Communication Networks

Spring 2020

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Materials inspired from Scott Shenker & Jennifer Rexford

Last week on

Communication Networks

Communication Networks

Part 1: General overview

#1 **What is a network made of?**

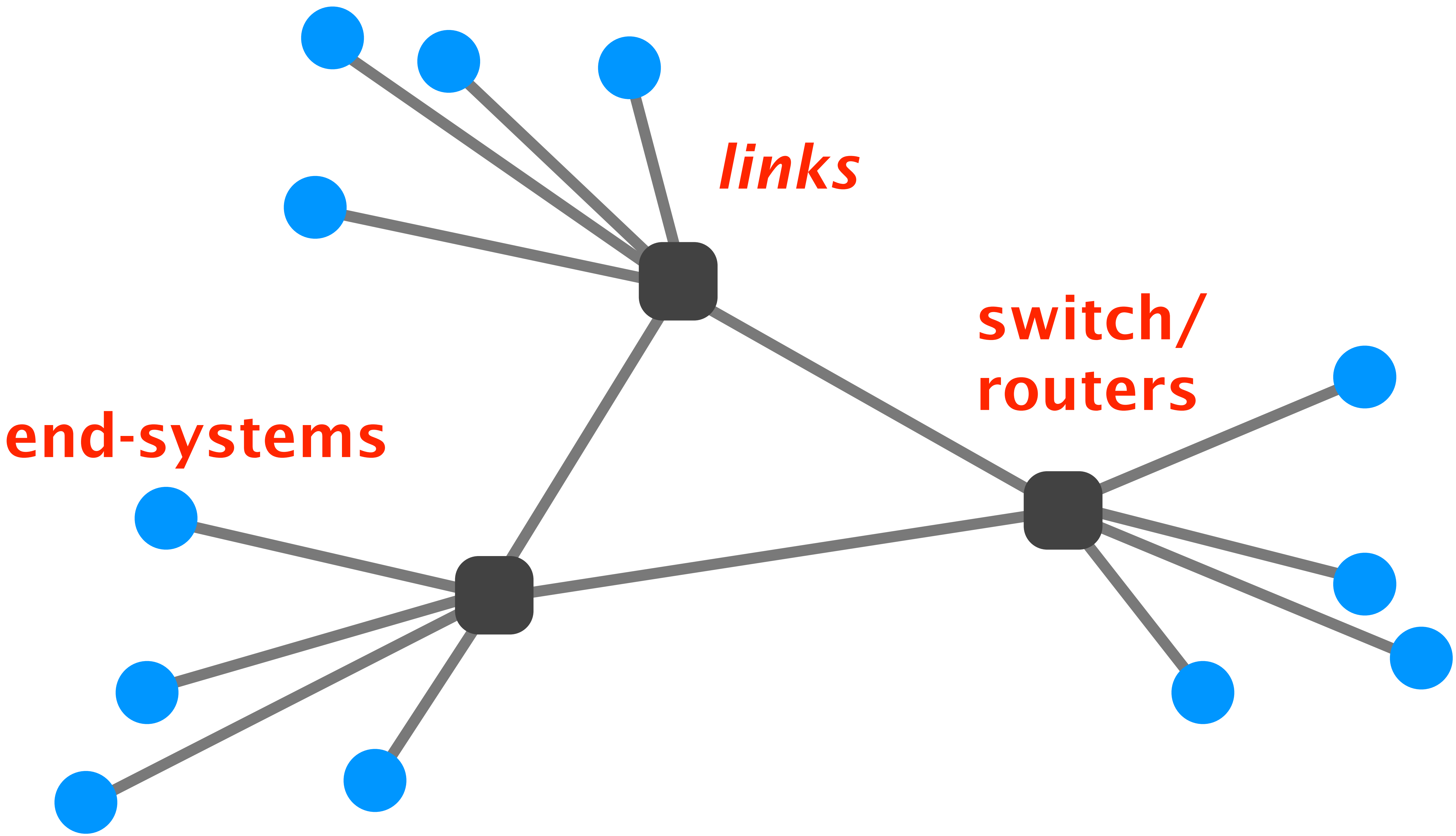
How is it shared?

How is it organized?

How does communication happen?

How do we characterize it?

Networks are composed of three basic components



Communication Networks

Part 1: General overview

What is a network made of?

#2 **How is it shared?**

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There exist two approaches to sharing:

reservation and **on-demand**

	Reservation	On-demand
principle	reserve the bandwidth you need in advance	send data when you need

In practice, the approaches are implemented using **circuit-switching** or **packet-switching**

	Reservation	On-demand
implem.	circuit-switching	packet-switching

Pros and cons of circuit switching

advantages

predictable performance

simple & fast switching

once circuit established

disadvantages

inefficient if traffic is bursty or short

complex circuit setup/teardown

which adds delays to transfer

requires new circuit upon failure

Pros and cons of packet switching

advantages	disadvantages
efficient use of resources	unpredictable performance
simpler to implement than circuit switching	requires buffer management and congestion control
route around trouble	

Communication Networks

Part 1: General overview

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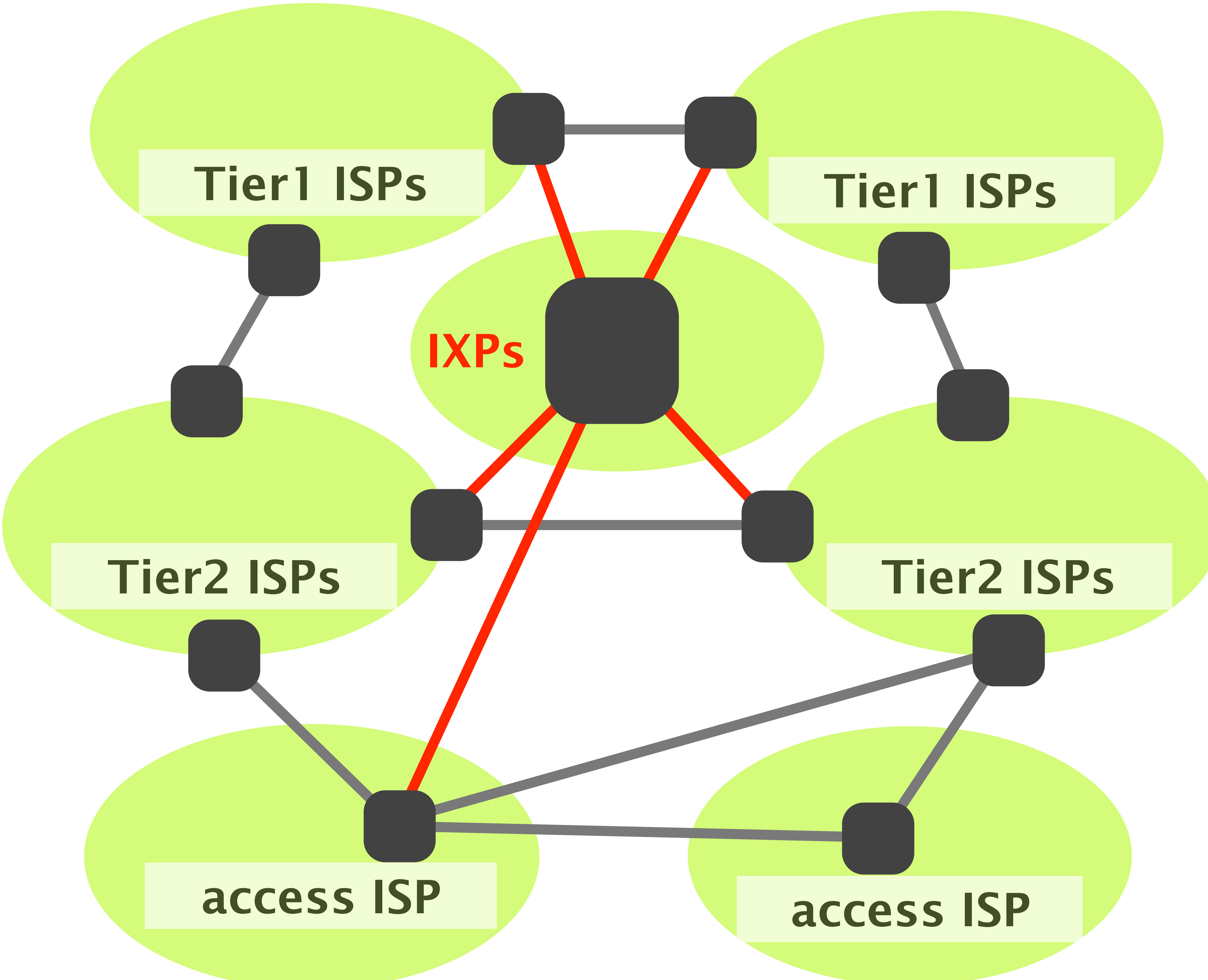
#3 How is it organized?

How does communication happen?

How do we characterize it?



Tier1 ISPs
Tier1 ISPs
IXPs
Tier2 ISPs
Tier2 ISPs
access ISP
access ISP

This week on

Communication Networks

Communication Networks

Part 1: General overview

What is a network made of?

How is it shared?

How is it organized?

#4 How does communication happen?

#5 How do we characterize it?

Communication Networks

Part 1: General overview

What is a network made of?

How is it shared?

How is it organized?

#4 How does communication happen?

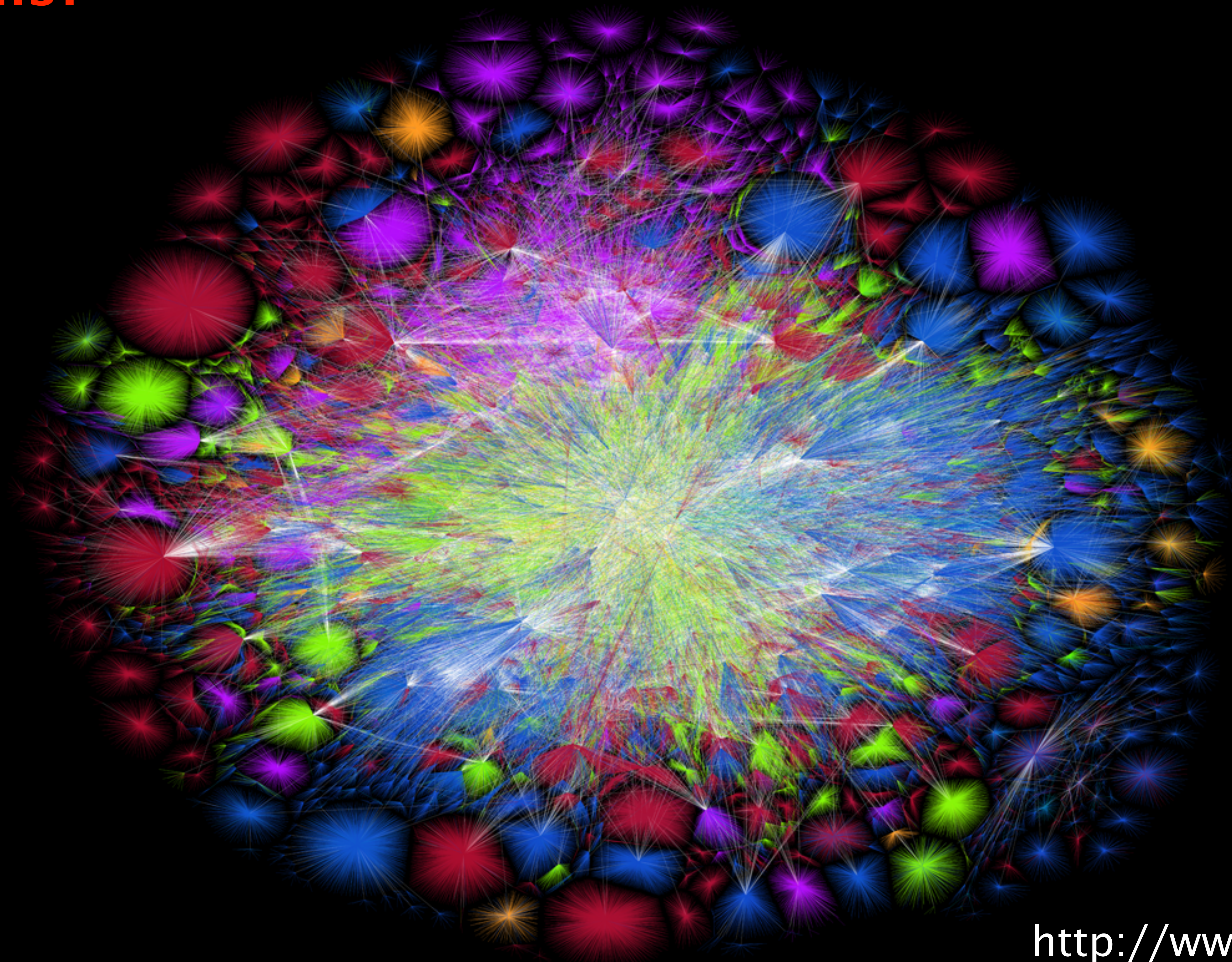
How do we characterize it?

The Internet should allow

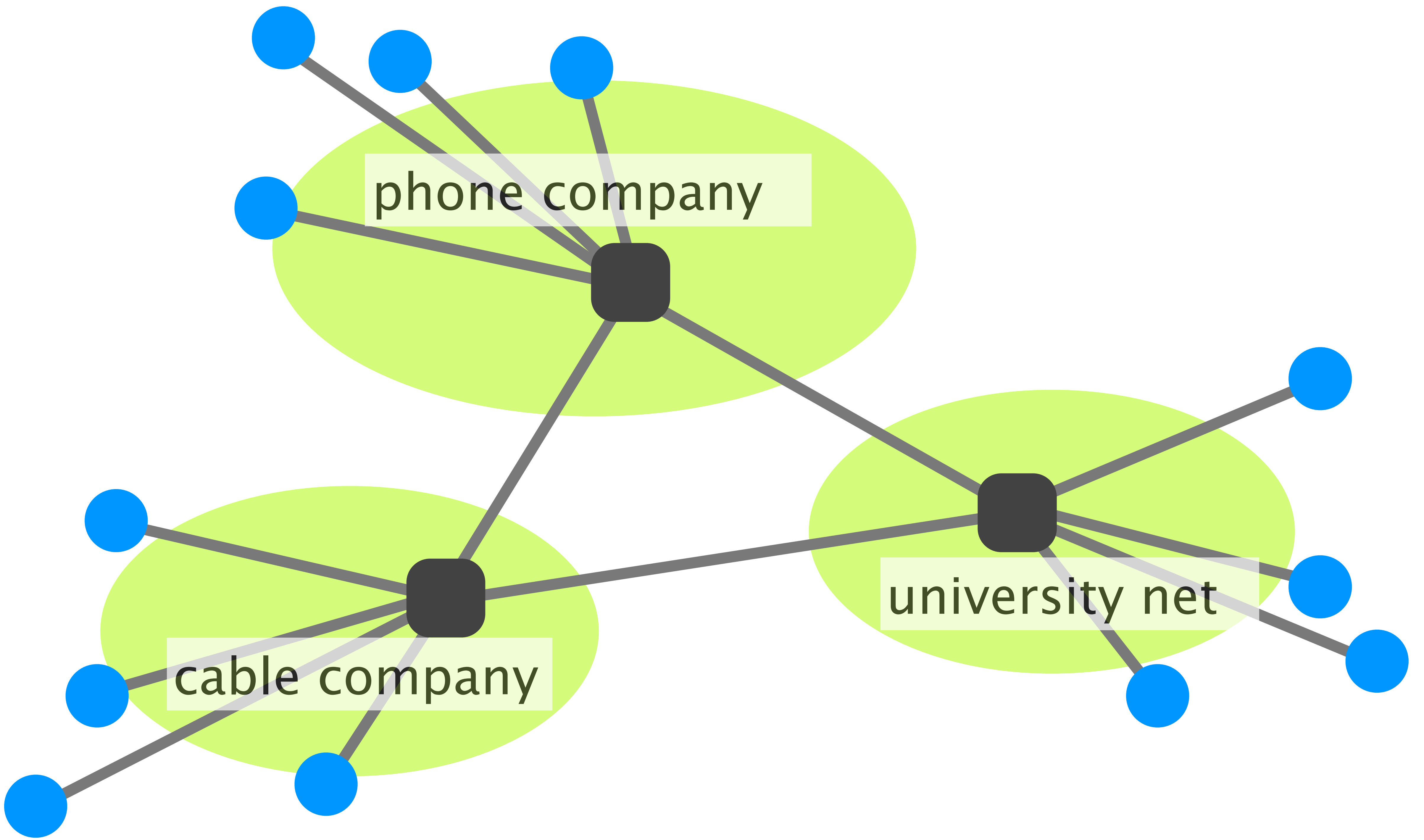
processes on different hosts
to exchange data

