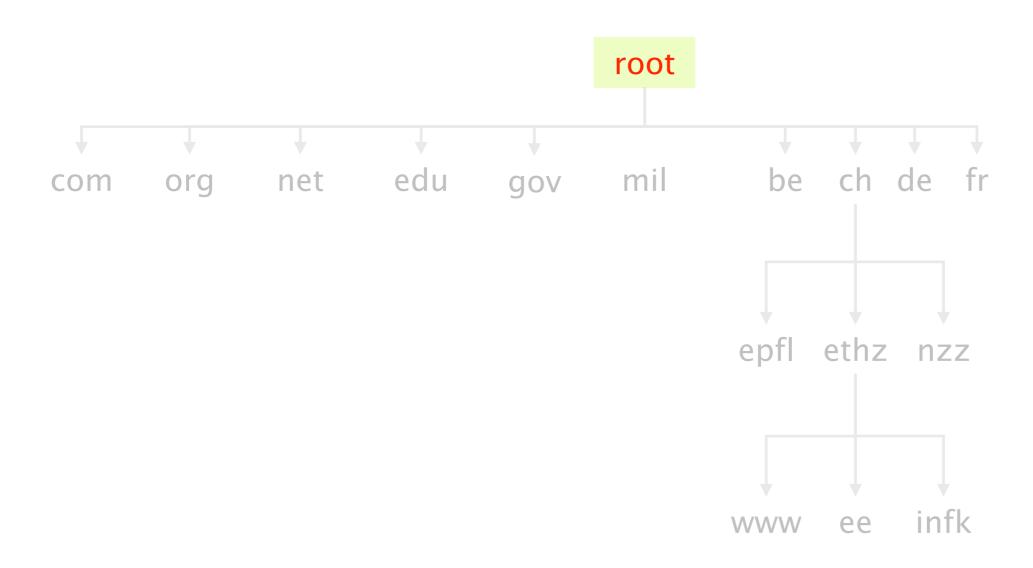
infrastructure

hierarchy of DNS servers

The DNS infrastructure is hierarchically organized



13 root servers (managed professionally) serve as root (*)



(*) see http://www.root-servers.org/

a. root-servers.net VeriSign, Inc.

