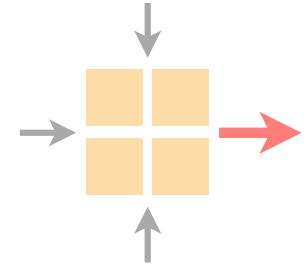


# Communication Networks

Spring 2020



Tobias Bühler

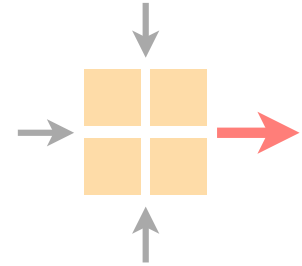
<https://comm-net.ethz.ch/>

ETH Zürich

March 5 2020

# Communication Networks

## Exercise 2



Quick Dijkstra repetition

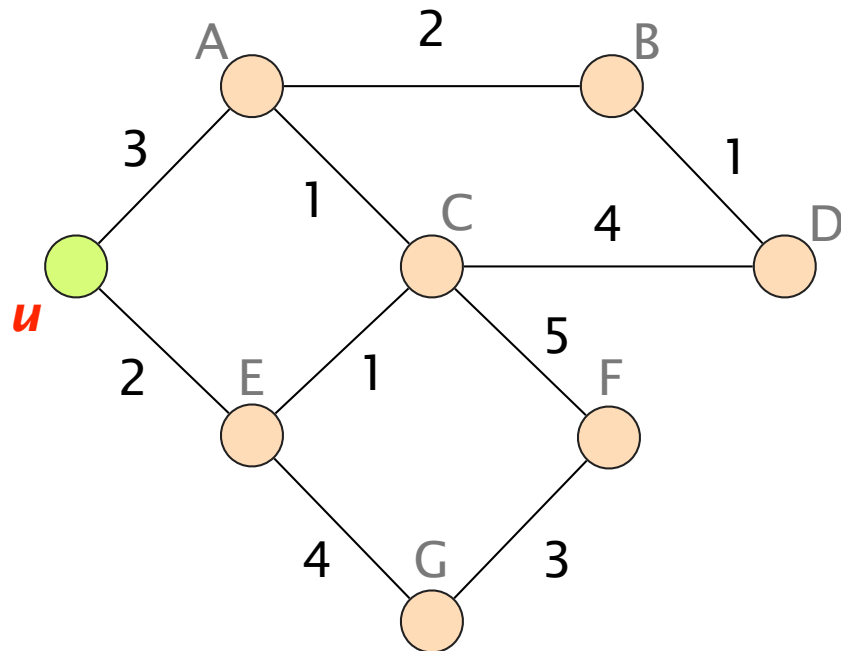
Overview current assignment

Questions?

Time for you to solve the tasks

Solutions will be published next week

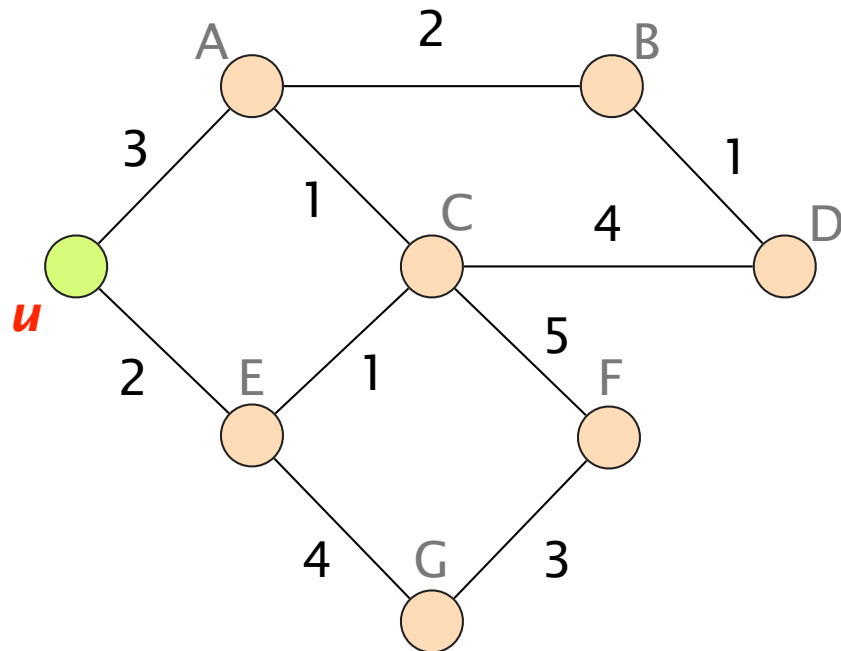
Let's compute the shortest-paths from  $u$  using Dijkstra's algorithm



$S = \{u\}$  — set of nodes for which we know the shortest-path

$D(v)$  — the smallest distance currently known by  $u$  to reach  $v$

$c(u, v)$  — the weight of the link connecting  $u$  and  $v$



### Initialization

$S = \{u\}$

for all nodes  $v$ :