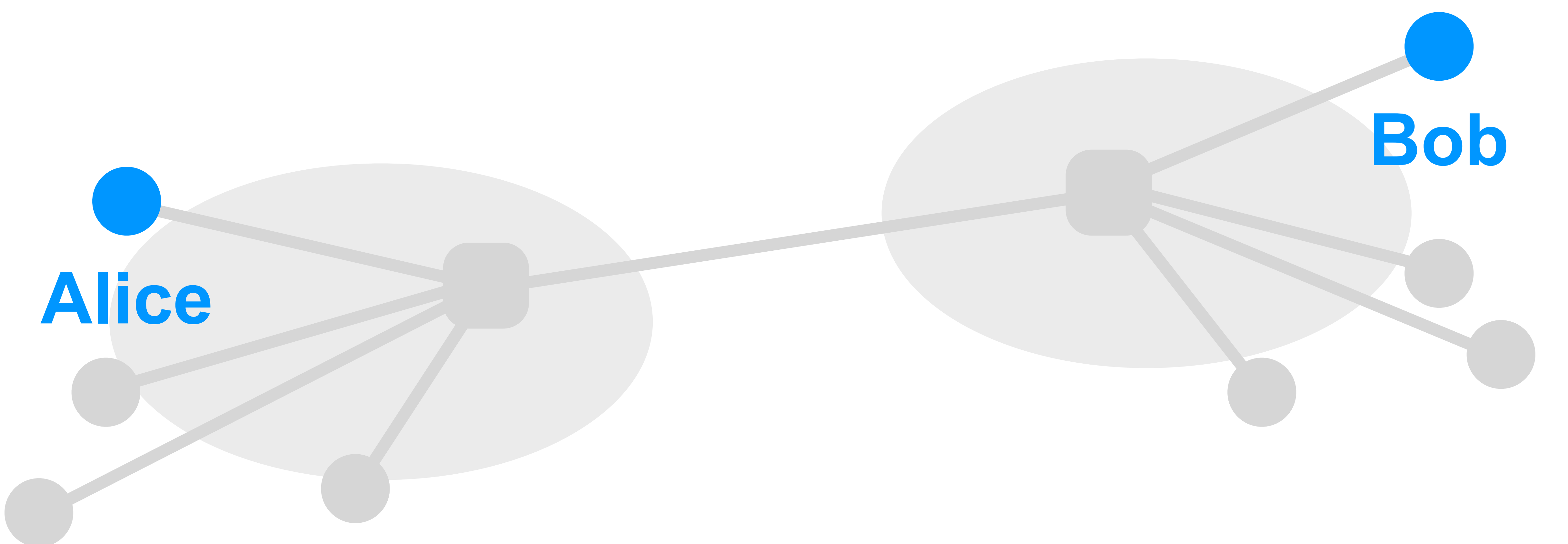
everything else is just commentary…

How do you exchange data in a network as complex as **this?**



http://www.opte.org


phone company
university net
cable company

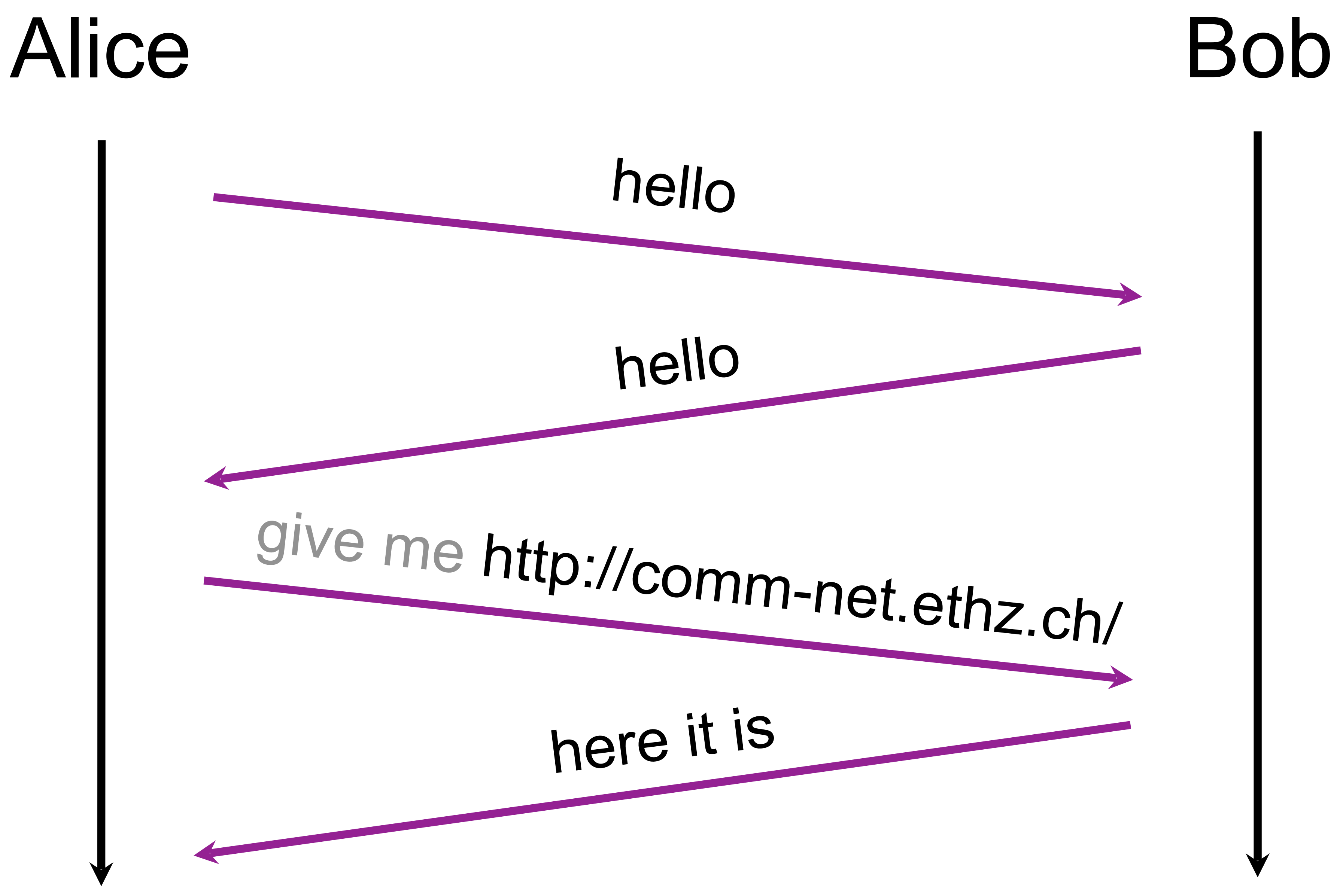


Bob
Alice

To exchange data, Alice and Bob use a set of network protocols

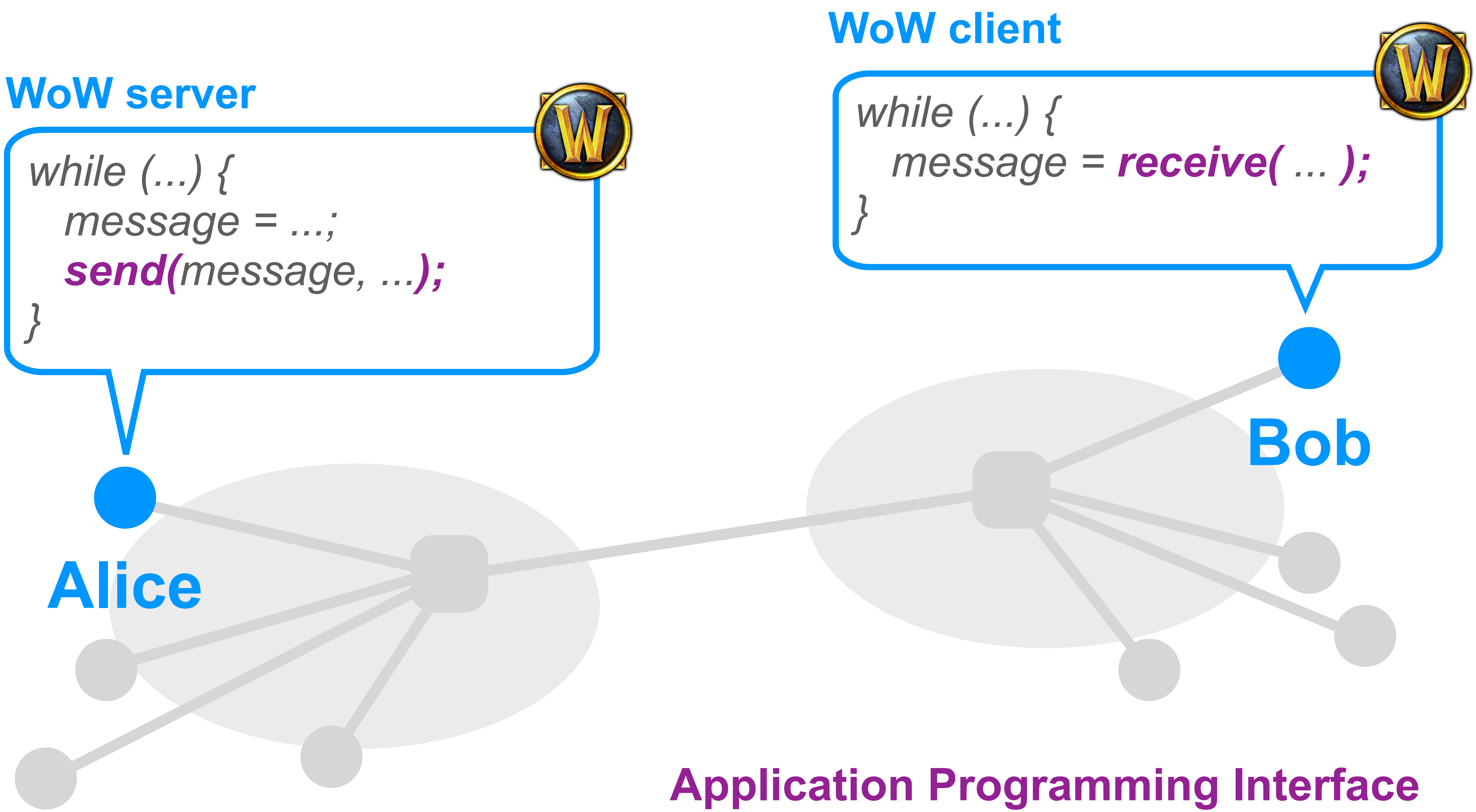
A protocol is like a conversational convention:

who should talk next and how they should respond



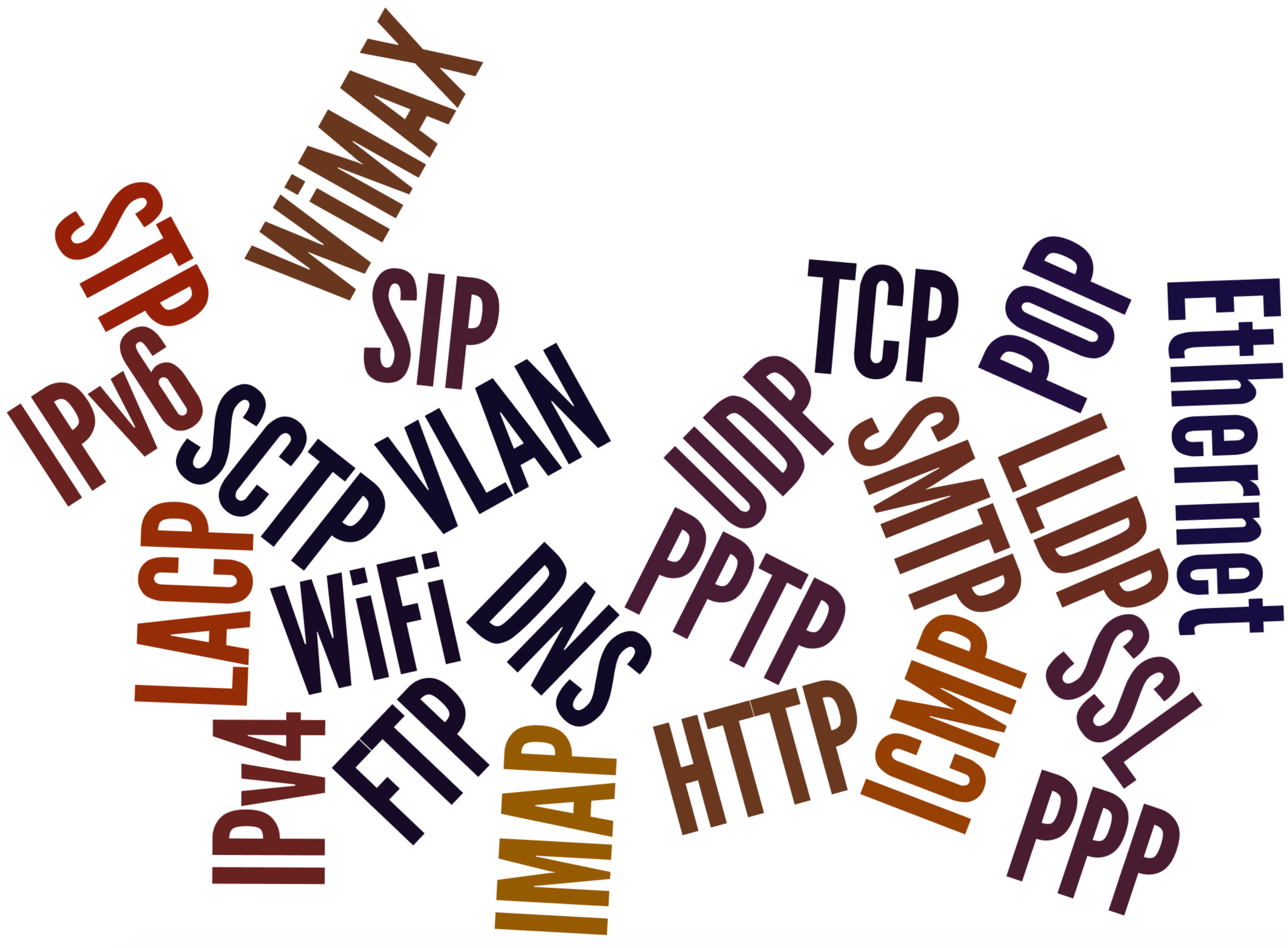
Sometimes implementations are not compliant...

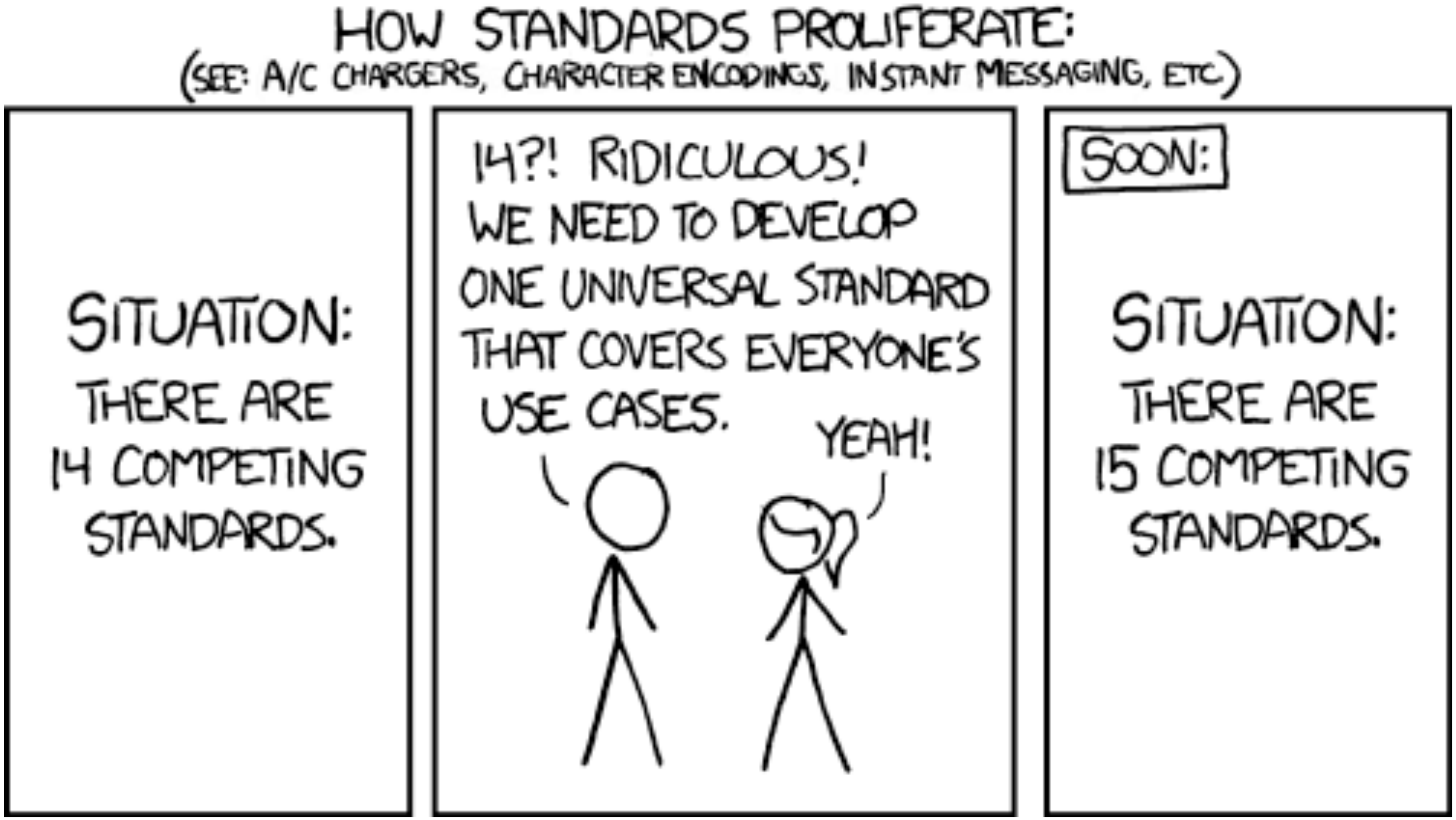
Each protocol is governed by a specific interface



In practice, there exists **a lot** of network protocols.

