Communication Networks Spring 2022



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https://comm-net.ethz.ch/

ETH Zürich

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Communication Networks Exercise 3



General information

Last week's exercise

Important lecture topics

Introduction to this week's exercise

Time to solve the exercise

Group registration is open

Please register your groups for the routing project: <u>https://comm-net.ethz.ch/registration/php/index.php</u>

Use the #group_search channel on Slack if you look for other group members

Let us know via Slack or email if there are any problems

GitLab setup

We will give each group a GitLab repository

If you did not already use the ETH GitLab in the past, please log in once: <u>https://gitlab.ethz.ch</u>

More information to the GitLab usage once the project starts

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Task 2.2 - Changing Weights



Link weight from A to C can be different than from C to A

Weights are dynamic and always represent the link load

Tie-breaking: path with lower (alphabetically) next hop

Network topology with directional link weights.

e is much bigger than 1



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				Link	Load				Ne	Next Ho B C				
	$A \rightarrow B$	$A \rightarrow C$	$B \rightarrow A$	$B \rightarrow D$	$C \rightarrow A$	$C \rightarrow D$	$D \rightarrow B$	$D \rightarrow C$	В	С	E			
0	0	0	0	0	0	0	0	0	А	А	B			
1	0	0	1 + e	0	1	0	е	0						

D

В





First, find the current next hops



				Link	Load				Ne	ext F
	$A \rightarrow B$	$A \rightarrow C$	$B \rightarrow A$	$B \rightarrow D$	$C \rightarrow A$	$C \rightarrow D$	$D \rightarrow B$	$D \rightarrow C$	В	С
0	0	0	0	0	0	0	0	0	А	А
1	0	0	1 + e	0	1	0	e	0		

Ne	ext H	op
В	С	D
А	А	В







First, find the current next hops

				Link	Load				Next Hop			
	$A \rightarrow B$	$A \rightarrow C$	$B \rightarrow A$	$B \rightarrow D$	$C \rightarrow A$	$C \rightarrow D$	$D \rightarrow B$	$D \rightarrow C$	В	С	D	
0	0	0	0	0	0	0	0	0	Α	А	В	
1	0	0	1 + e	0	1	0	е	0	D	Α	C	





Then update the current link loads



2+e

				Link	Load				Ne	ор	
	$A \rightarrow B$	$A \rightarrow C$	$B \rightarrow A$	$B \rightarrow D$	$C \rightarrow A$	$C \rightarrow D$	$D \rightarrow B$	$D \rightarrow C$	B	С	D
0	0	0	0	0	0	0	0	0	A	А	B
1	0	0	1 + e	0	1	0	e	0	D	Α	C
2	0	0	0	1	2 + e	0	0	1 + e			



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Once again find the next hops



				Link	Load				Ne	op	
	$A \rightarrow B$	$A \rightarrow C$	$B \rightarrow A$	$B \rightarrow D$	$C \rightarrow A$	$C \rightarrow D$	$D \rightarrow B$	$D \rightarrow C$	В	C	D
0	0	0	0	0	0	0	0	0	А	А	В
1	0	0	1 + e	0	1	0	e	0	D	Α	C
2	0	0	0	1	2 + e	0	0	1 + e	<u> </u>		<u>.</u>











Once again find the next hops

				Link	Load				Next Ho		
	$A \rightarrow B$	$A \rightarrow C$	$B \rightarrow A$	$B \rightarrow D$	$C \rightarrow A$	$C \rightarrow D$	$D \rightarrow B$	$D \rightarrow C$	В	С	D
0	0	0	0	0	0	0	0	0	Α	А	B
1	0	0	1 + e	0	1	0	e	0	D	А	C
2	0	0	0	1	2 + e	0	0	1 + e	Α	D	B



2+e



Update the link loads



				Link	Load				Ne	Next Hop B C D A A B			
	$A \rightarrow B$	$A \rightarrow C$	$B \rightarrow A$	$B \rightarrow D$	$C \rightarrow A$	$C \rightarrow D$	$D \rightarrow B$	$D \rightarrow C$	В	C	D		
0	0	0	0	0	0	0	0	0	А	A	В		
1	0	0	1 + e	0	1	0	e	0	D	Α	C		
2	0	0	0	1	2 + e	0	0	1 + e	Α	D	В		
3	0	0	2 + e	0	0	1	1 + e	0					