b. root-servers.net University of Southern California

c. root-servers.net Cogent Communications

d. root-servers.net University of Maryland

e. root-servers.net NASA

f. root-servers.net Internet Systems Consortium

g. root-servers.net US Department of Defense

h. root-servers.net US Army

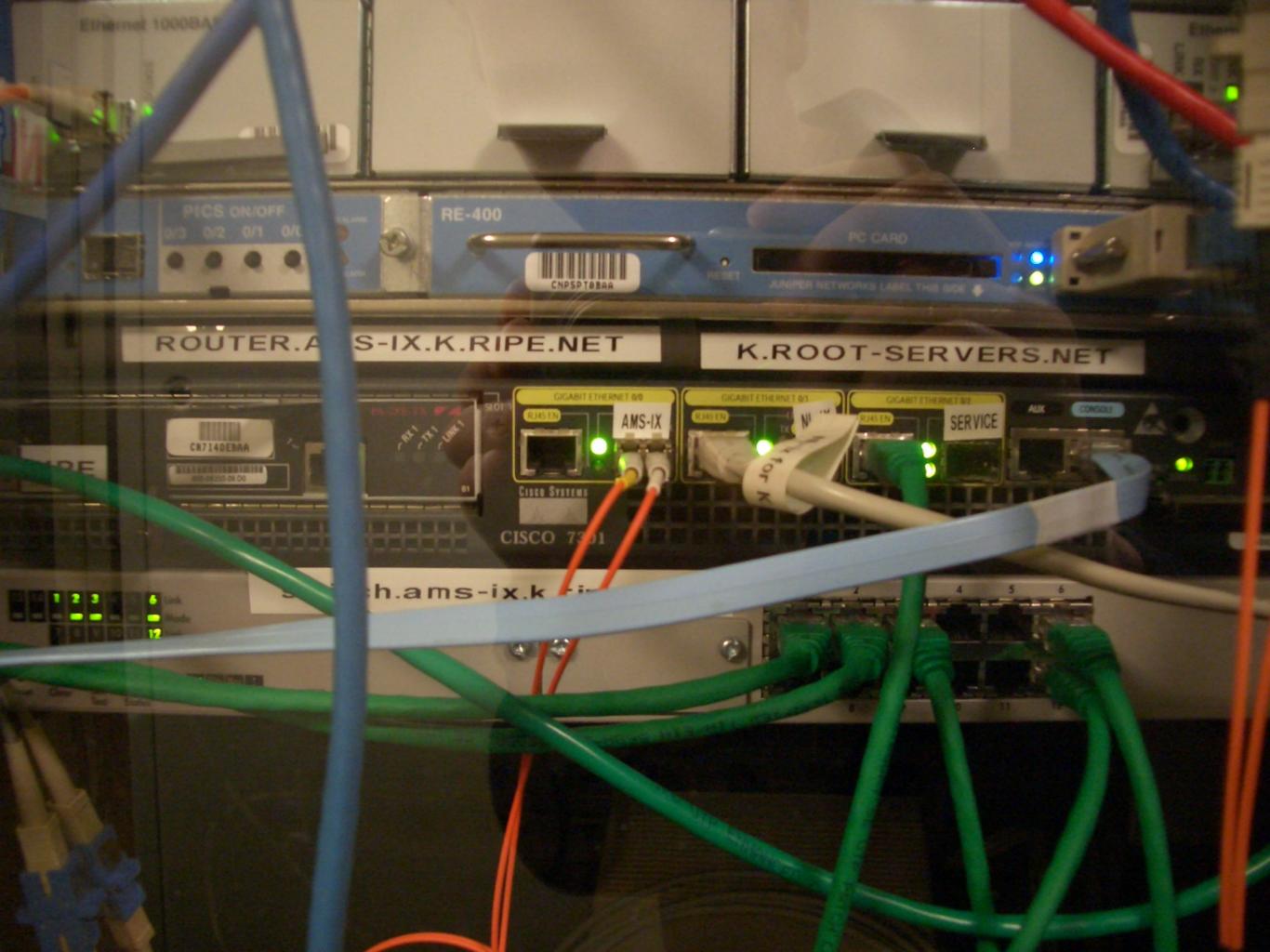
root-servers.net
 Netnod

. root-servers.net VeriSign, Inc.

k. root-servers.net RIPE NCC

I. root-servers.net ICANN

m. root-servers.net WIDE Project



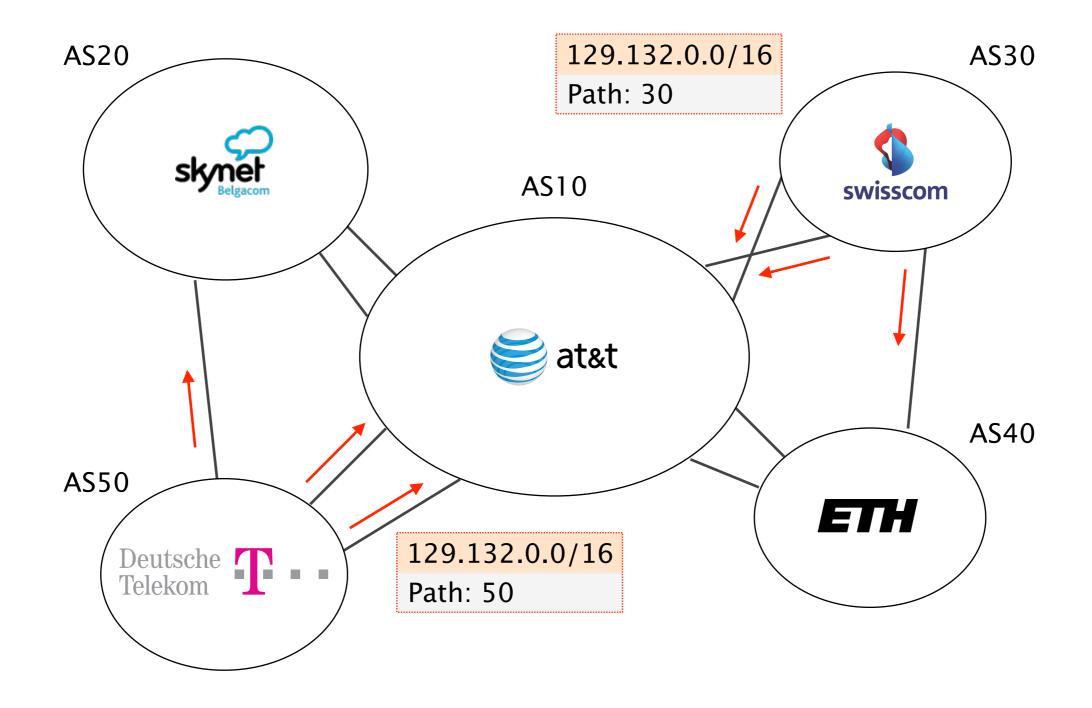
To scale root servers, operators rely on BGP anycast