How does the Internet organize **this?**





https://xkcd.com/927/

Modularity is a key component of any good system

Problem

can’t build large systems out of spaghetti code

hard (if not, impossible) to understand, debug, update

need to bound the scope of changes

evolve the system without rewriting it from scratch

Solution

Modularity is how we do it

…and understand the system at a higher-level

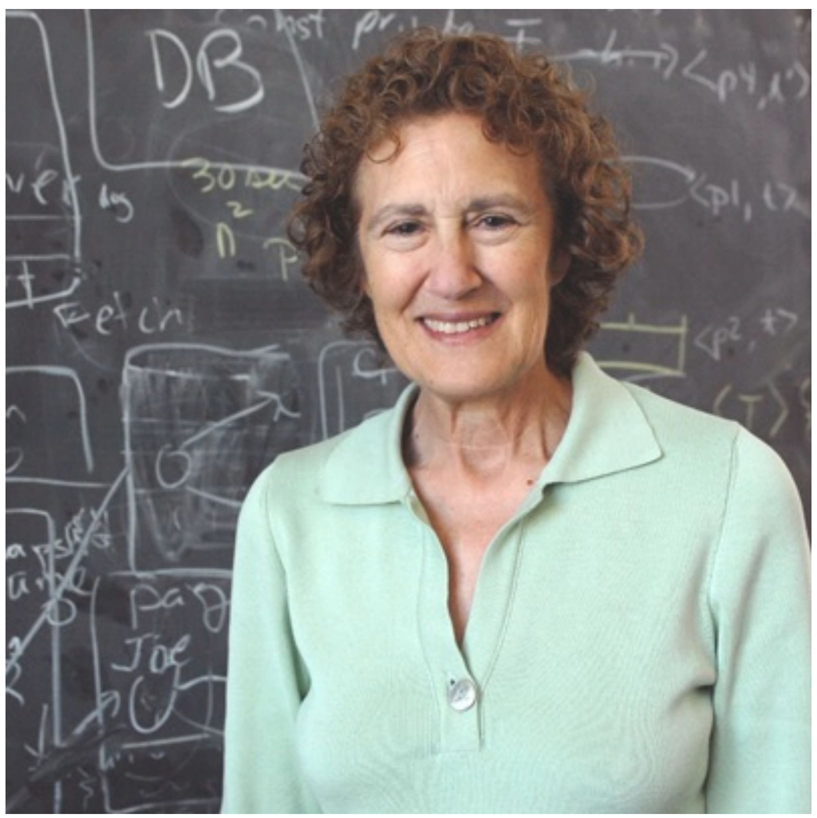


Photo: Donna Coveney

Modularity,
based on abstraction,
is *the* way things get done

— *Barbara Liskov*, MIT

To provide structure to the design of network protocols, network designers organize protocols in layers

and the network hardware/software that implement them

Internet communication can be decomposed in 5 independent layers (or 7 layers for the OSI model)

	layer
L5	Application
L4	Transport
L3	Network
L2	Link
L1	Physical

Each layer provides a service to the layer above

	layer	service provided:
L5	Application	network access
L4	Transport	end-to-end delivery (reliable or not)
L3	Network	global best-effort delivery
L2	Link	local best-effort delivery
L1	Physical	physical transfer of bits

Each layer provides a service to the layer above

by using the services of the layer directly below it

Applications

…built on…

Reliable (or unreliable) transport

…built on…

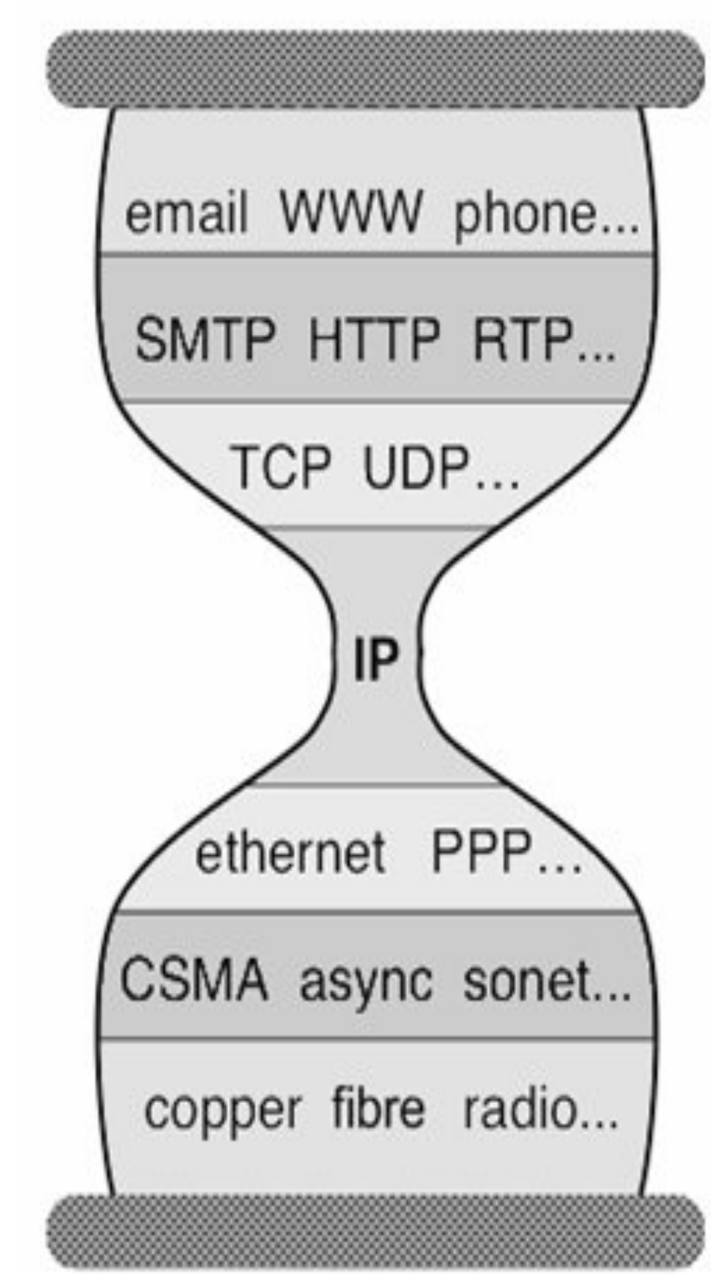
Best-effort global packet delivery

…built on…

Best-effort local packet delivery

…built on…

Physical transfer of bits



Each layer has a unit of data

	layer	role
L5	Application	exchanges messages between processes
L4	Transport	transports segments between end-systems
L3	Network	moves packets around the network
L2	Link	moves frames across a link
L1	Physical	moves bits across a physical medium

Each layer (except for L3) is implemented with different protocols

	layer	protocol
L5	Application	HTTP, SMTP, FTP, SIP, …
L4	Transport	TCP, UDP, SCTP
L3	Network	IP
L2	Link	Ethernet, Wifi, (A/V)DSL, WiMAX, LTE, …
L1	Physical	Twisted pair, fiber, coaxial cable, …

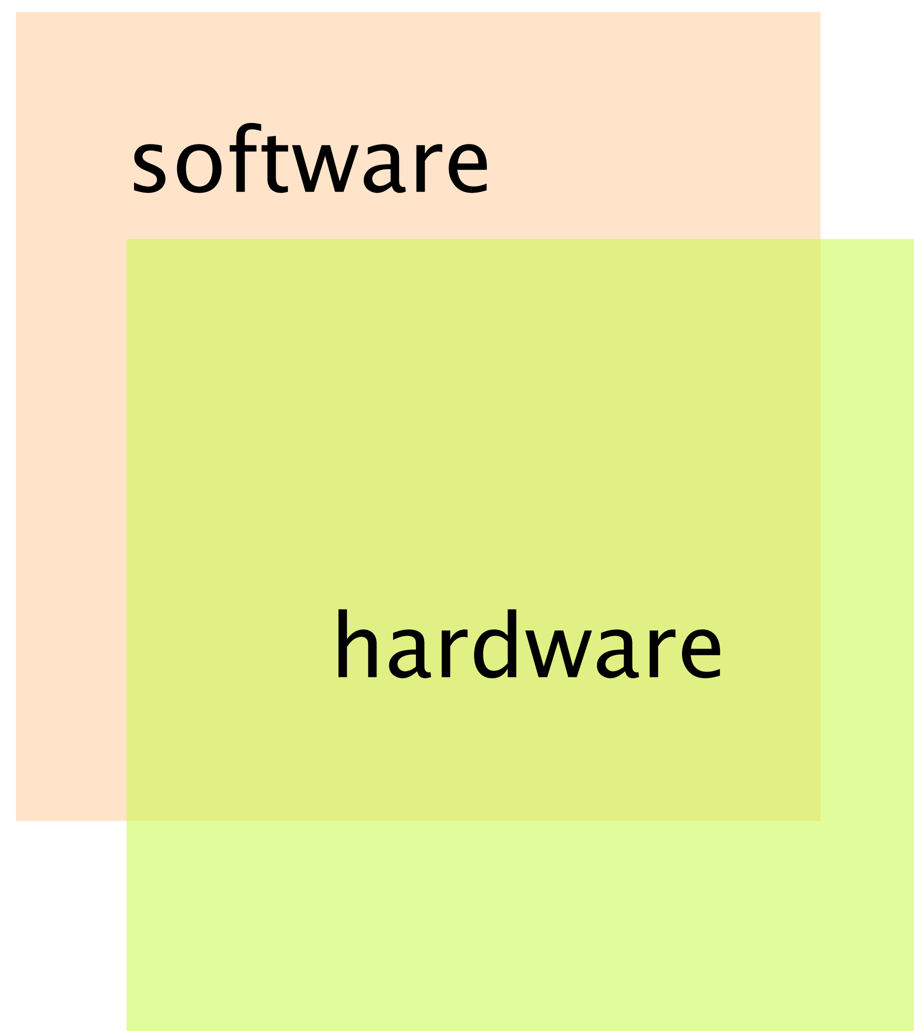
The Internet Protocol (IP) acts as an unifying, network, layer

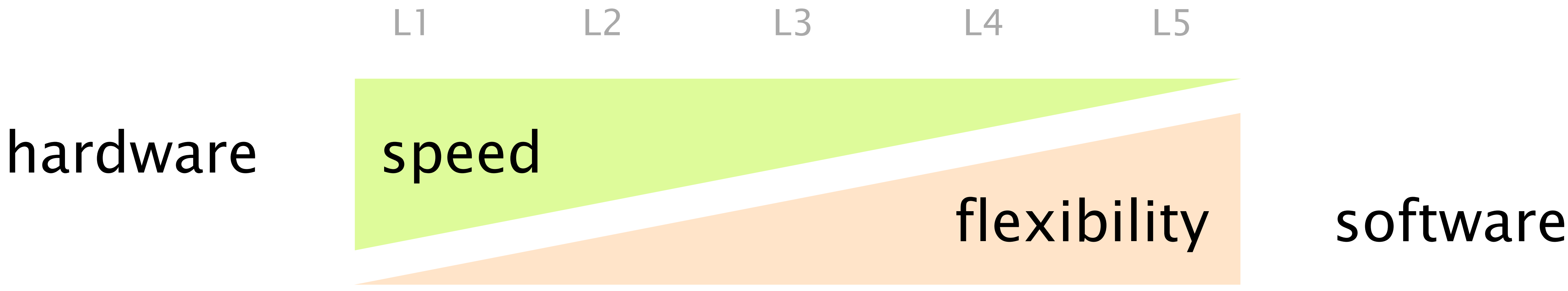
	layer	protocol
L5	Application	HTTP, SMTP, FTP, SIP, …
L4	Transport	TCP, UDP, SCTP
L3	Network	IP
L2	Link	Ethernet, Wifi, (A/V)DSL, Cable, LTE, …
L1	Physical	Twisted pair, fiber, coaxial cable, …

Each layer is implemented with different protocols and technologies

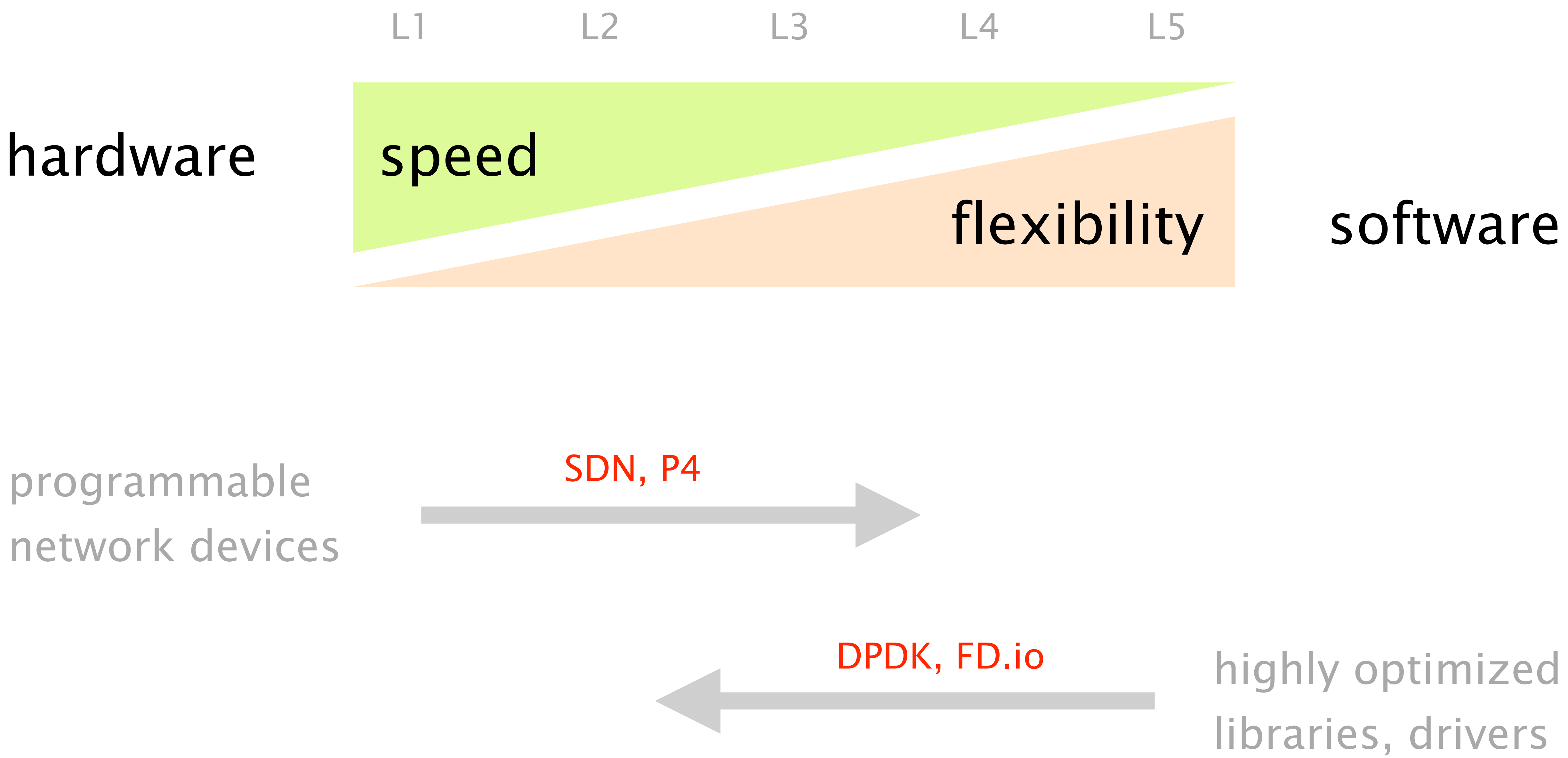
	layer
L5	Application
L4	Transport
L3	Network
L2	Link
L1	Physical