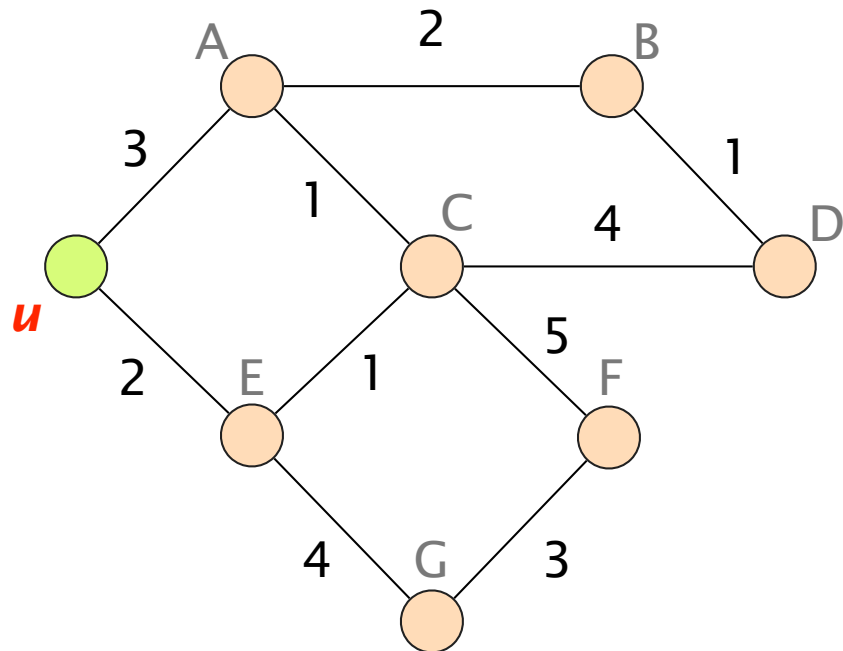
if ( $v$  is adjacent to  $u$ ):

$$D(v) = c(u, v)$$

else:

$$D(v) = \infty$$

D is initialized based on u's weight,  
and S only contains u itself



$D(.) =$

$S = \{u\}$

A 3

B  $\infty$

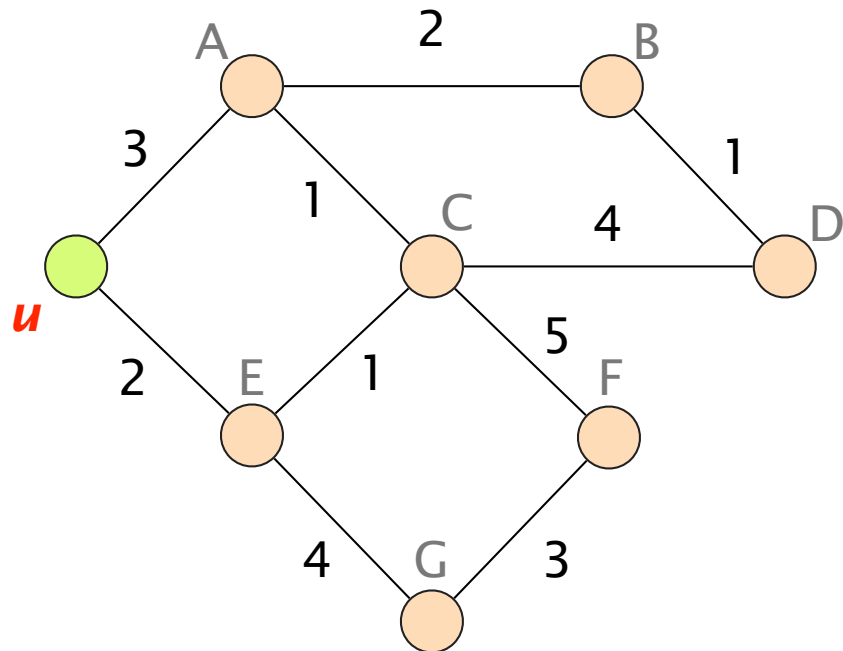
C  $\infty$

D  $\infty$

E 2

F  $\infty$

G  $\infty$



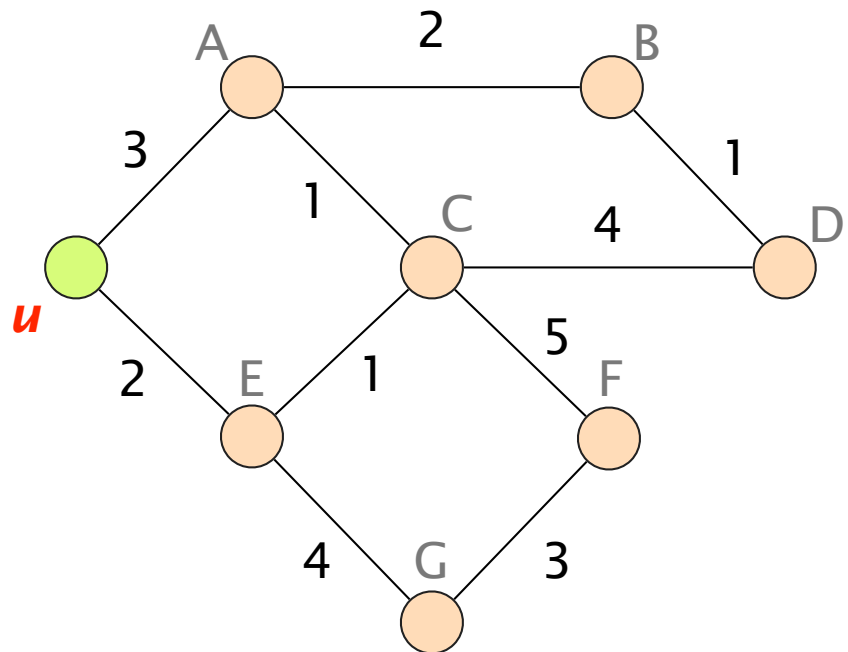
Loop

**while** *not* all nodes in S:

**add**  $w$  with the smallest  $D(w)$  to S

**update**  $D(v)$  for all adjacent  $v$  not in S:

$$D(v) = \min\{D(v), D(w) + c(w, v)\}$$



$D(.) =$

$S = \{u\}$

A 3

B  $\infty$