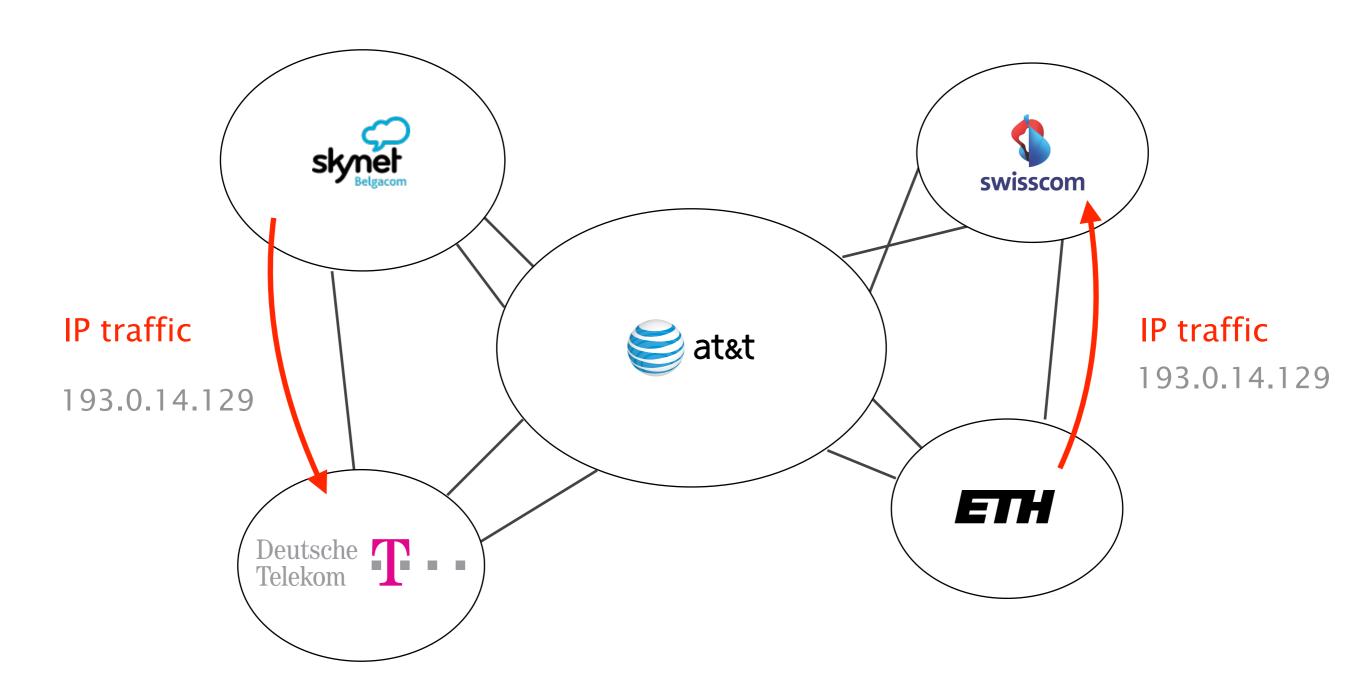
Intuition

Routing finds shortest-paths

If several locations announce the same prefix, then routing will deliver the packets to the "closest" location

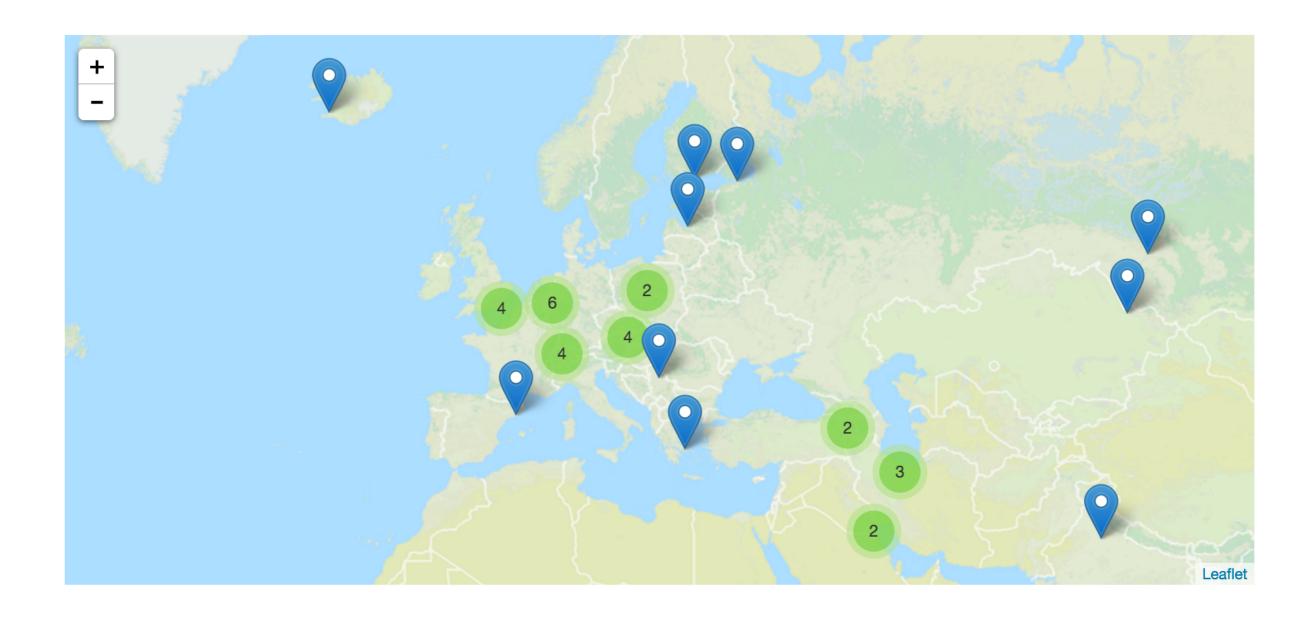
This enables seamless replications of resources





Do you see any problems in performing load-balancing this way?