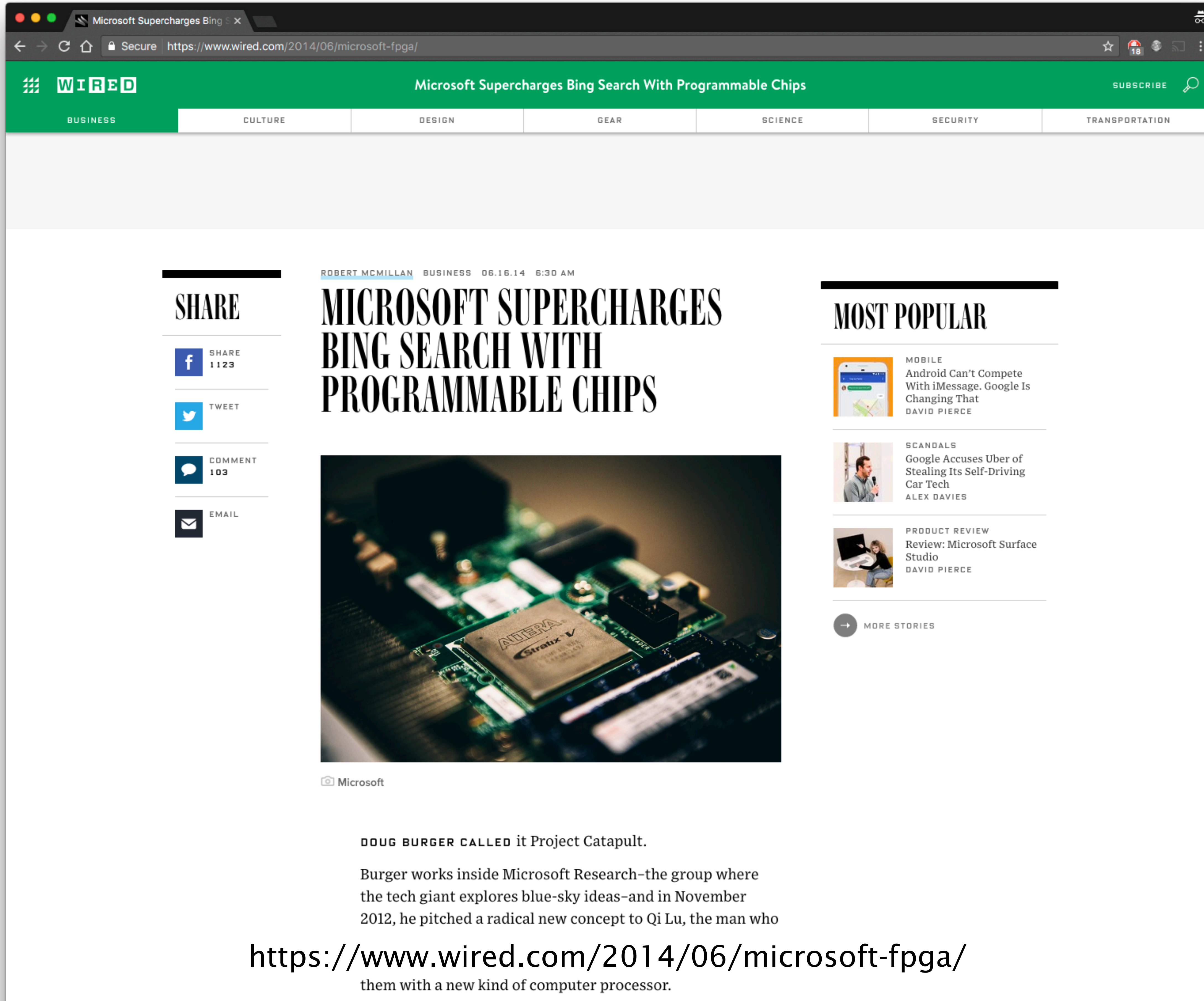
technology




L1
L2
L3
L4
L5
hardware
speed
flexibility
software

Software and hardware advancements





Microsoft Supercharges Bing Search With Programmable Chips

Secure https://www.wired.com/2014/06/microsoft-fpga/

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ROBERT MCMILLAN BUSINESS 06.16.14 6:30 AM

MICROSOFT SUPERCHARGES BING SEARCH WITH PROGRAMMABLE CHIPS

SHARE

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Microsoft

DOUG BURGER CALLED it Project Catapult.

Burger works inside Microsoft Research–the group where the tech giant explores blue-sky ideas–and in November 2012, he pitched a radical new concept to Qi Lu, the man who

them with a new kind of computer processor.

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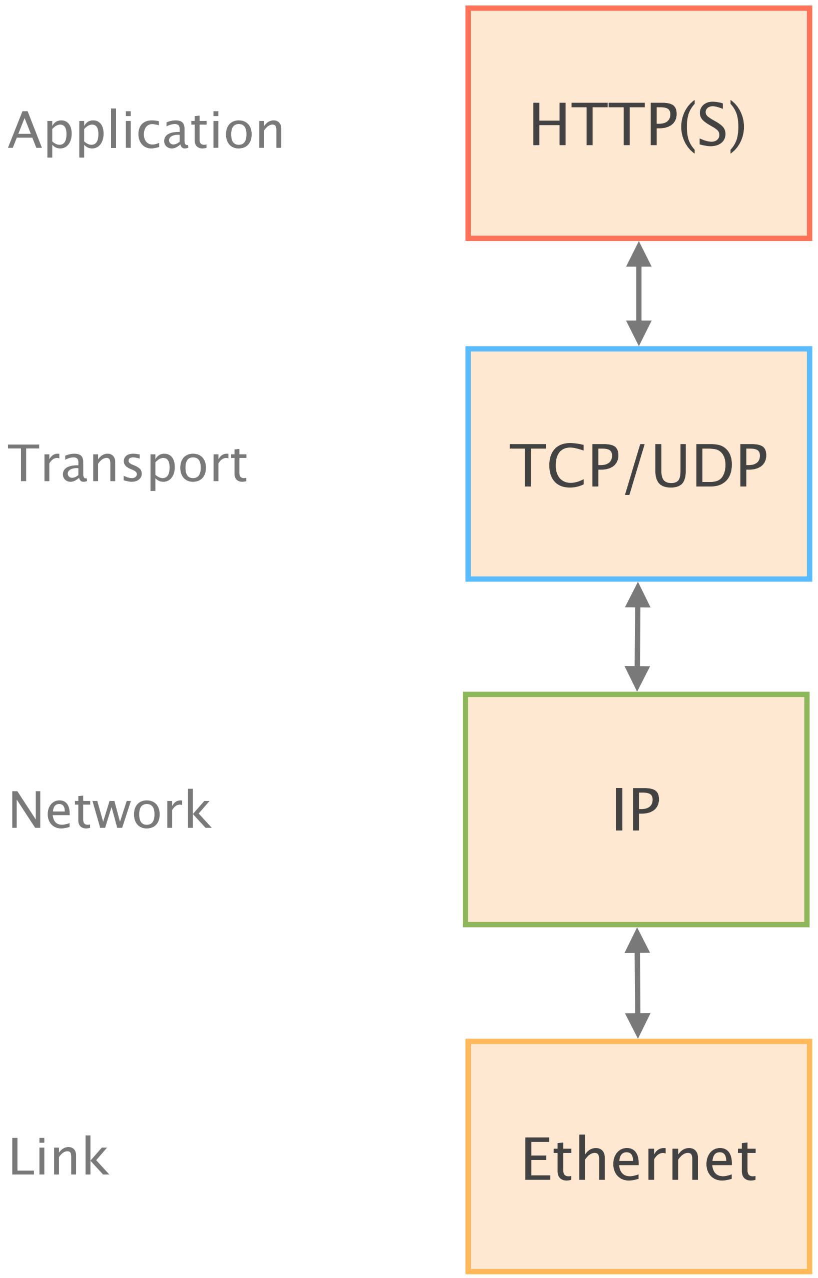
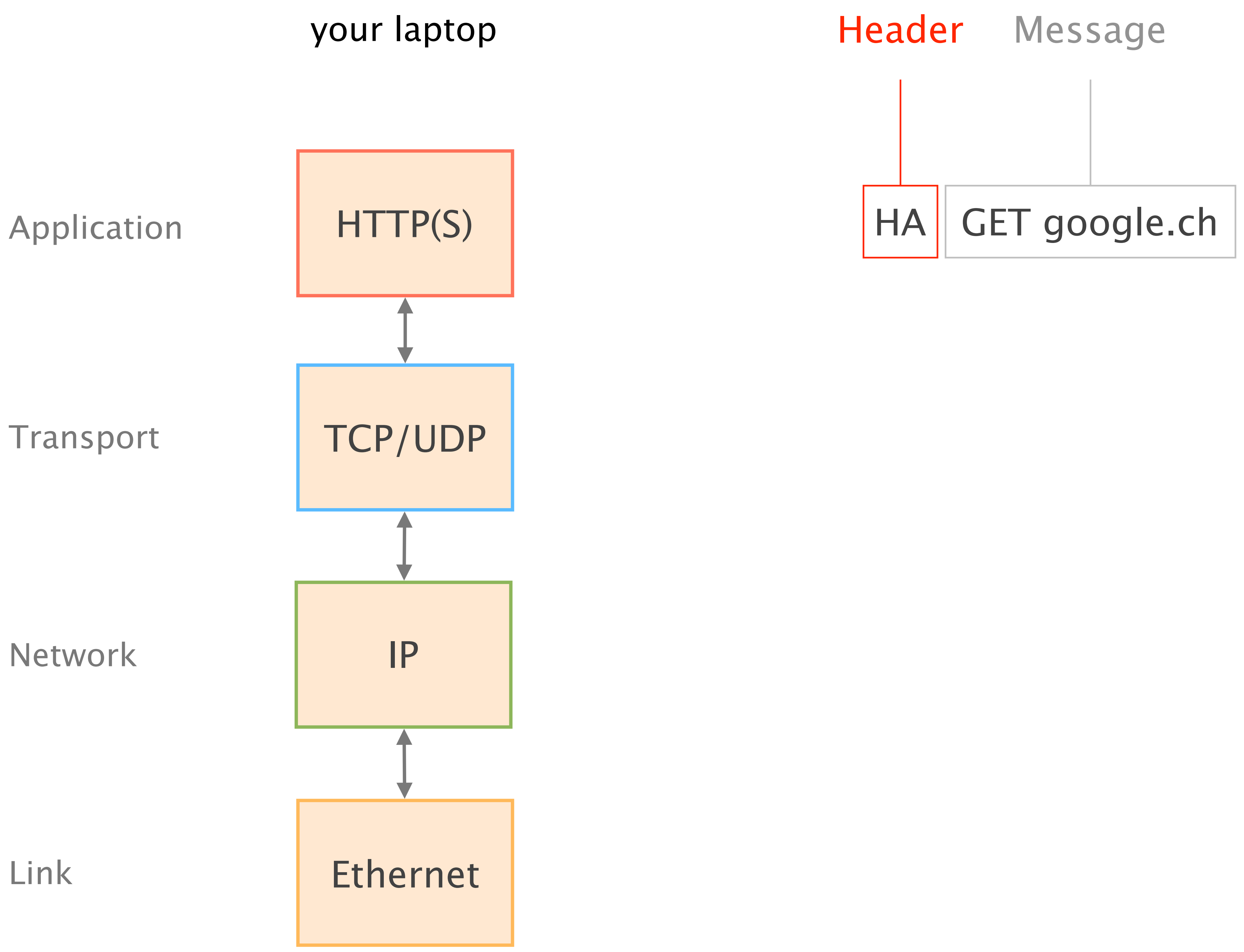
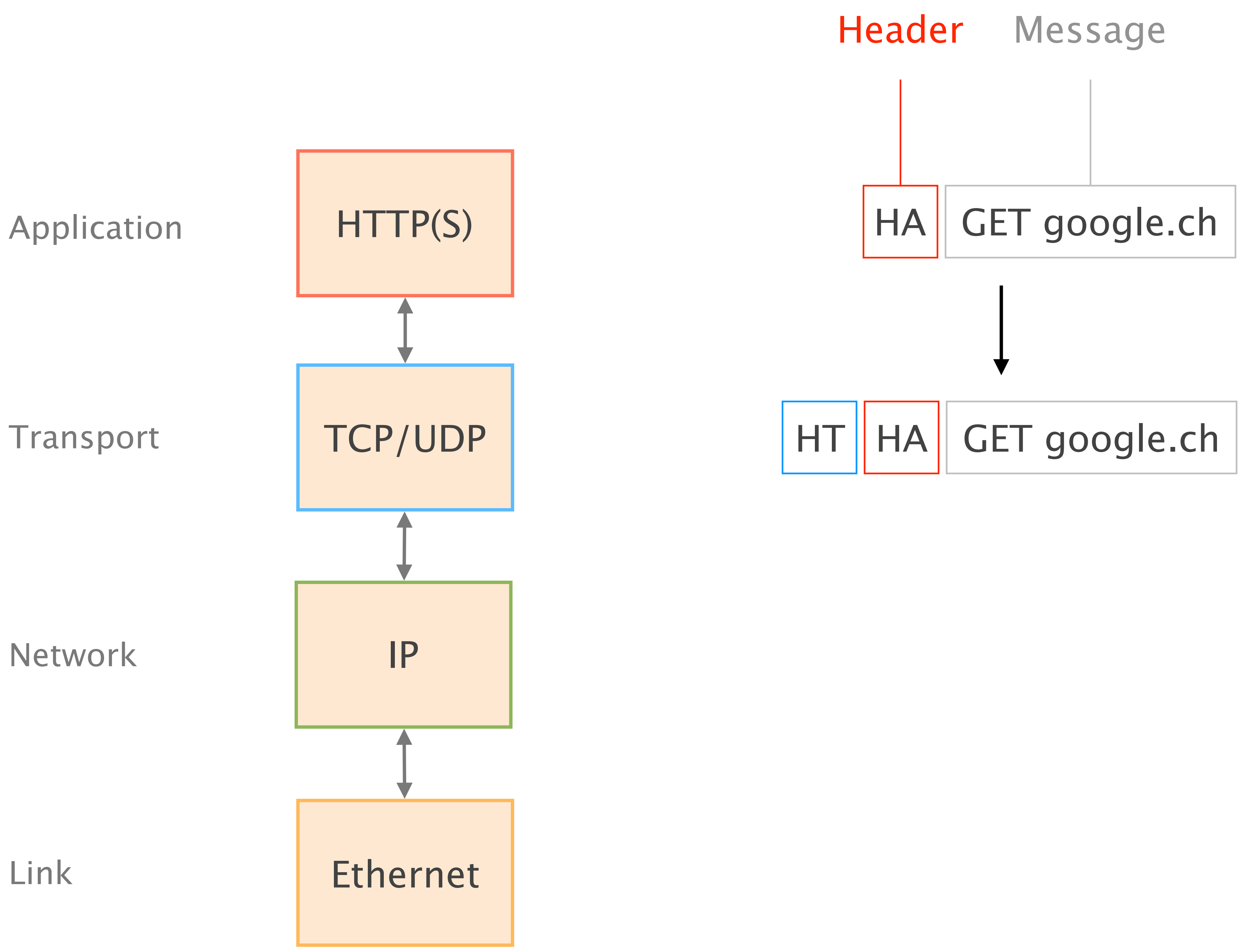
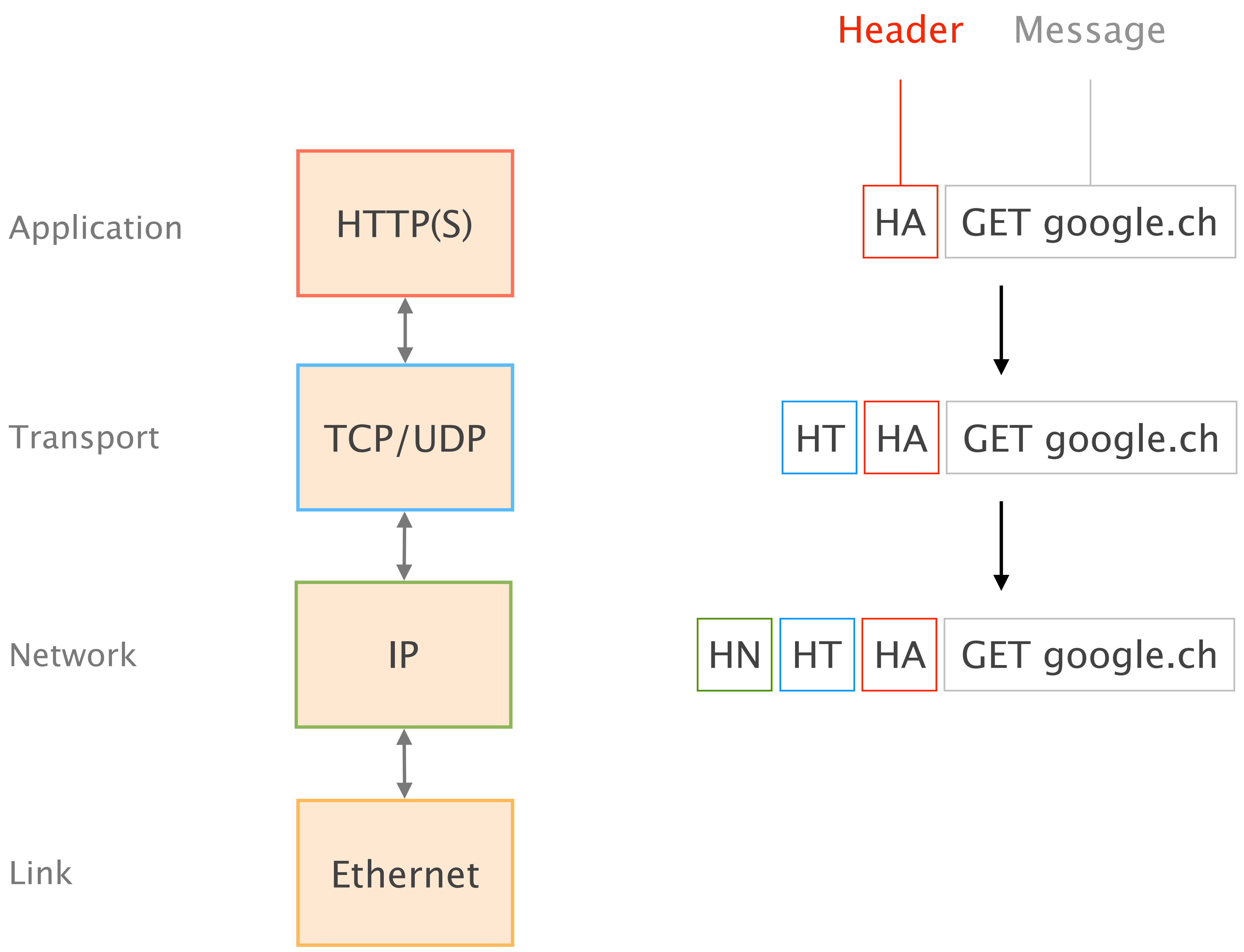
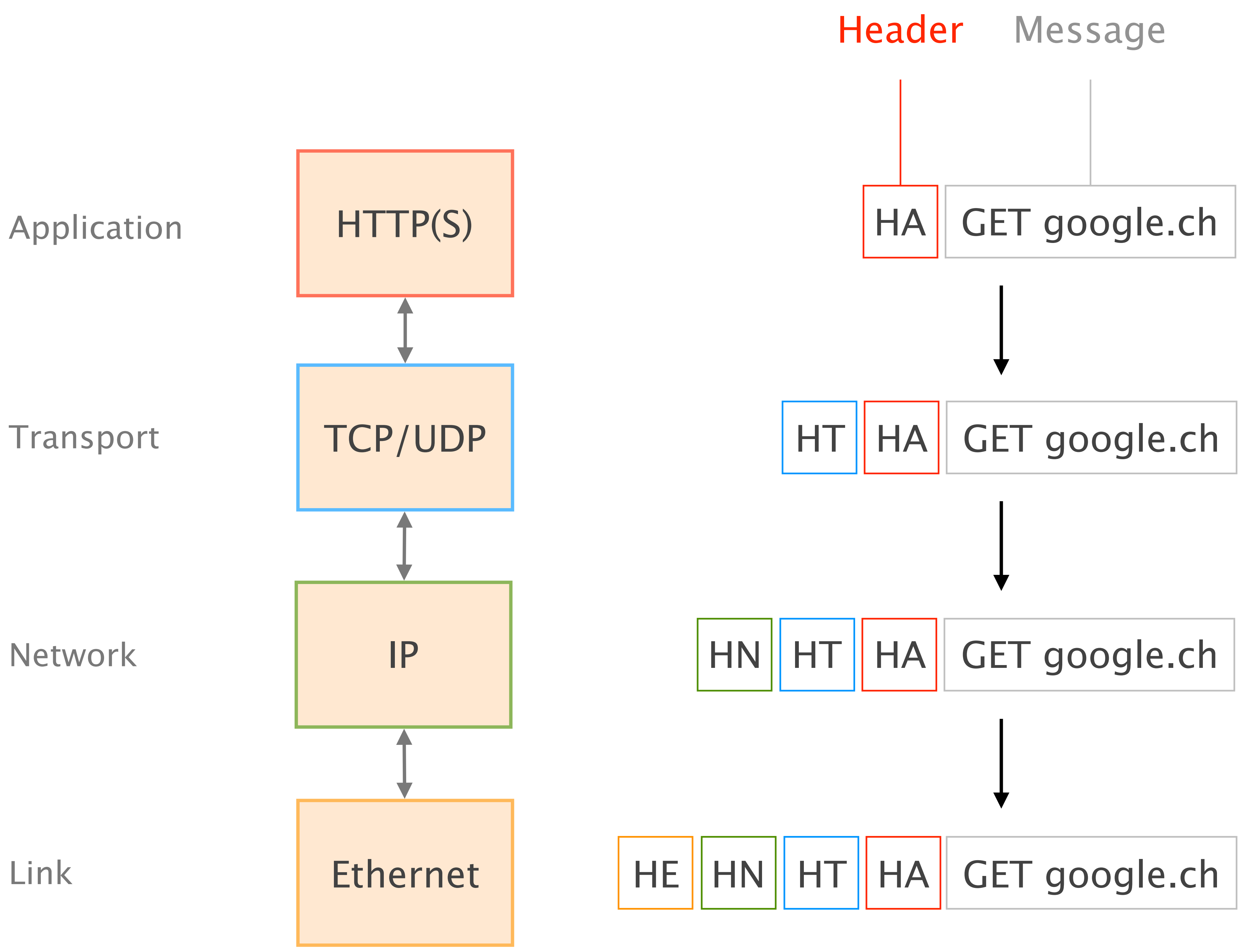
Each layer takes messages from the layer above, and *encapsulates* with its own header and/or trailer