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Continue in the same way



				Link	Load				Ne	ext H	op
	$A \rightarrow B$	$A \rightarrow C$	$B \rightarrow A$	$B \rightarrow D$	$C \rightarrow A$	$C \rightarrow D$	$D \rightarrow B$	$D \rightarrow C$	В	C	D
0	0	0	0	0	0	0	0	0	A	A	В
1	0	0	1 + e	0	1	0	e	0	D	A	C
2	0	0	0	1	2 + e	0	0	1 + e	Α	D	B
3	0	0	2 + e	0	0	1	1 + e	0	D	Α	C
4	0	0	0	1	2 + e	0	0	1 + e	A	D	B

e — incoming traffic

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We already see the oscillation

In general, such a frequent weight adaptation is way too sensitive

е

All traffic anti-clockwise

Link Load $A \rightarrow C$ $B \rightarrow D$ $\mathsf{C} \to \mathsf{A}$ $C \rightarrow D$ $D \rightarrow C$ $A \rightarrow B$ $B \rightarrow A$ $D \rightarrow B$ 0 0 0 0 0 0 0 0 0 0 1 0 0 0 1 1 + e0 e 2 + e 2 0 0 1 0 0 0 1 + e 3 0 0 2 + e 0 1 + e 0 0 1 4 0 0 0 1 2 + e 0 0 1 + e

All traffic clockwise

Ne	Next Hop								
В	С	D							
А	А	В							
D	А	С							
Α	D	В	•						
D	А	С							
A	D	В							

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In the lecture we go through the layers bottom-up



Another possible approach would be top-down



We face a common problem

No matter the direction, often concepts of other layers are needed to understand the current one

Unfortunately, we cannot prevent that completely

We saw that when speaking about MAC addresses, suddenly we also care about IP addresses

MAC addresses identify sender and receiver adapters

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In general, we therefore use IP addresses (L3) to address arbitrary hosts

MAC addresses are then used on a hop-by-hop basis to eventually reach the corresponding host

In fact, for humans domain names are even easier to remember

domain name \rightarrow DNS (L5) \rightarrow IP (L3) \rightarrow ARP \rightarrow MAC (L2)of destinationof destinationof next hop

We currently only consider IP addresses which are reachable over a given link

> That simplifies the whole process, we only need to be able to translate from IP to MAC address

> > $IP (L3) \longrightarrow ARP \longrightarrow MAC (L2)$ of destination of destination

Who are you? IP-to-MAC binding Given an IP address reachable on a link, How do I find out what MAC to use?

Address Resolution Protocol

That can only work if hosts can get an IP address

Who am I? MAC-to-IP binding How do I acquire an IP address?

Dynamic Host Configuration Protocol

We will explore both concepts (ARP and DHCP) in today exercise

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Two more questions related to routing concepts

Task 3.1 Distance Vector



Compute shortest-paths using a distance vector algorithm

Tie-breaking: path with lower amount of links

Compared to link-state algorithms, paths are now computed in a distributed fashion

Task Pijkstra's Algithm with Link Failure



Back to Dijkstra (link-state)

We assume that the link between d and R3 fails

R3 detects that quickly but what about the other nodes?

What happens if the local network view does not match with the reality?

And three questions related to Ethernet & Switching



As a reminder, let's look at this simple example

A switch learns how to map MACs to ports



Switch learns how to map A to port 1



Dst D unknown: broadcast



Switch learns how to map D to port 4



Dst A known, no broadcast required



What happens if you have duplicated MAC addresses?

Task 3.4: Imposter

Put your knowledge about DHCP and ARP together

Who am I? MAC-to-IP binding How do I acquire an IP address? Dynamic Host Configuration Protocol (DHCP)

Who are you? IP-to-MAC binding Given an IP address reachable on a link, how do I find out what MAC to use? Address Resolution Protocol (ARP)

Task 3.5: MAC-Learning (exam question 2021)

We asked this question in the summer exam of 2021

You find the full exam (as well as other ones) on the website

Use your knowledge from task 3.3 to solve this one

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