Instances of the k-root server (*) are hosted in more than 80 locations worldwide



Two of these locations are in Switzerland: in Zürich and in Geneva



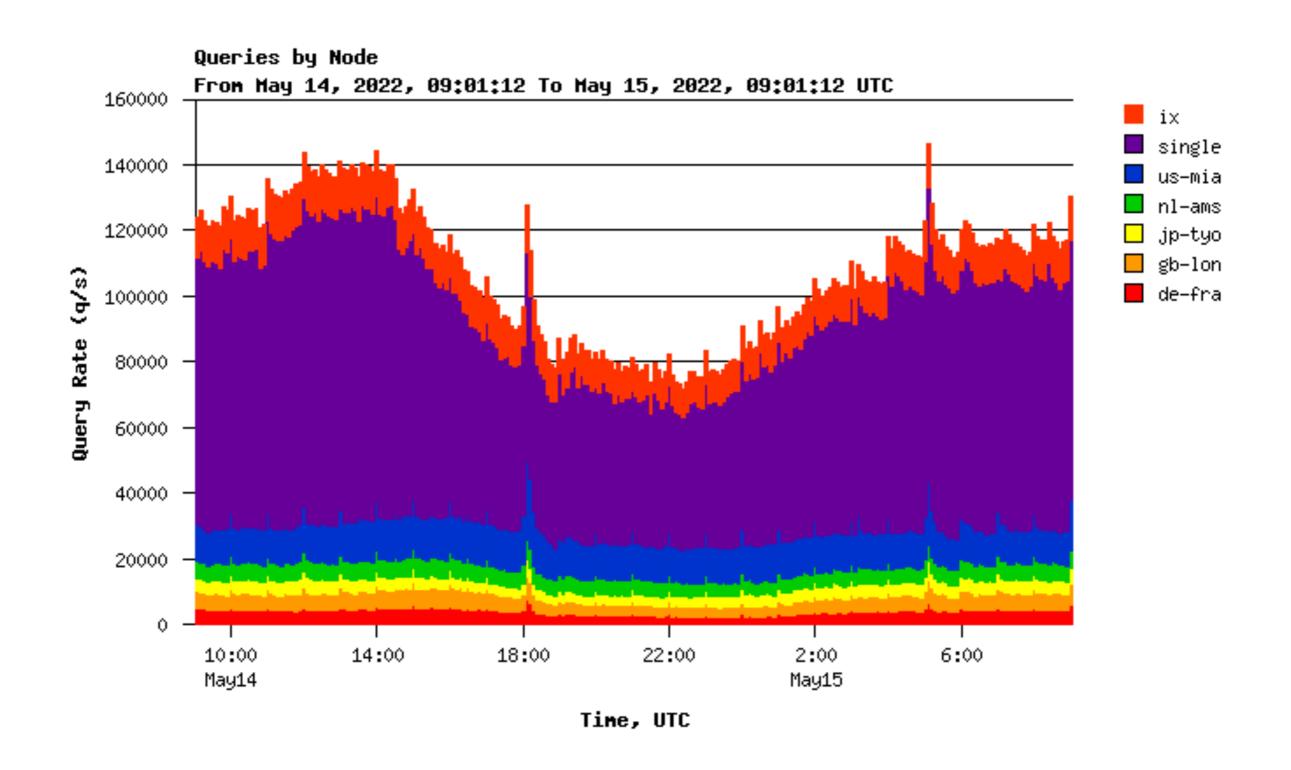
All locations announce 193.0.14.0/23 in BGP, with 193.0.14.129 being the IP of the server

Two of these locations are in Switzerland: in Zürich and in Geneva

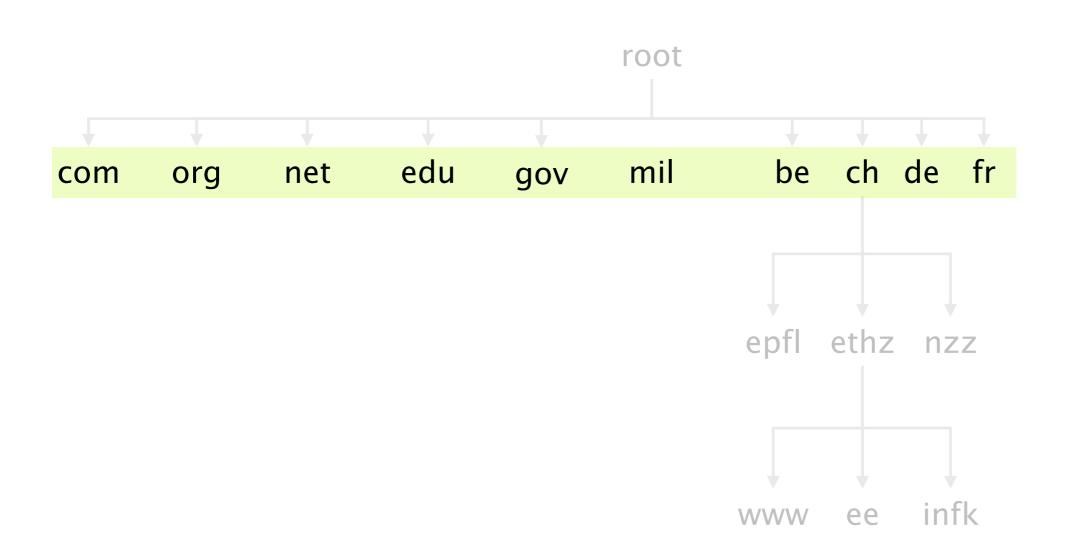
Do you mind guessing which one we use, here... in Zürich?