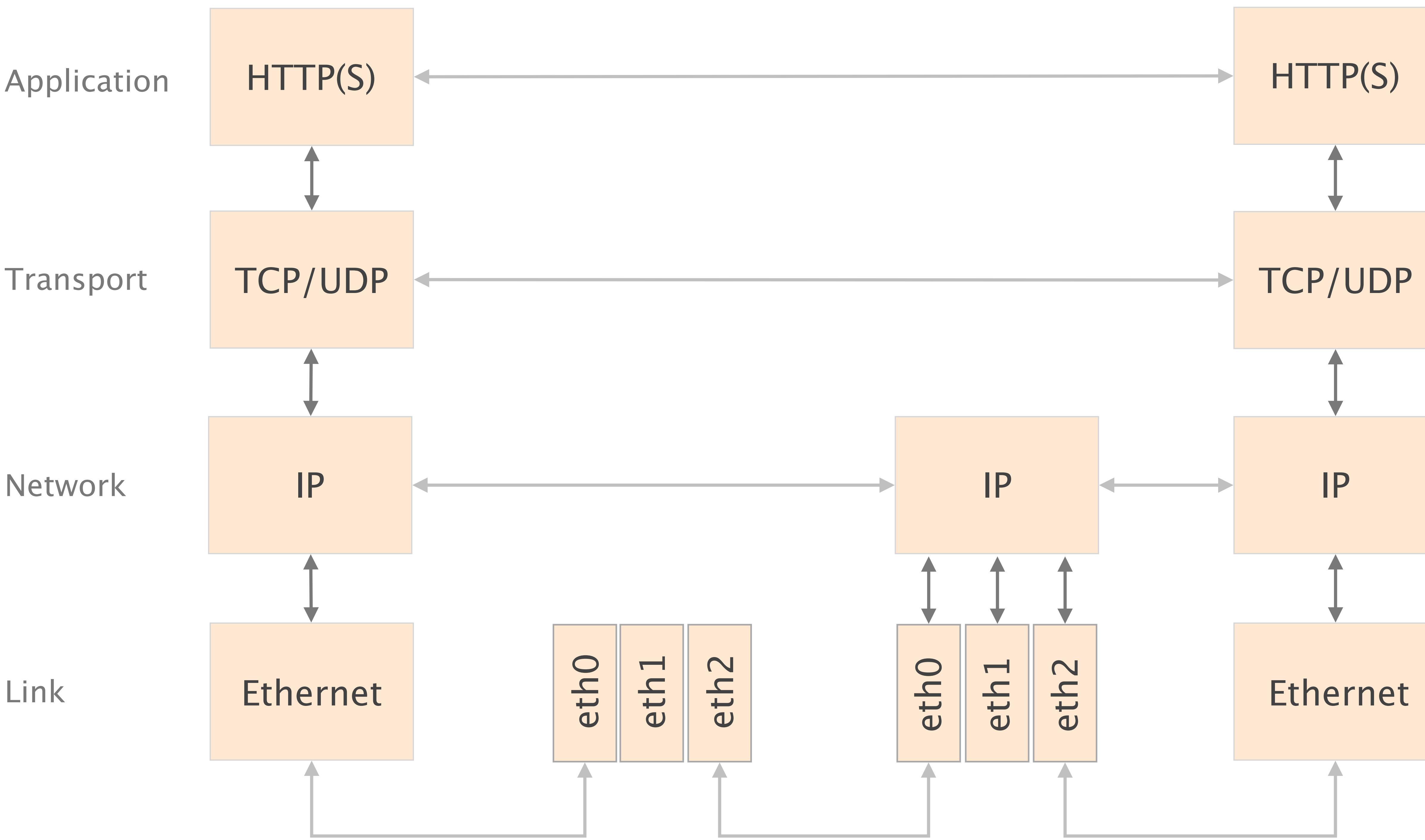

your laptop
Header
Message
Application
HTTP(S)
HA
GET google.ch
Transport
TCP/UDP
Network
IP
Link
Ethernet


Header
Message
Application
HTTP(S)
HA
GET google.ch
Transport
TCP/UDP
HT
HA
GET google.ch
Network
IP
Link
Ethernet

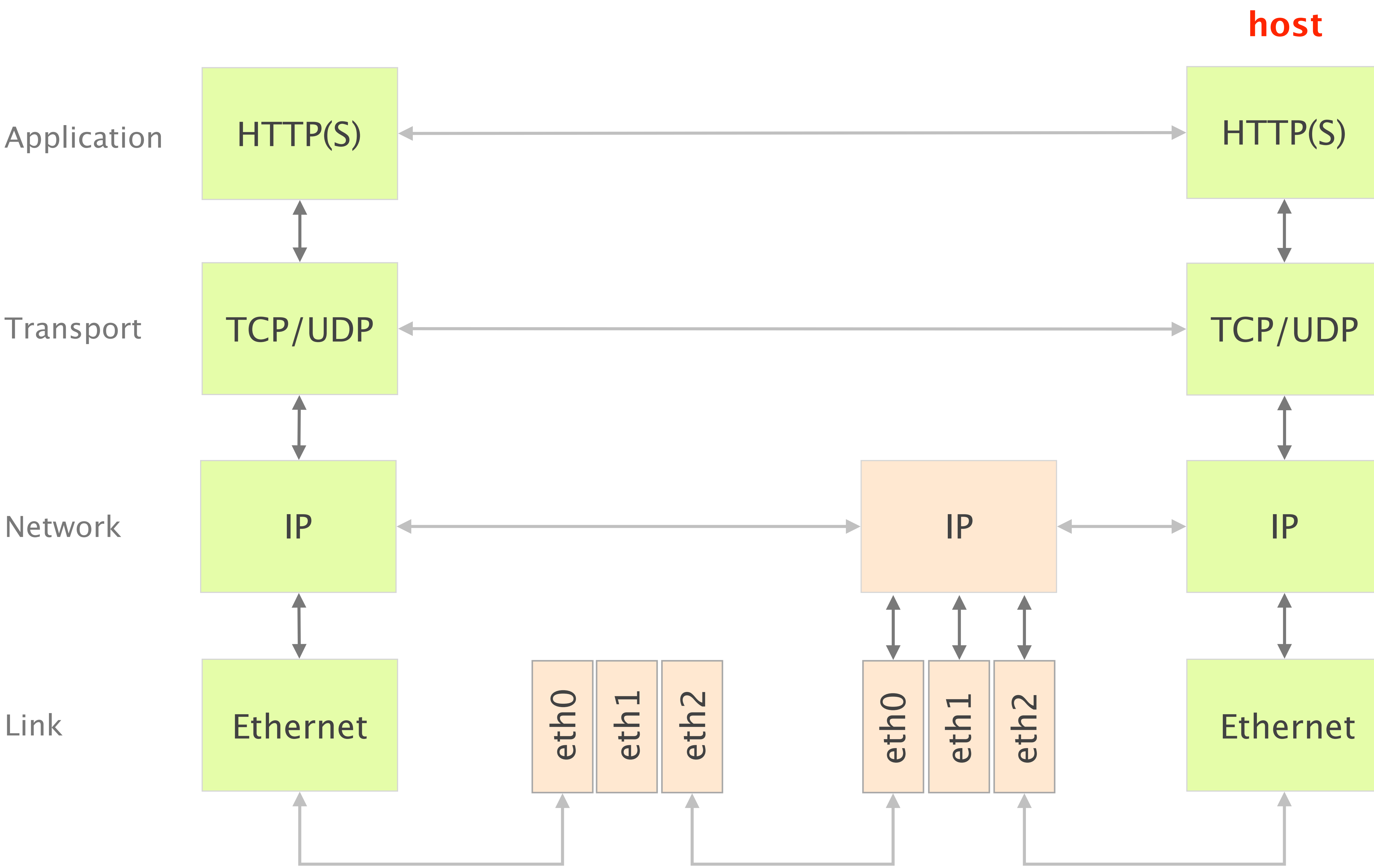

Header
Message
Application
HTTP(S)
HA
GET google.ch
Transport
TCP/UDP
HT
HA
GET google.ch
Network
IP
HN
HT
HA
GET google.ch
Link
Ethernet


Header
Message
Application
HTTP(S)
HA
GET google.ch
Transport
TCP/UDP
HT
HA
GET google.ch
Network
IP
HN
HT
HA
GET google.ch
Link
Ethernet
HE
HN
HT
HA
GET google.ch

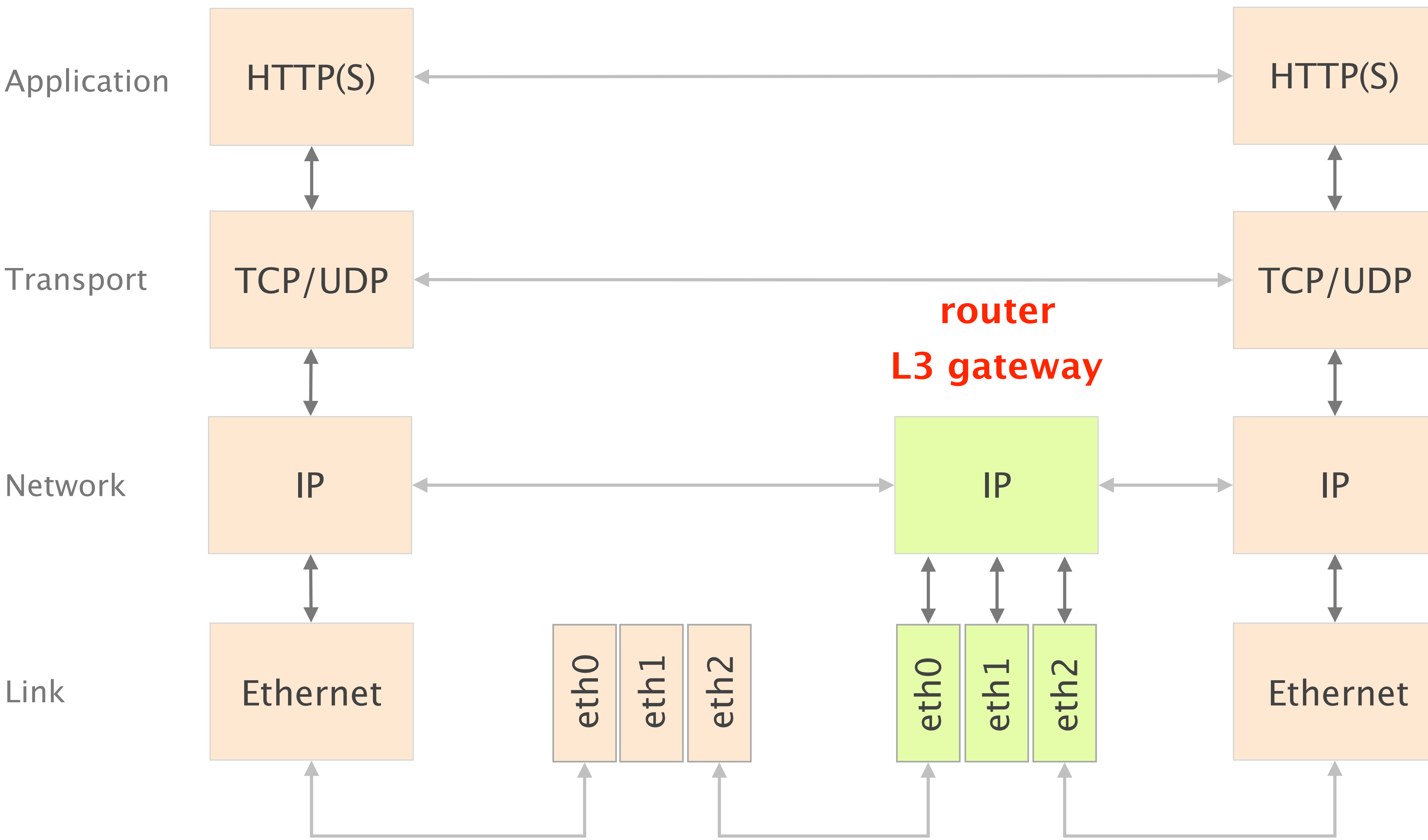
In practice, layers are distributed on every network device