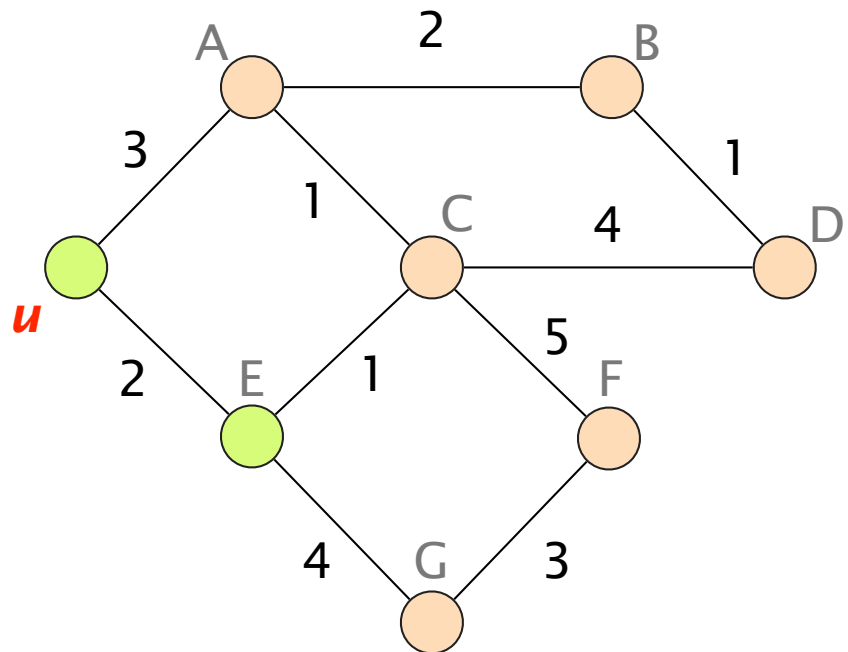
C  $\infty$

D  $\infty$

E 2 — smallest  $D(w)$

F  $\infty$

G  $\infty$



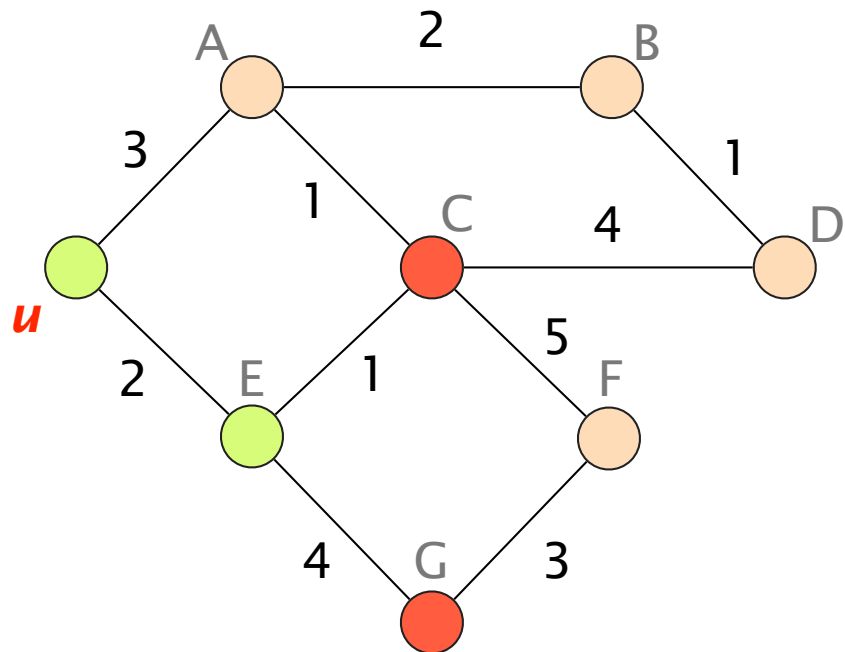
$D(.) =$

A	3
B	$\infty$
C	$\infty$
D	$\infty$
E	2
F	$\infty$
G	$\infty$

add E to S

$S = \{u, E\}$





$D(.) =$

$S = \{u, E\}$

A 3

B  $\infty$

C 3 —  $D(v) = \min\{\infty, 2 + 1\}$

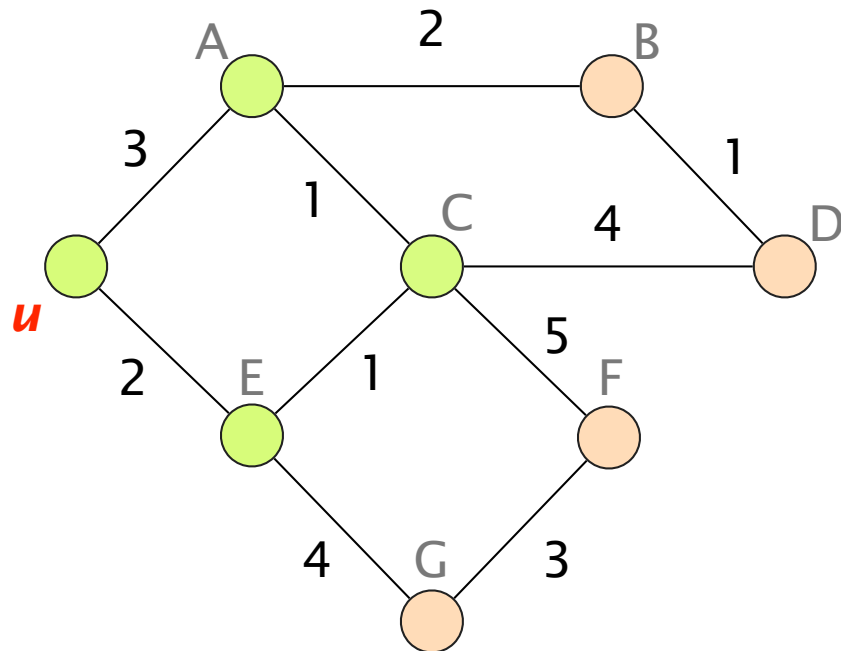
D  $\infty$

E 2

F  $\infty$

G 6 —  $D(v) = \min\{\infty, 2 + 4\}$

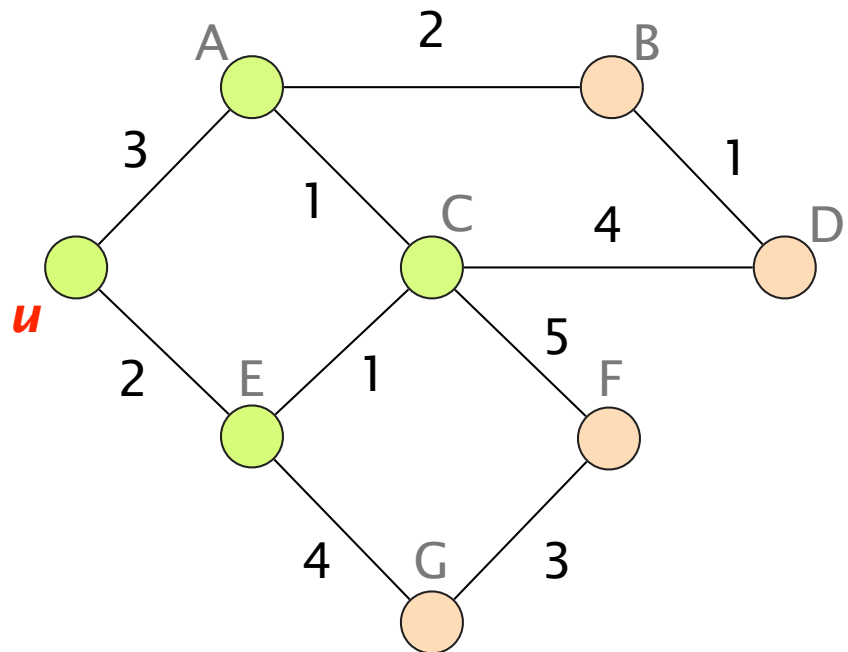
Skipping a few steps...