Each instance receives up to 80k queries per second

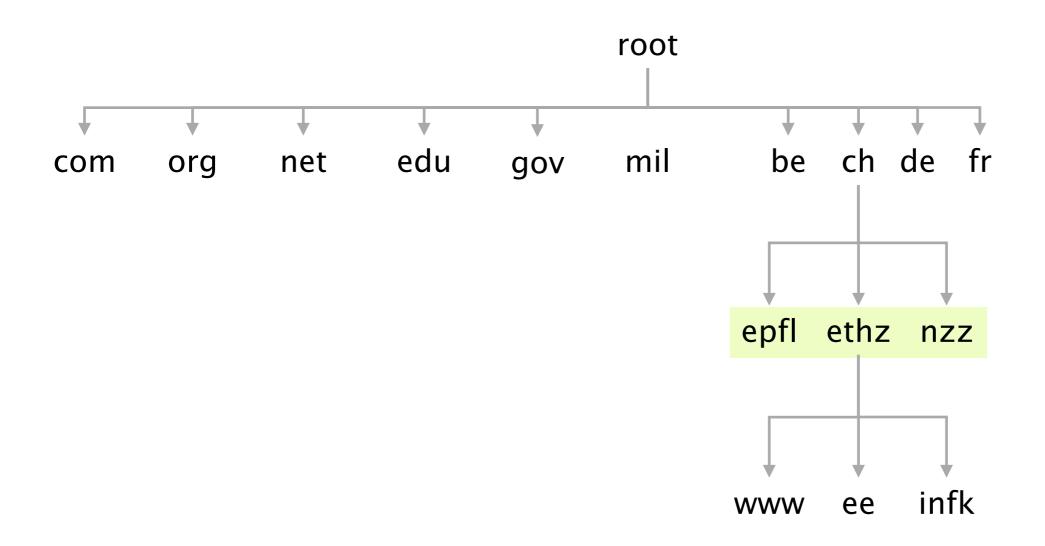
summing up to a few billions queries per day



TLDs server are also managed professionally by private or non-profit organization



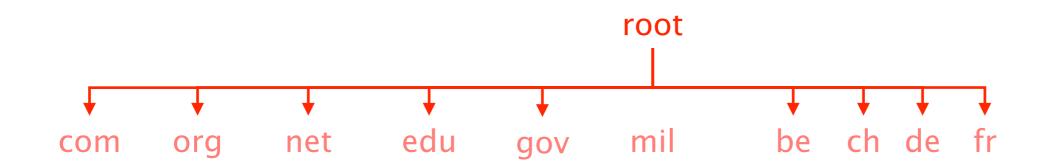
The bottom (and bulk) of the hierarchy is managed by Internet Service Provider or locally



Every server knows the address of the root servers (*) required for bootstrapping the systems

(*) see https://www.internic.net/domain/named.root

Each root server knows the address of all TLD servers



lvanbever:~\$ dig @a.root-servers.net ch.

ch.	172800	IN	NS	a.nic.ch.
ch.	172800	IN	NS	b.nic.ch.
ch.	172800	IN	NS	c.nic.ch.
ch.	172800	IN	NS	d.nic.ch.
ch.	172800	IN	NS	e.nic.ch.
ch.	172800	IN	NS	f.nic.ch.
ch.	172800	IN	NS	h.nic.ch.

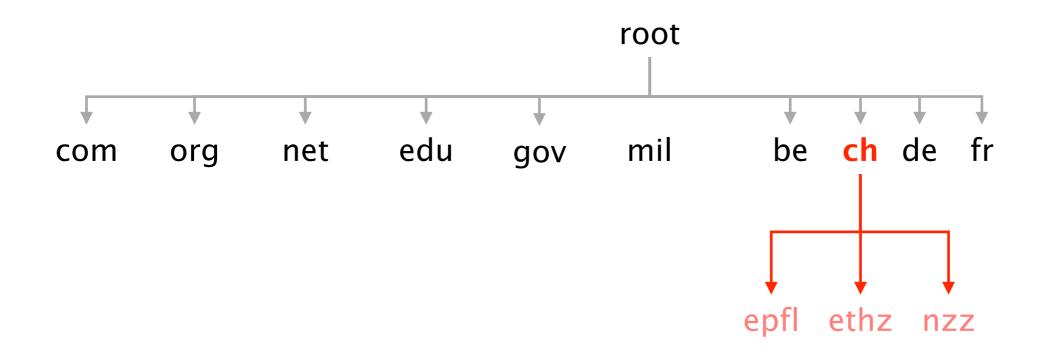
also see https://www.iana.org/domains/root/db/ch.html

If you want to learn more on ".ch" take a look at SWITCH's annual report



https://www.nic.ch/export/shared/.content/files/SWITCH_Report_Registry_2021.pdf

Any .ch DNS server knowns the addresses of the DNS servers of all sub-domains



lvanbever:~\$ dig @a.nic.ch ethz.ch

```
ethz.ch. 3600 IN NS ns2.ethz.ch. ethz.ch. 3600 IN NS ns1.ethz.ch.
```