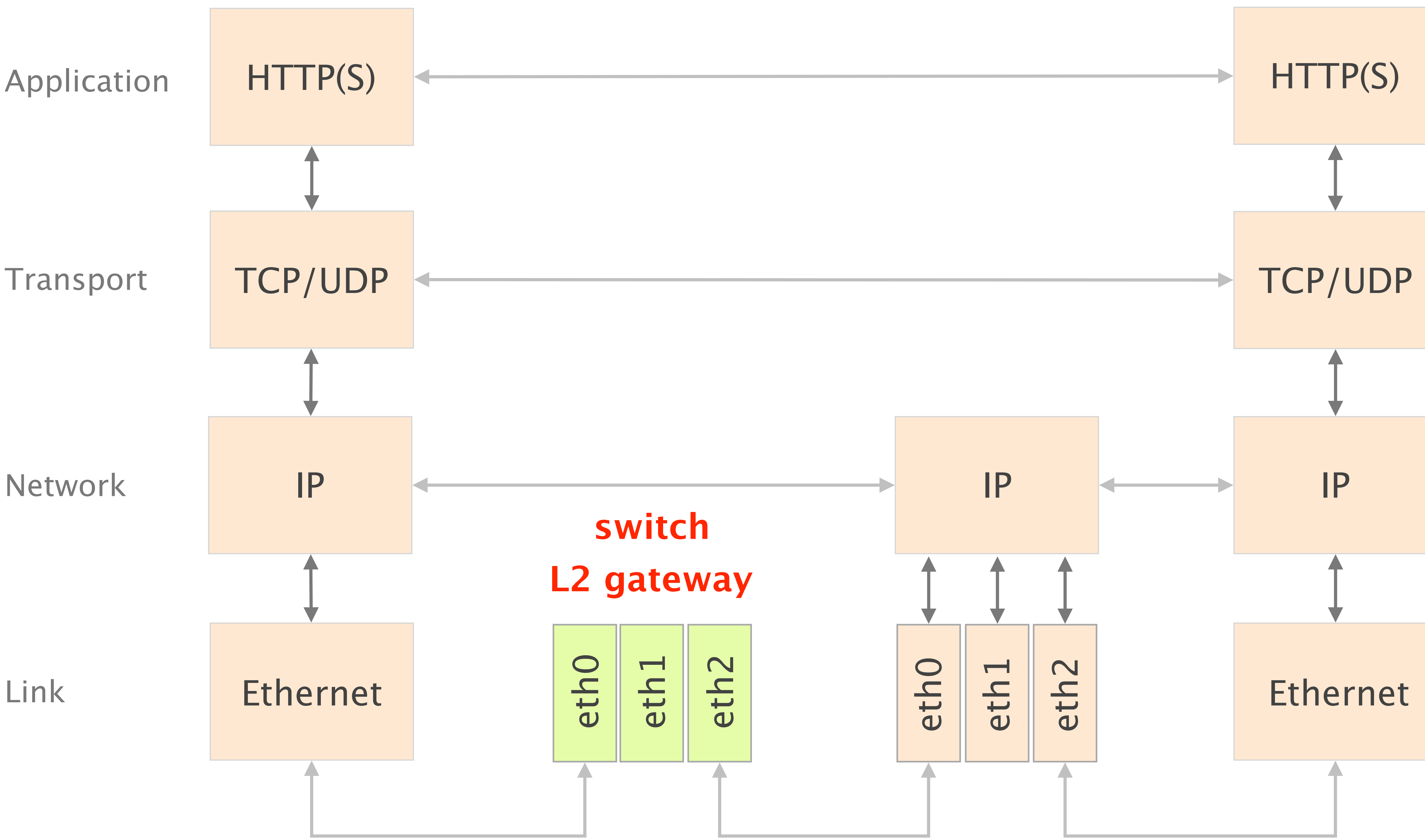
Since when bits arrive they must make it to the application, all the layers exist on a host



Routers act as L3 gateway
as such they implement L2 and L3



Switches act as L2 gateway
as such they only implement L2



Let's see how it looks like in practice

on a host, using Wireshark https://www.wireshark.org

Communication Networks

Part 1: General overview

What is a network made of?

How is it shared?

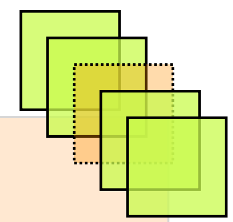
How is it organized?

How does communication happen?

#5 How do we characterize it?

A network *connection* is characterized by its delay, loss rate and throughput

delay



loss

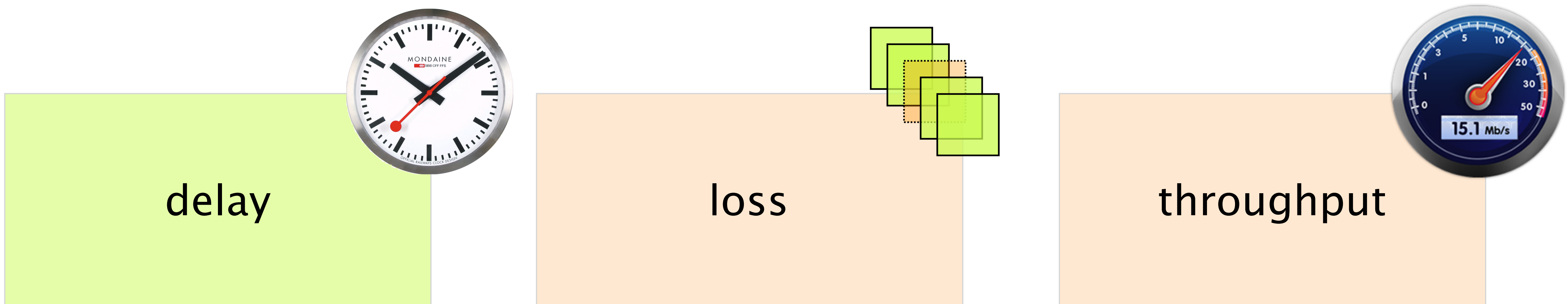
throughput

How long does it take for a packet to reach the destination

What fraction of packets sent to a destination are dropped?

At what rate is the destination receiving data from the source?

A network *connection* is characterized by its delay, loss rate and throughput

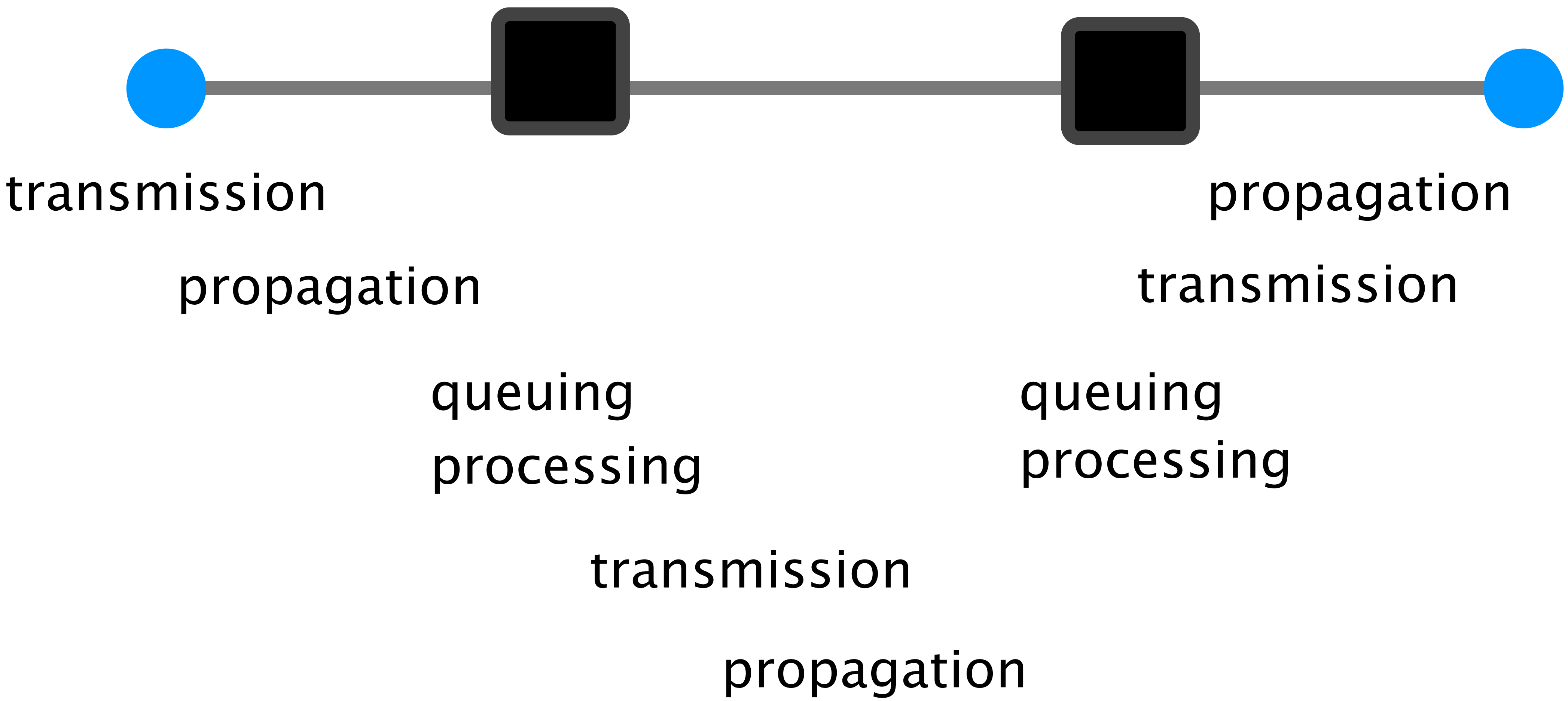


Each packet suffers from several types of delays at *each node* along the path

	transmission	delay	due to link properties
⊕	propagation	delay	
⊕	processing	delay	due to traffic mix & switch internals
⊕	queuing	delay	
=	total	delay	

Overall, the main culprits for the overall delay are the transmission, propagation and queuing delays

	transmission	delay	
⊕	propagation	delay	
⊕	processing	delay	*tend to be tiny*
⊕	queuing	delay	
=	total	delay	


transmission
propagation
queuing
processing
transmission
propagation
queuing
processing
transmission
propagation

The transmission delay is the amount of time required to push all of the bits onto the link

$$\text{Transmission delay [sec]} = \frac{\text{packet size [\#bits]}}{\text{link bandwidth [\#bits/sec]}}$$

Example

$$\frac{1000 \text{ bits}}{100 \text{ Gbps}} \quad 10 \text{ ns}$$

The propagation delay is the amount of time required for a bit to travel to the end of the link

$$\text{Propagation delay [sec]} = \frac{\text{link length [m]}}{\text{propagation speed [m/sec]}}$$

(fraction of speed of light)

Example

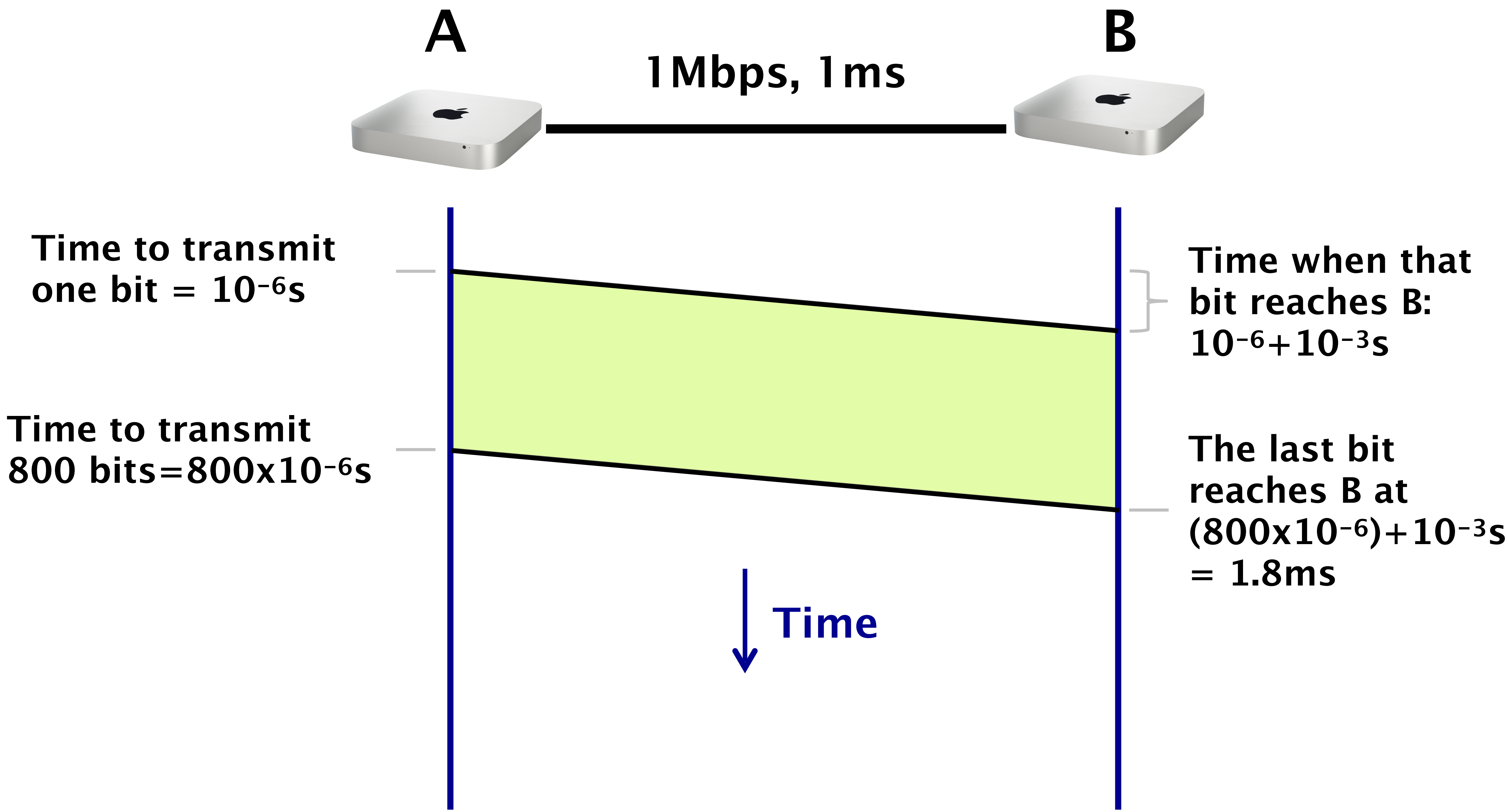
$$\frac{30\,000 \text{ m}}{2 \times 10^8 \text{ m/sec}} \quad 150\ \mu\text{sec}$$

(speed of light in fiber)

How long does it take for a packet to travel from A to B?