$D(.) =$

A	3
B	5
C	3
D	7
E	2
F	8
G	6

$S = \{u, E, A, C\}$



$D(.) =$

$S = \{u, E, A, C\}$

A 3

B 5

C 3

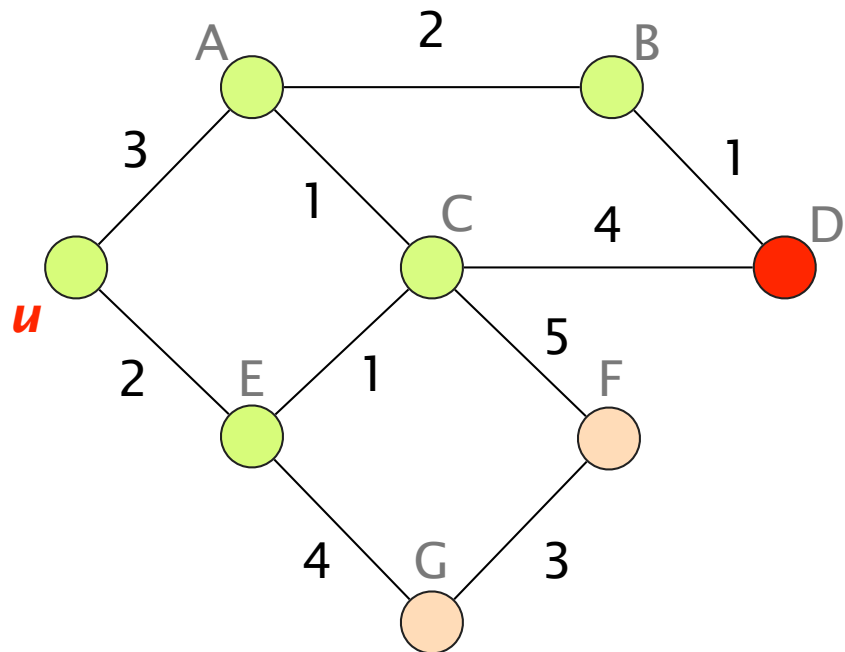
D 7

E 2

F 8

G 6

— smallest  $D(w)$



$D(.) =$

$S = \{u, E, A, C, B\}$

A 3

B 5

C 3

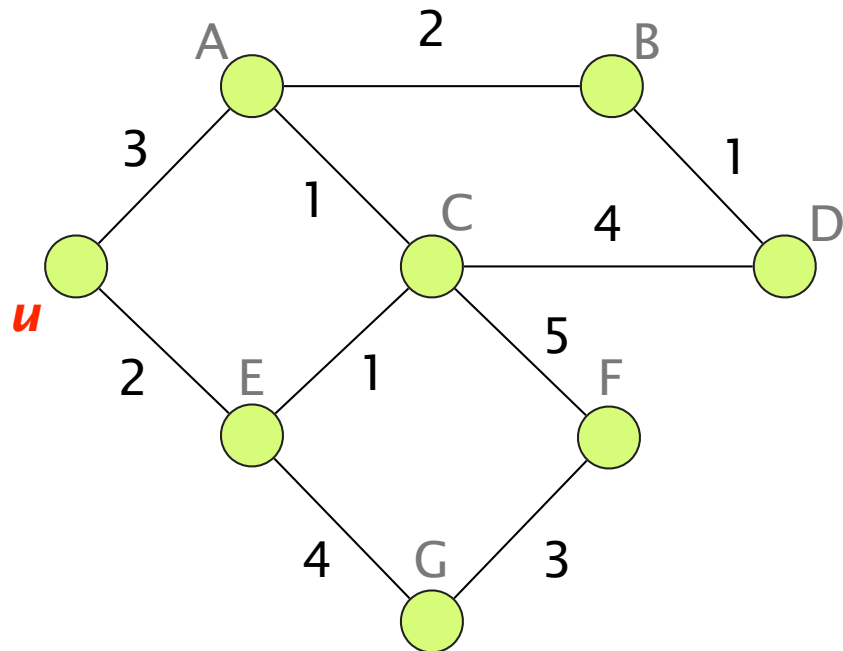
D 6 —  $D(v) = \min\{7, 5 + 1\}$

E 2

F 8

G 6

Here is the final state

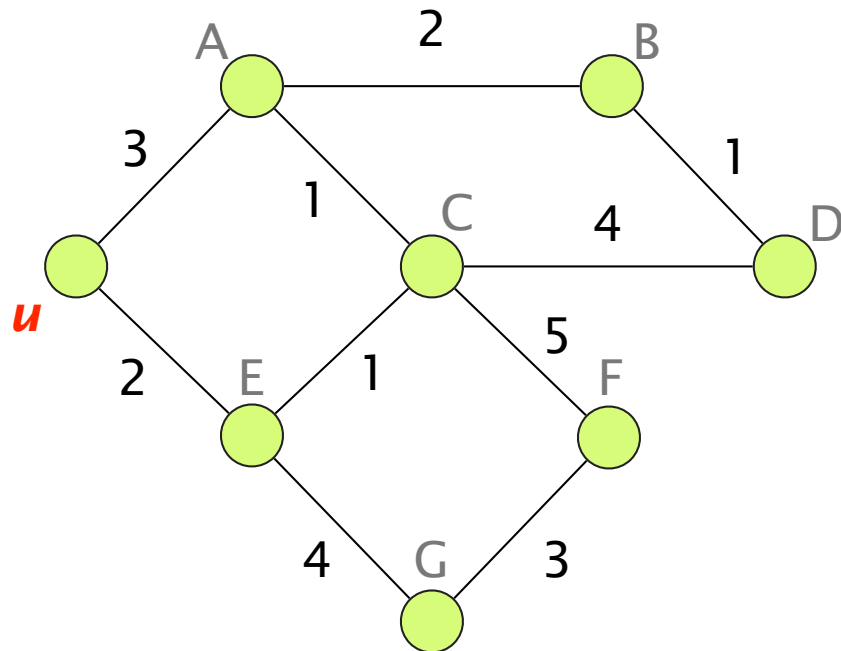


$D(.) =$

A	3
B	5