Once arrived at the leaf of the hierarchy (*.ethz.ch), each DNS server knows the IP address of all children

lvanbever:~\$ dig @ns1.ethz.ch comm-net.ethz.ch

comm-net.ethz.ch. 3600 IN CNAME virt07.ethz.ch. virt07.ethz.ch. 3600 IN A 82.130.102.71

To scale, DNS adopt three intertwined hierarchies

naming structure addresses are hierarchical

https://www.ee.ethz.ch/de/departement/

management hierarchy of authority

over names

infrastructure hierarchy of DNS servers

To ensure availability, each domain must have at least a primary and secondary DNS server

Ensure name service availability as long as one of the servers is up

DNS queries can be load-balanced across the replicas

On timeout, client use alternate servers exponential backoff when trying the same server

Overall, the DNS system is highly scalable, available, and extensible

scalable #names, #updates, #lookups, #users,

but also in terms of administration

available domains replicate independently

of each other

extensible any level (including the TLDs)

can be modified independently

You've founded next-startup.ch and want to host it yourself, how do you insert it into the DNS?

You register next-startup.ch at a registrar *X e.g.* Swisscom, OVH, GoDaddy and *many* more

Provide *X* with the name and IP of your DNS servers *e.g.*, [ns1.next-startup.ch,129.132.19.253]

You set-up a DNS server @129.132.19.253 define A records for www, MX records for next-startup.ch...

A DNS server stores Resource Records composed of a (name, value, type, TTL)

Records Name Value

A hostname IP address

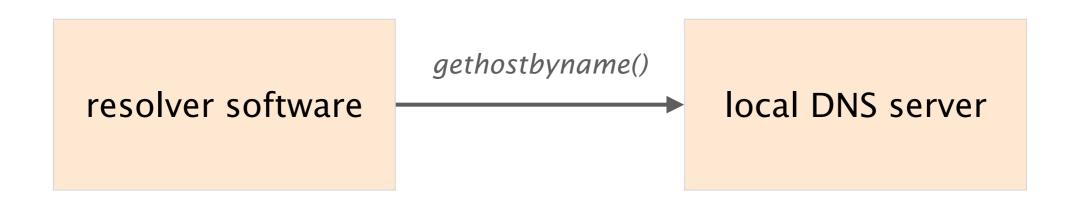
NS domain DNS server name

MX domain Mail server name

CNAME alias canonical name

PTR IP address corresponding hostname

Using DNS relies on two components



trigger resolution process send request to local DNS server usually, near the endhosts

configured statically (resolv.conf)

or dynamically (DHCP)

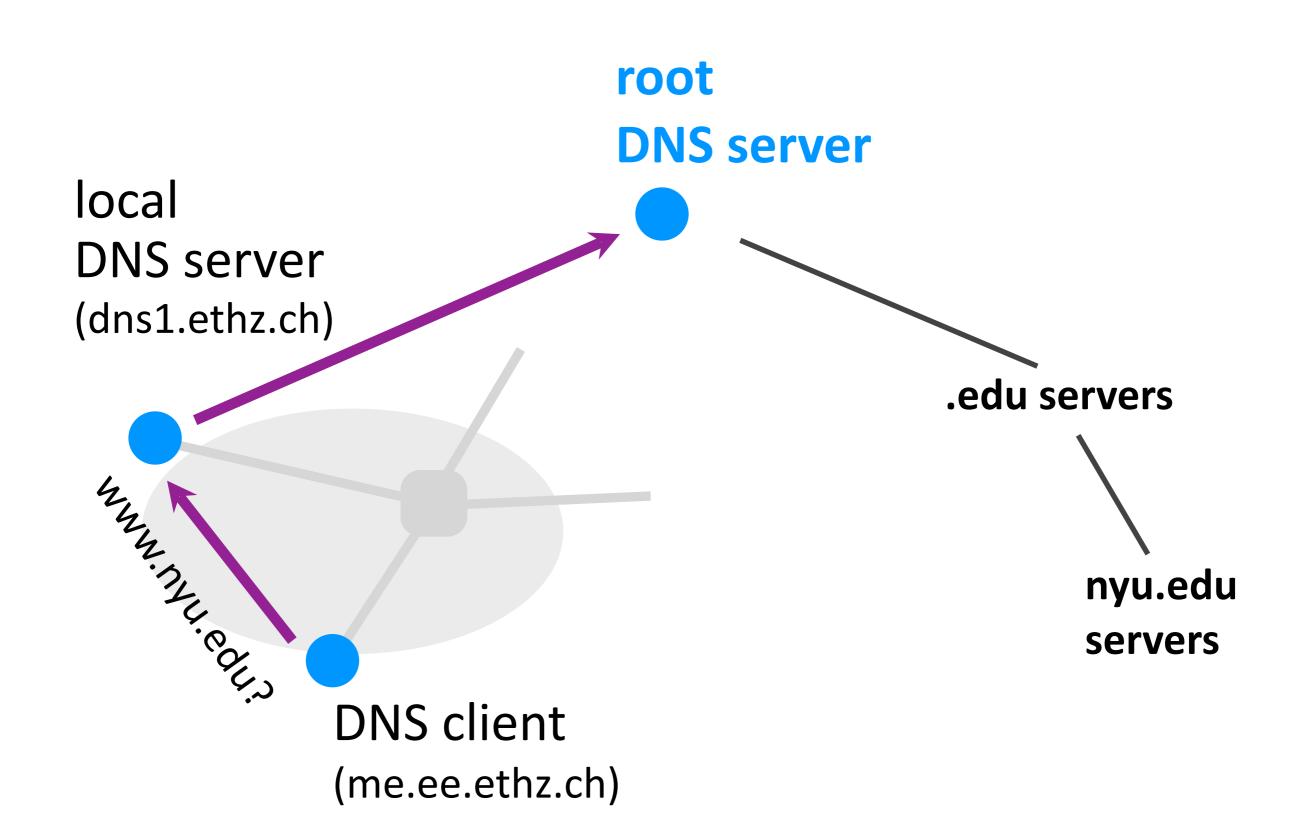
DNS query and reply uses UDP (port 53), reliability is implemented by repeating requests (*)