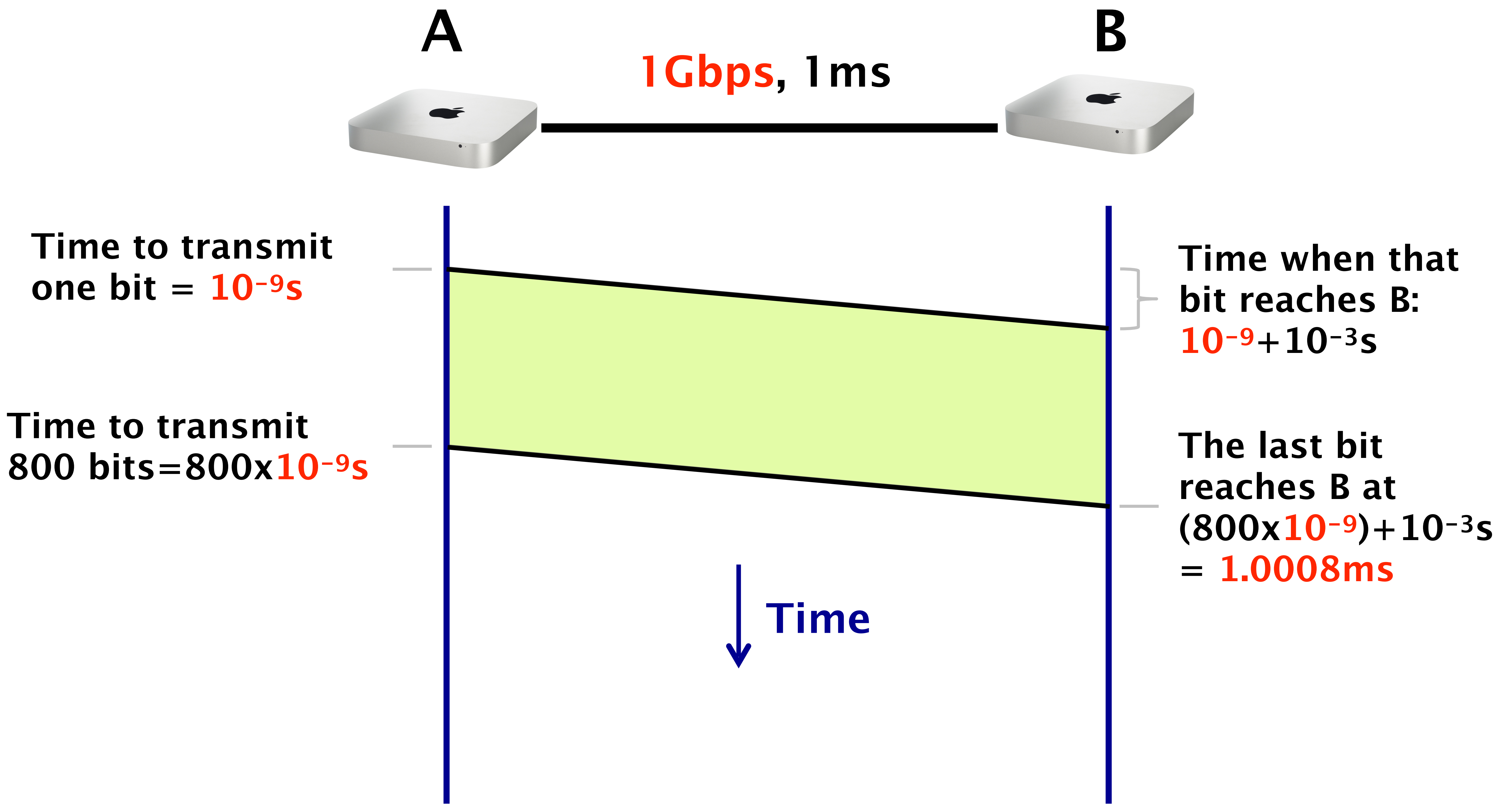
(not considering queuing for now)

How long does it take to exchange 100 Bytes packet?

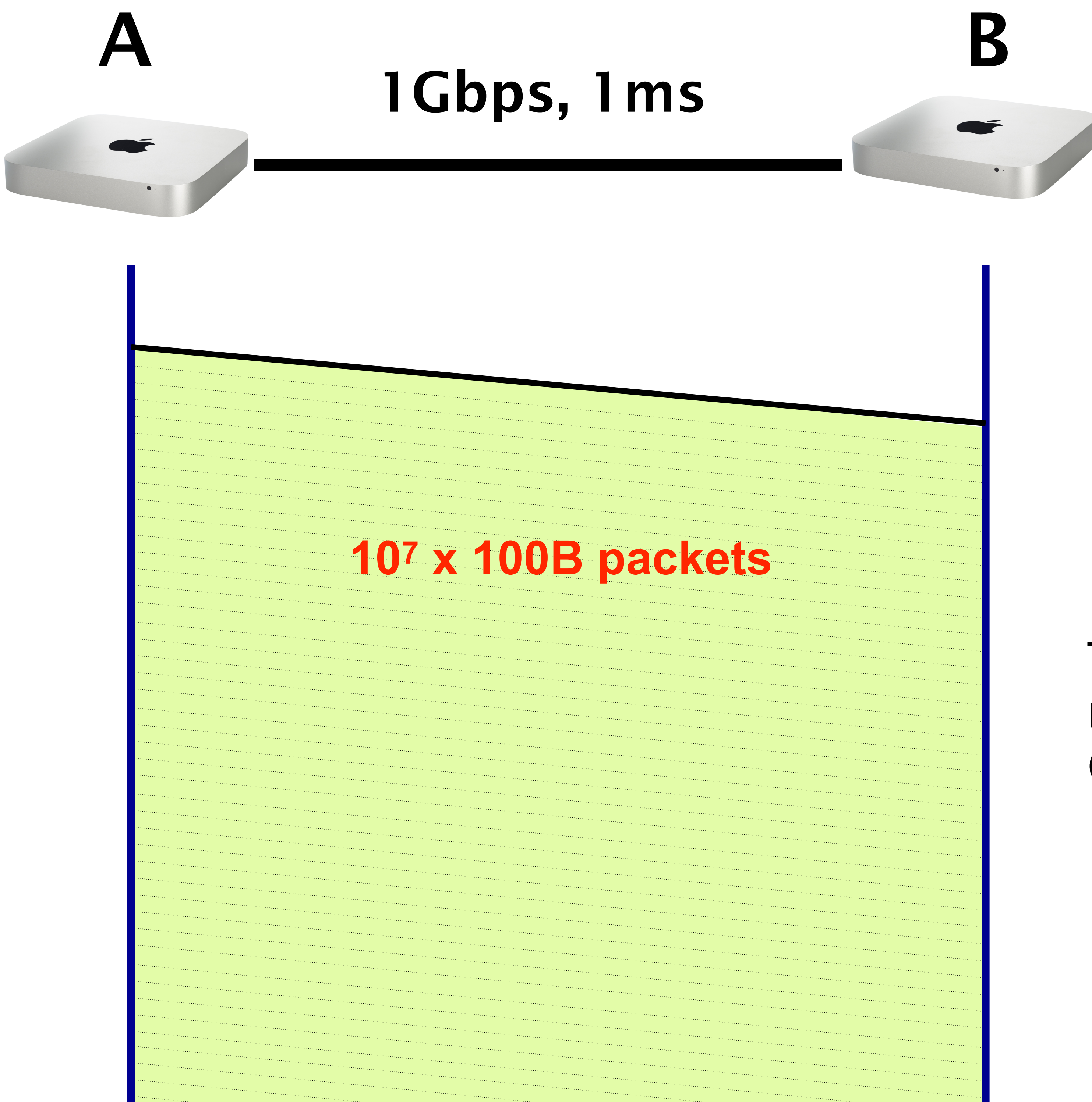


If we have a 1 Gbps link,

the total time decreases to 1.0008ms



If we now exchange a 1GB file split in 100B packets



The last bit reaches B at $(10^7 \times 800 \times 10^{-9}) + 10^{-3}s$ = 8001ms

Different transmission characteristics imply different tradeoffs in terms of which delay dominates

10^7x100B	pkt	1Gbps link	transmission delay dominates
1x100B	pkt	1Gbps link	propagation delay dominates
1x100B	pkt	1Mbps link	both matter

In the Internet, we **can't know** in advance which one matters!

The queuing delay is the amount of time a packet waits (in a buffer) to be transmitted on a link

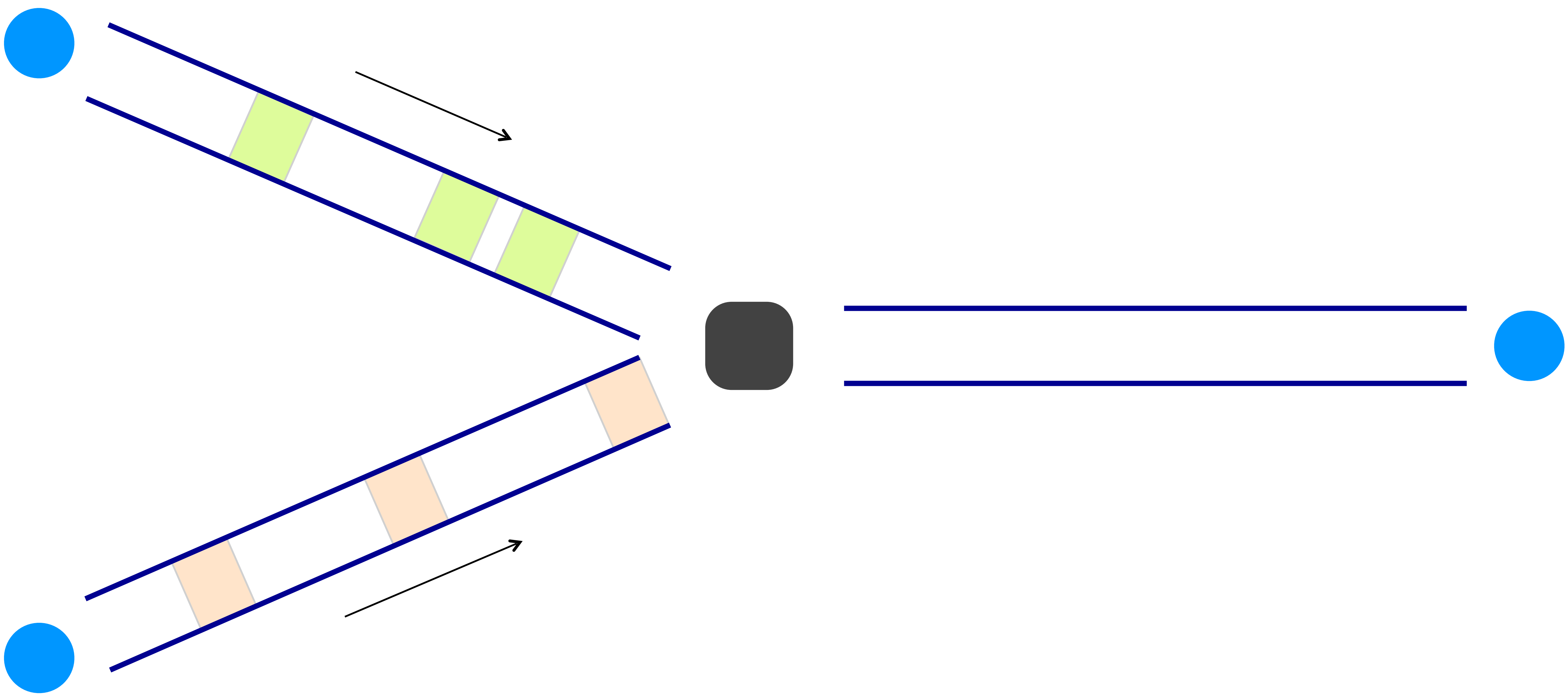
Queuing delay is the hardest to evaluate

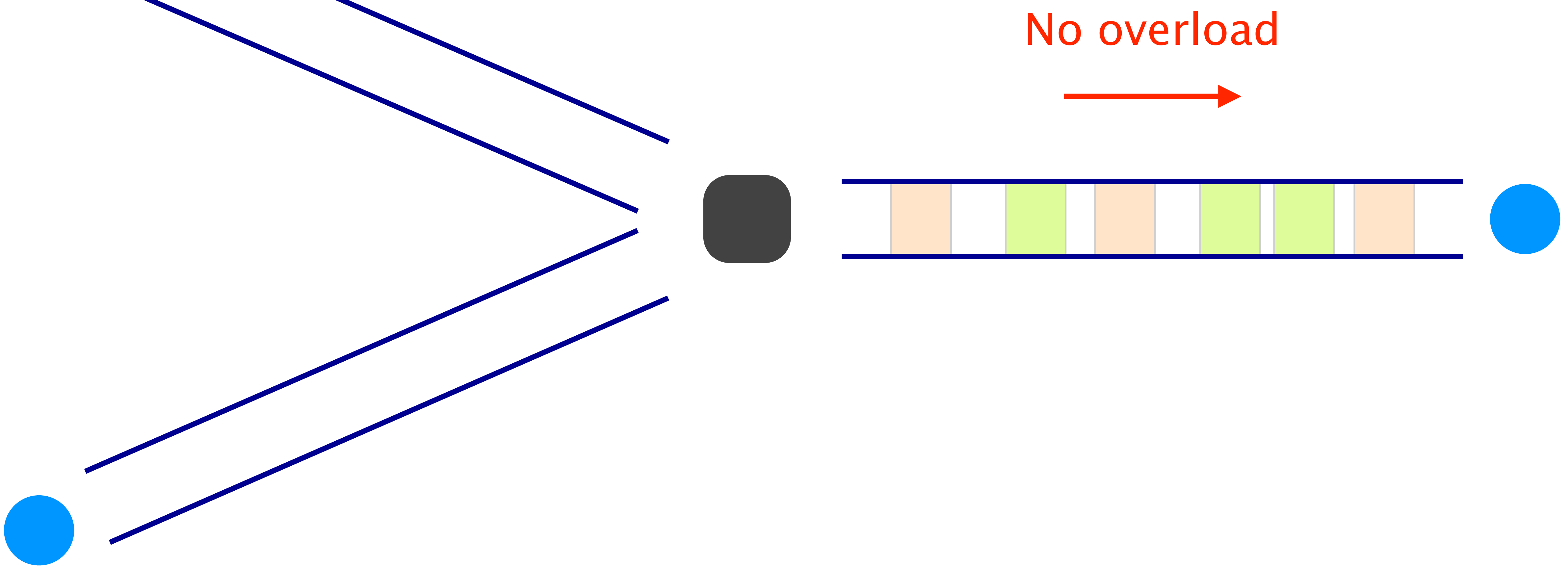
as it varies from packet to packet

It is characterized with statistical measures

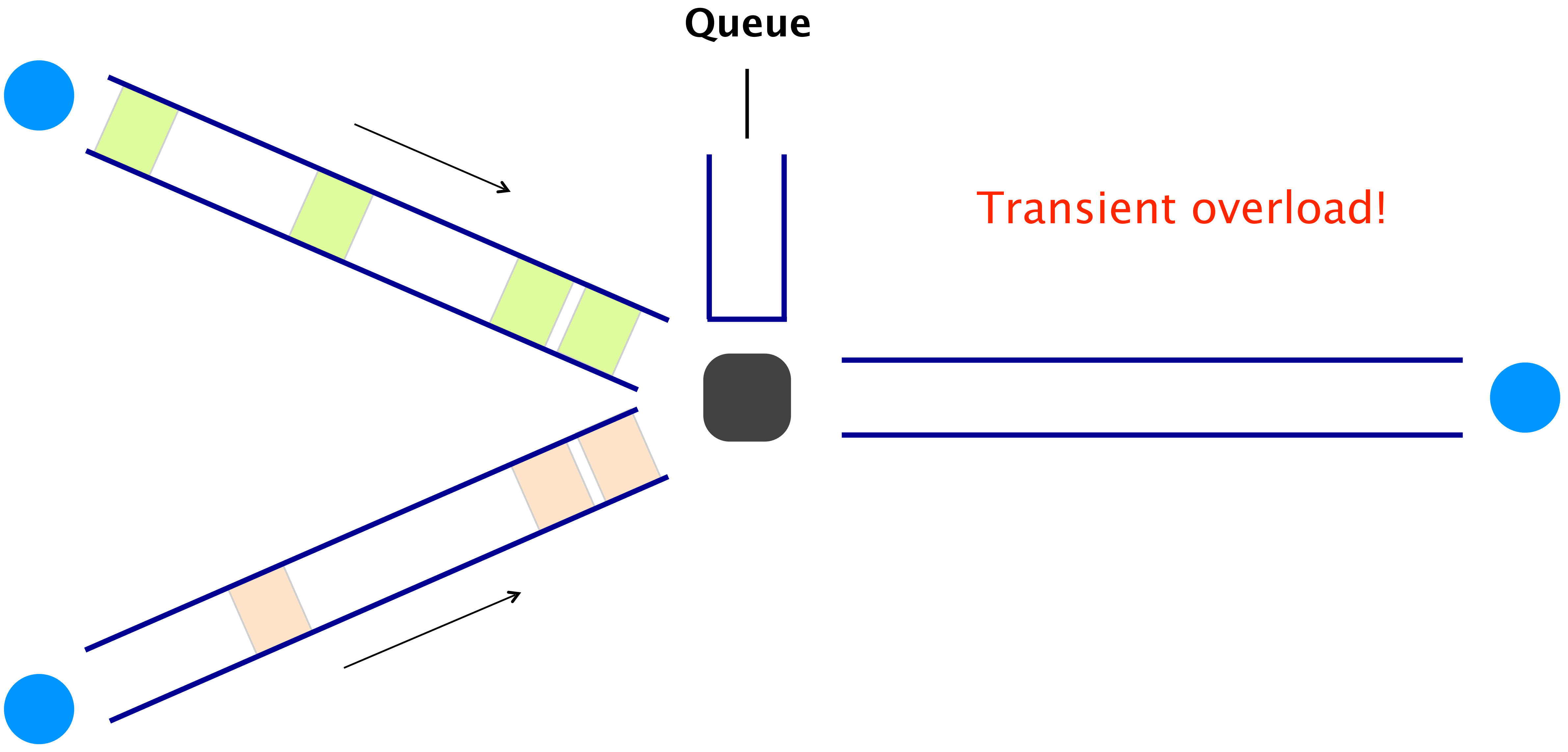
e.g., average delay & variance, probability of exceeding *x*