C	3
D	6
E	2
F	8
G	6

$S = \{u, A,$   
B, C, D, E,  
F, G}

From the shortest-paths,  
 $u$  can directly compute its forwarding table

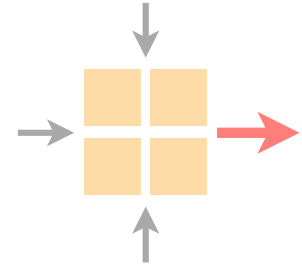


### Forwarding table

<i>destination</i>	<i>next-hop</i>
A	A
B	A
C	E
D	A
E	E
F	E
G	E

# Communication Networks

## Exercise 2



Quick Dijkstra repetition

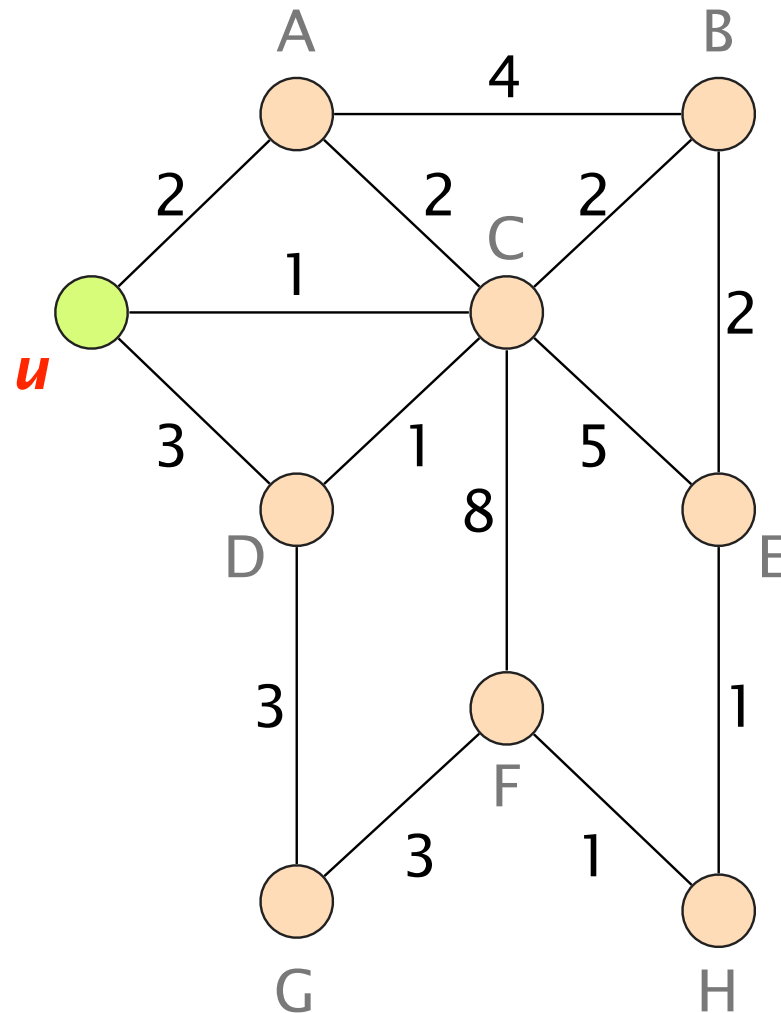
Overview current assignment

Questions?