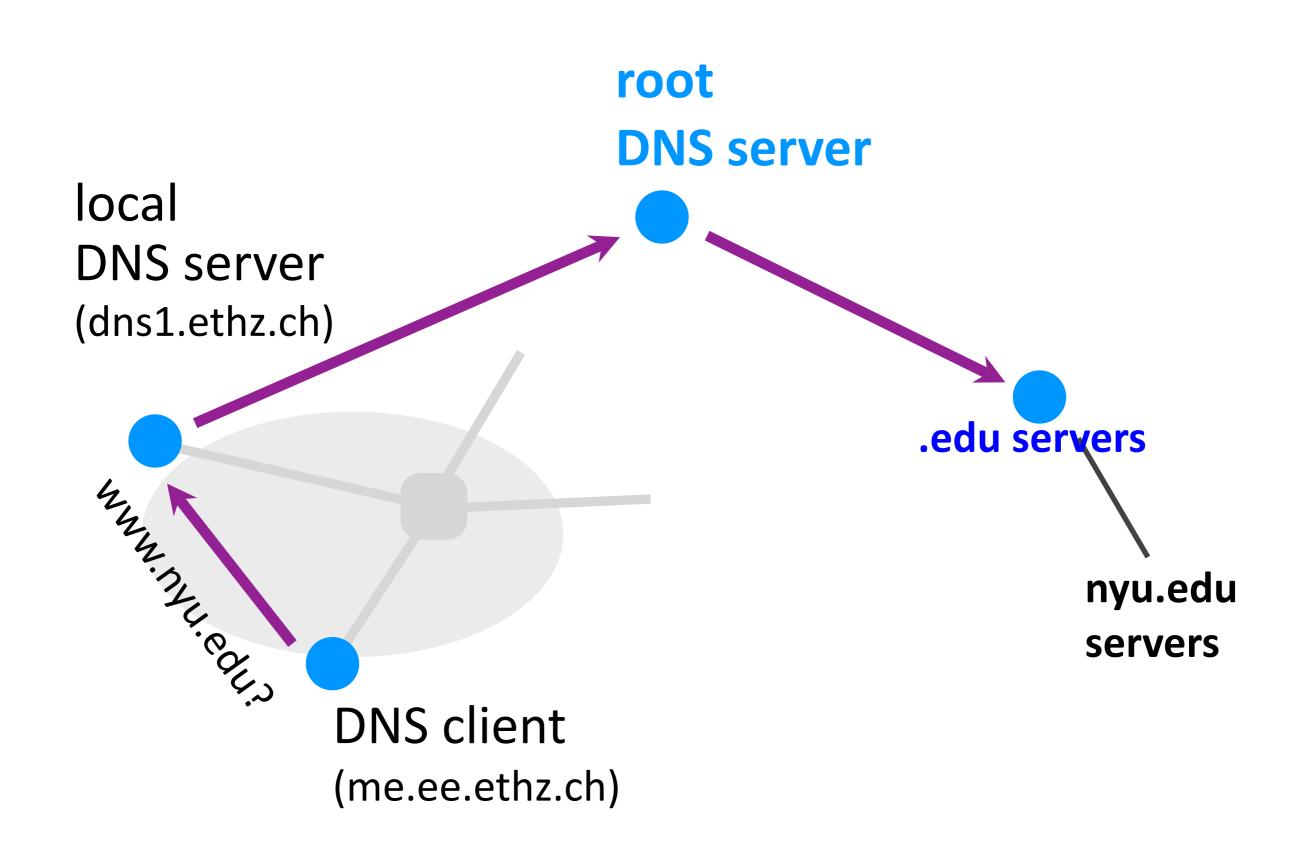
^(*) see Book (Section 5)