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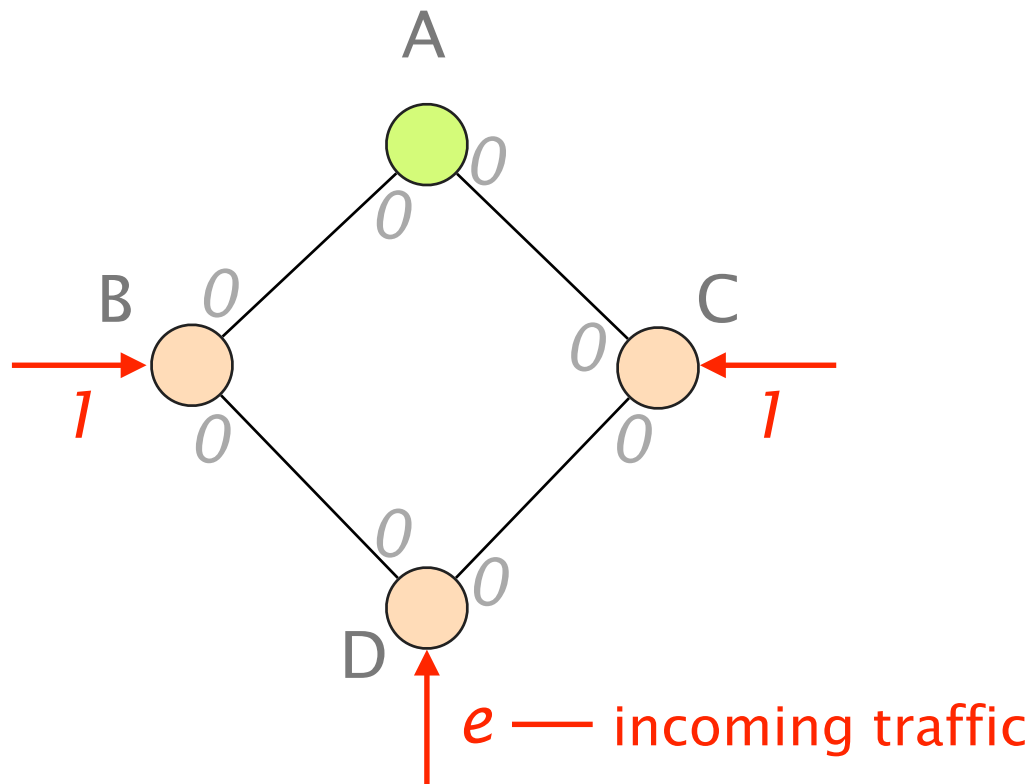
Solutions will be published next week

## Task 1: Dijkstra's algorithm





## Task 2: Changing weights

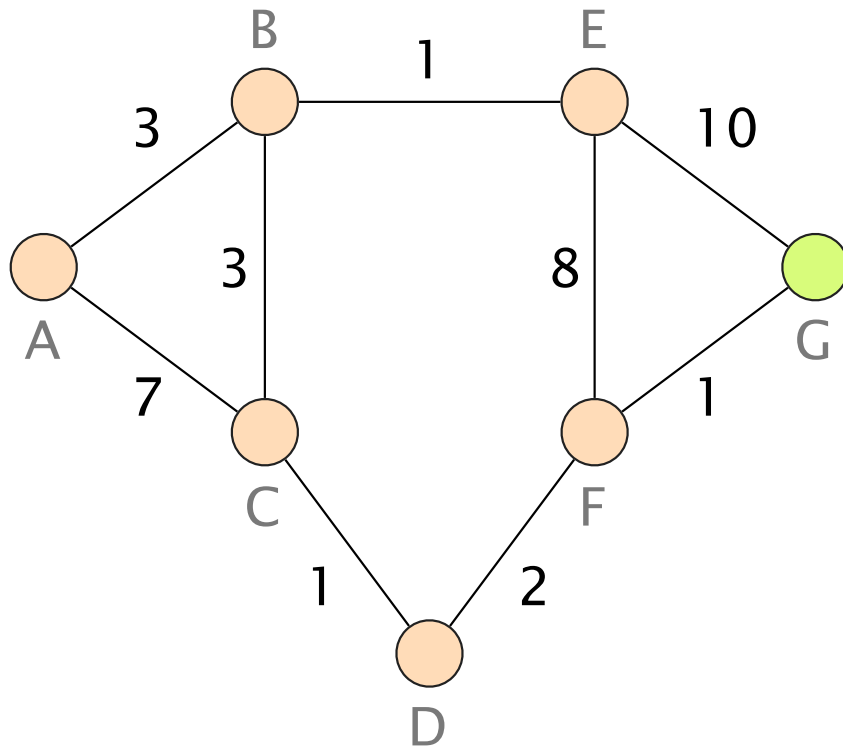


Consider dynamic weights

**Important:** different weights  
for both link directions

Next-hop as tie-break value

## Task 3: Distance Vector



Now we look at a **distance vector** approach (Dijkstra: link-state)

Each node sends a distance vector repeatedly to all neighbors

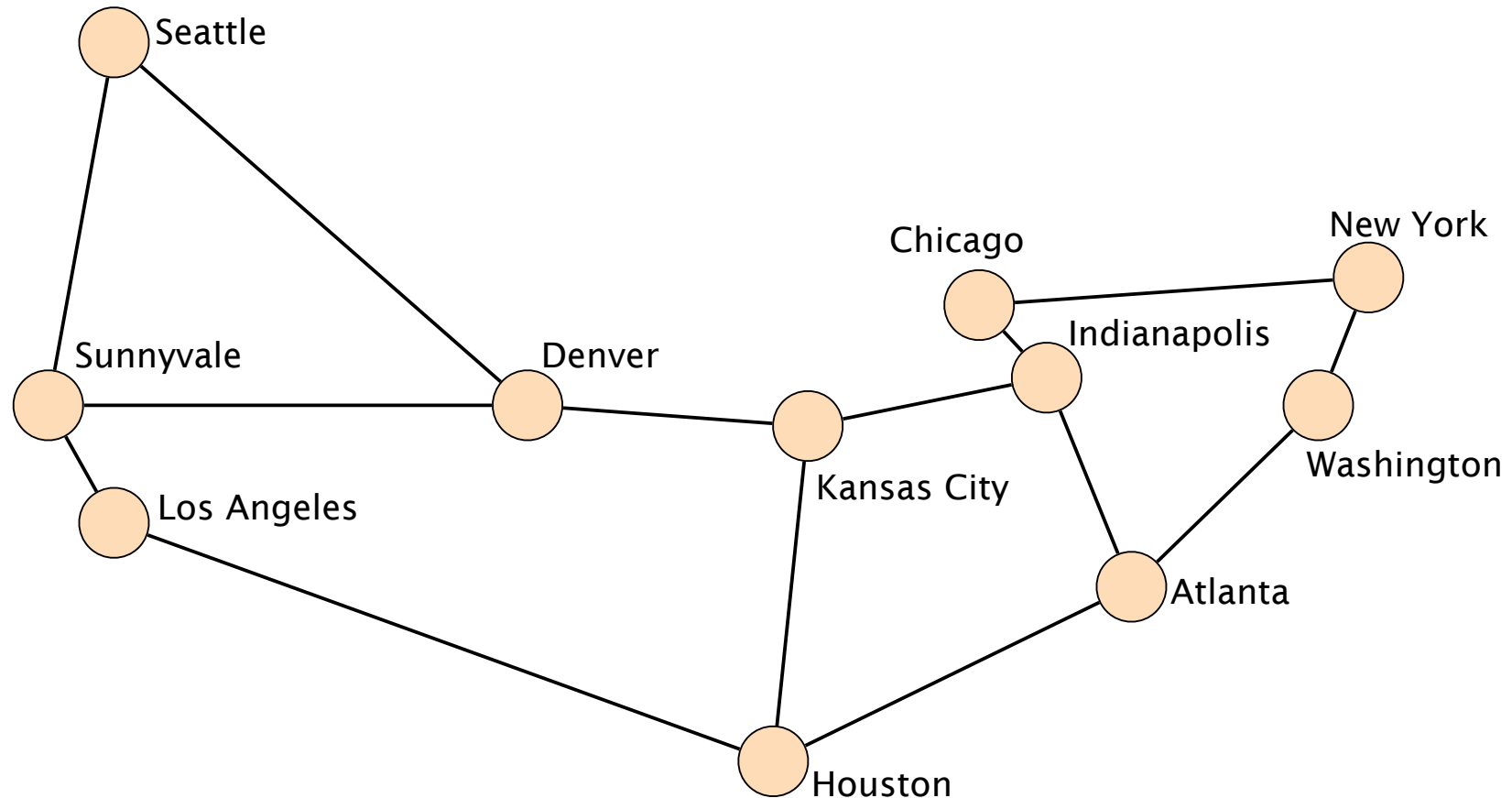
No global network view required

## Task 4: Source-and-Destination-Based Routing

Is it possible to route packets based on the source address?

What are advantages/disadvantages?

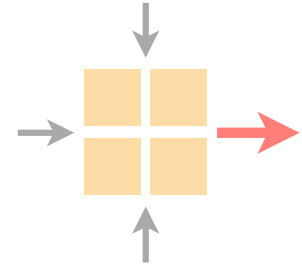
## Task 5: Link Weight Configuration



The Abilene network in the US

# Communication Networks

## Exercise 2



Quick Dijkstra repetition

Overview current assignment

Questions?

Time for you to solve the tasks

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# Group registration is open

Register on <https://comm-net.ethz.ch/registration>

Enter your nethz usernames

Please build groups of **three** students

Join the *#group\_search* Slack channel if you need more members



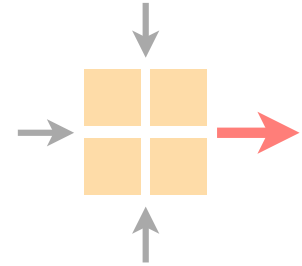
# “Internet Hackathon”

March 26, 6 - 10 pm



# Communication Networks

## Exercise 2



Quick Dijkstra repetition

Overview current assignment

Questions?

**Time for you to solve the tasks**

Solutions will be published next week