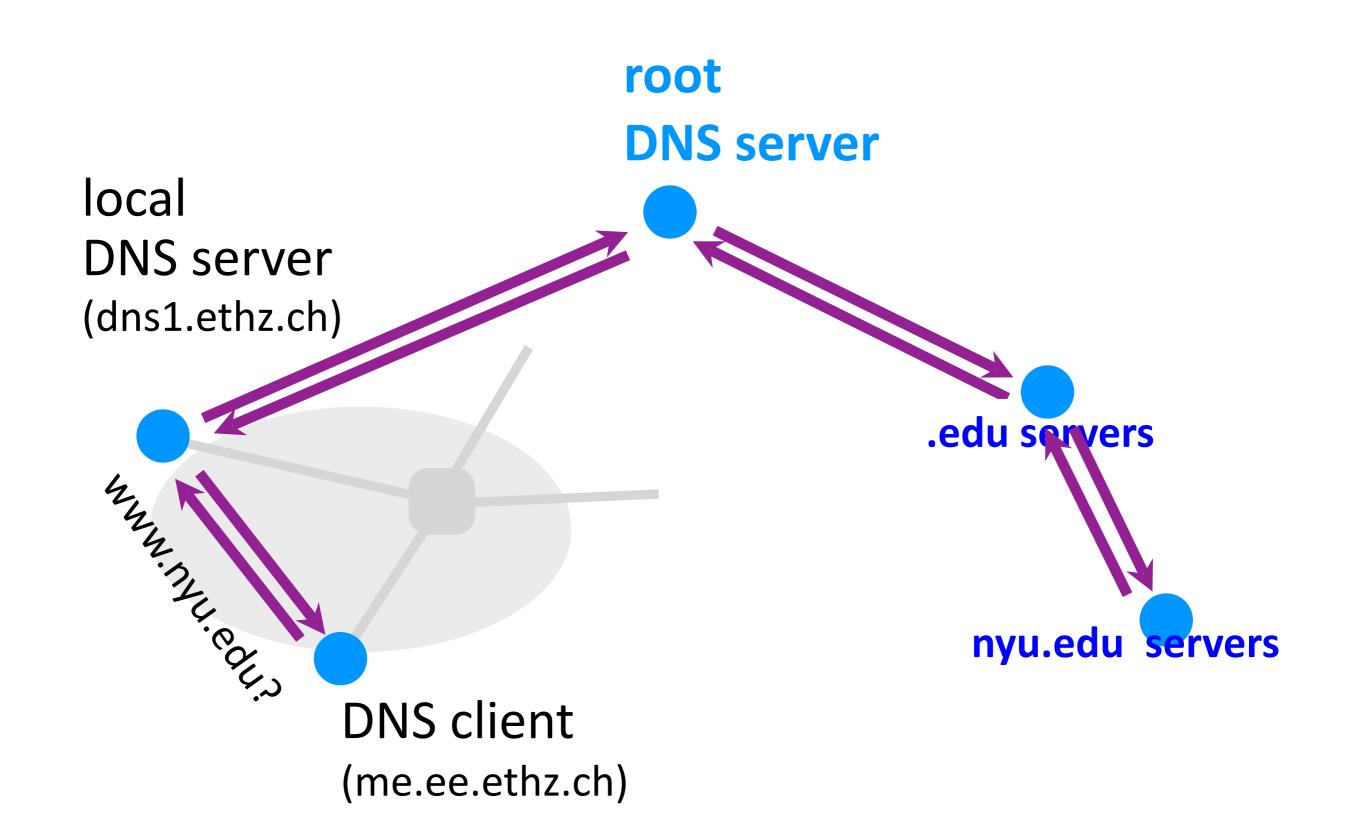
DNS resolution can either be recursive or iterative

When performing a recursive query, the client offload the task of resolving to the server local root servers **DNS** server (dns1.ethz.ch) .edu servers nyu.edu servers **DNS** client

(me.ee.ethz.ch)





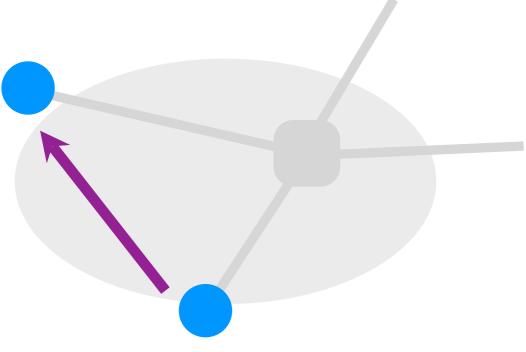


When performing a iterative query, the server only returns the address of the "next server"

root **DNS** server



local DNS server



DNS client (me.ee.ethz.ch)





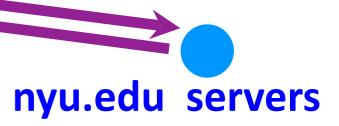


local Where is .edu? DNS server

Where is nyu.edu?



Where is www.nyu.edu?



DNS client (me.ee.ethz.ch)

What about resolving speeds?
Waiting for servers all over the globe is not fast...

To reduce resolution times, DNS relies on caching

DNS servers cache responses to former queries and your client and the applications (!)

Authoritative servers associate a lifetime to each record Time-To-Live (TTL)

DNS records can only be cached for TTL seconds after which they must be cleared

As top-level servers rarely change & popular website visited often, caching is very effective (*)

Top 10% of names account for 70% of lookups

9% of lookups are unique

Limit cache hit rate to 91%

Practical cache hit rates ~75%

Congestion Control

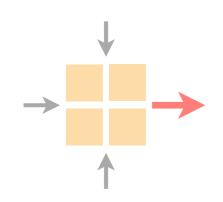
DNS

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reliable transport starts *today*!

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