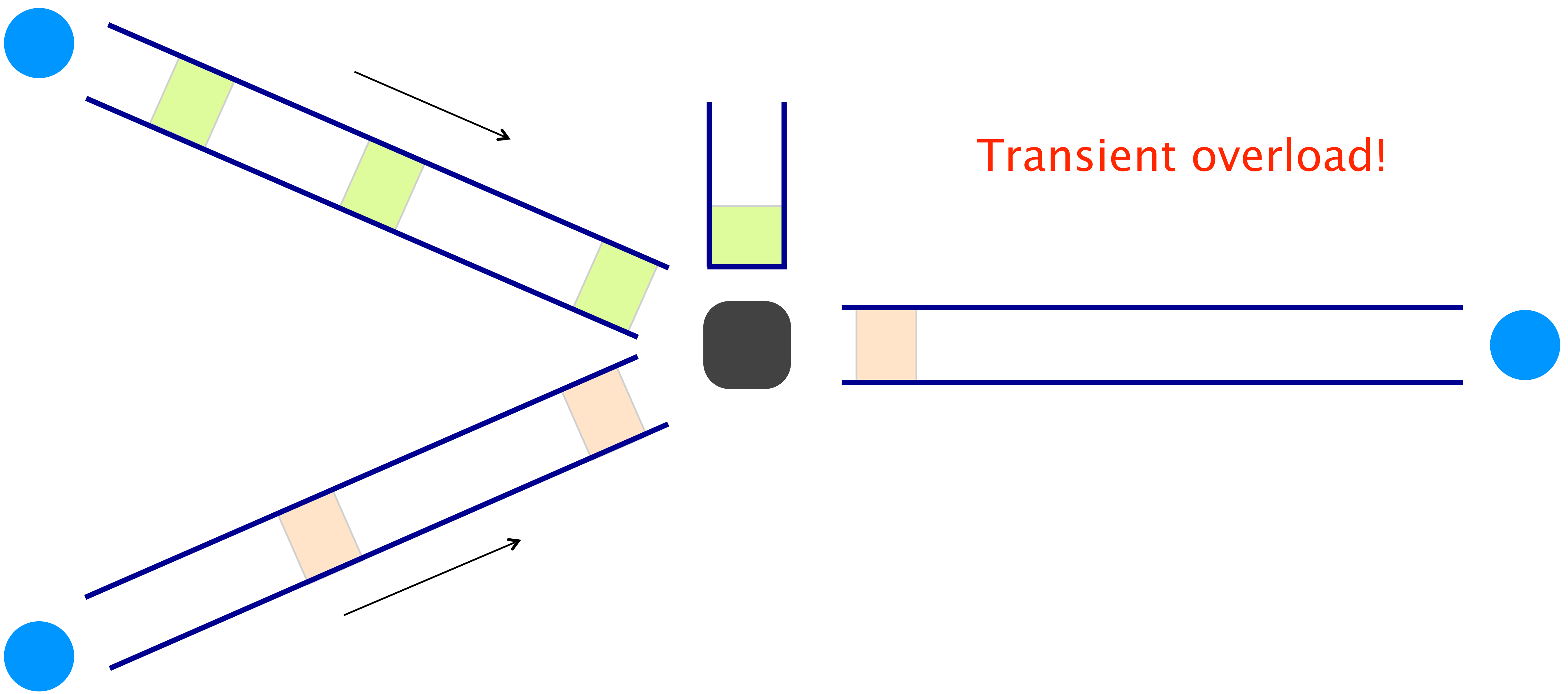
Queuing delay depends on the traffic pattern

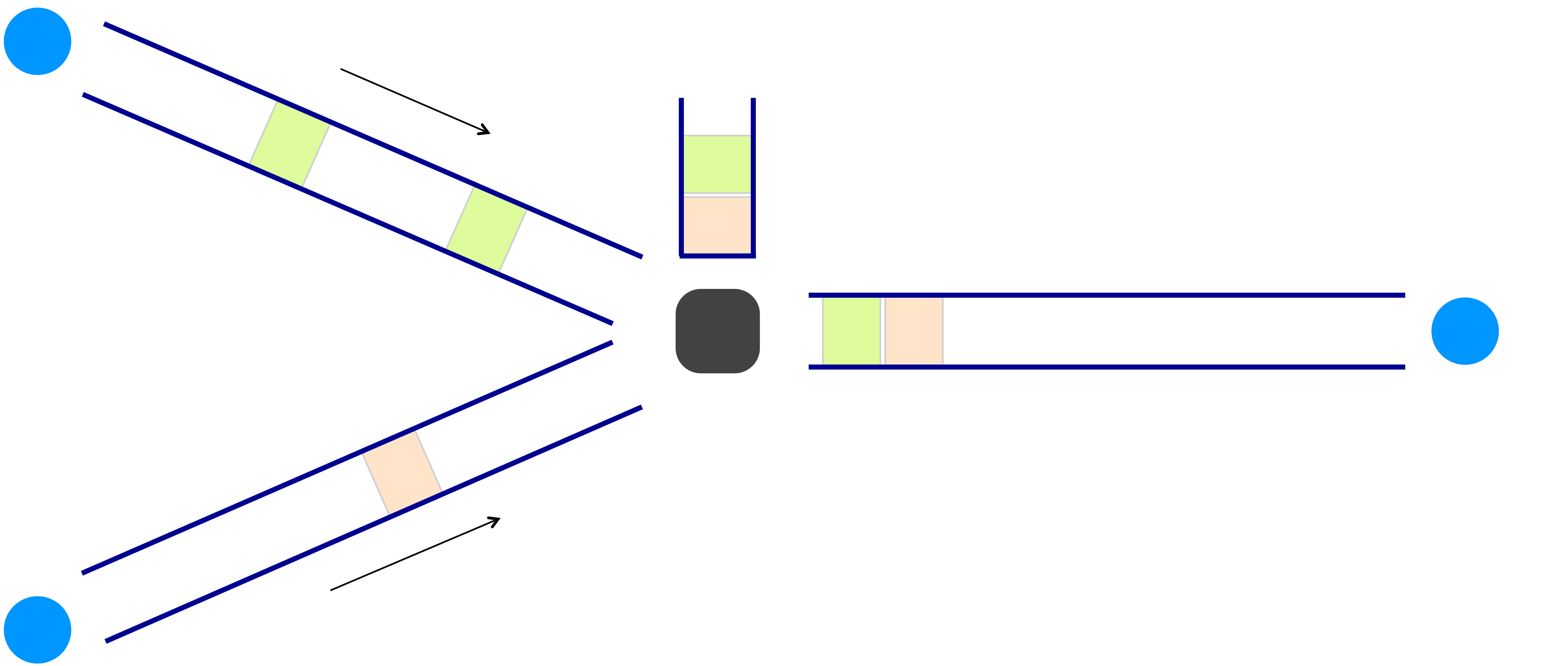


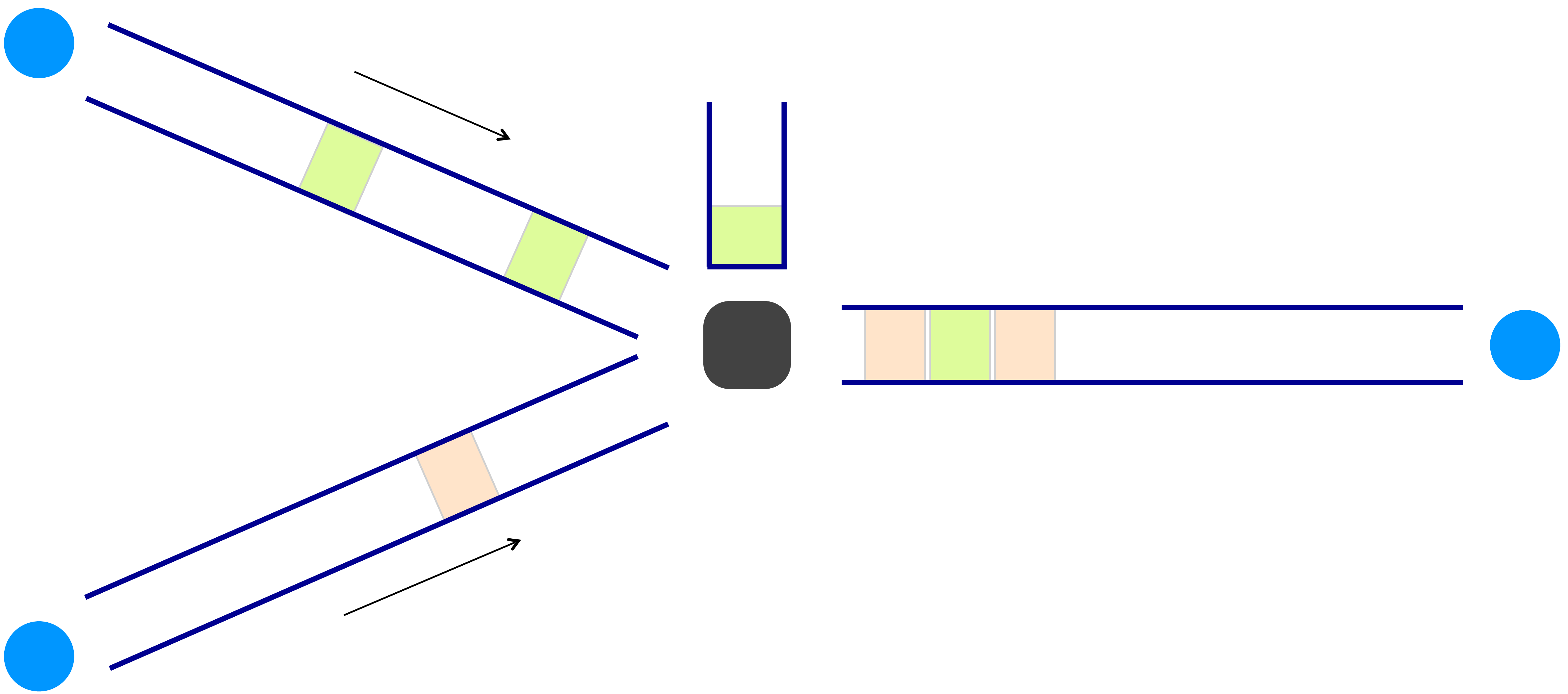

No overload

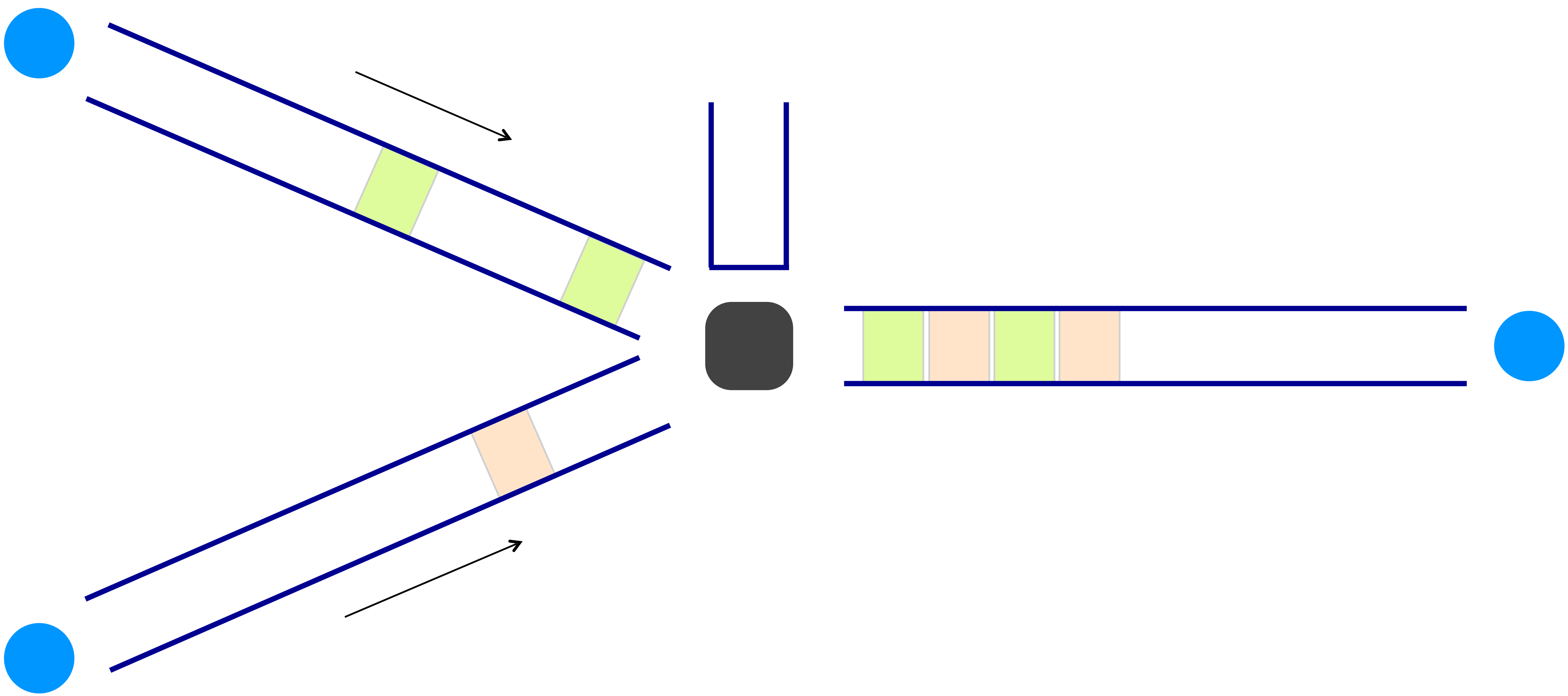
Queuing delay depends on the traffic pattern



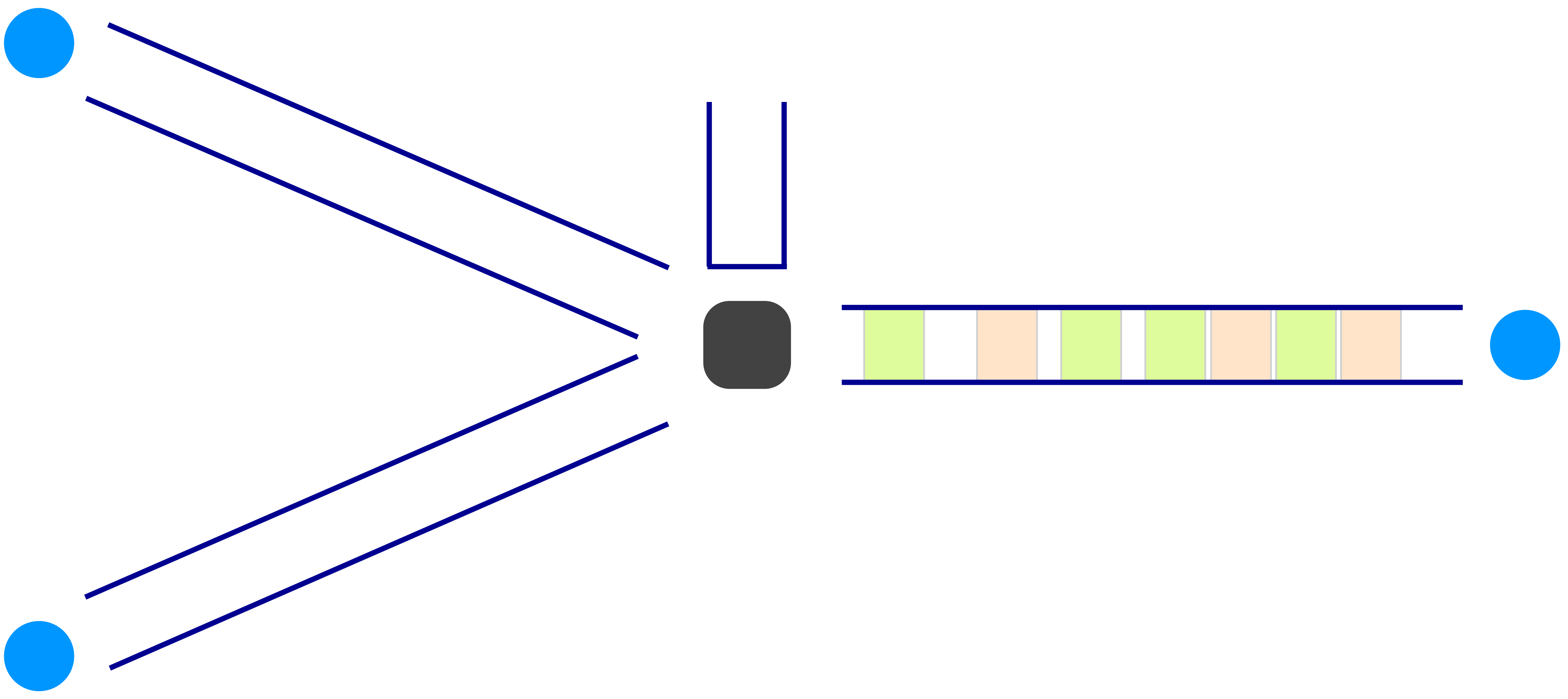

Transient overload!







Queues absorb transient bursts,
but introduce queueing delays



The time a packet has to sit in a buffer before being processed depends on the traffic pattern

Queueing delay depends on:

- arrival rate at the queue
- transmission rate of the outgoing link
- traffic burstiness

average packet arrival rate	a	[packet/sec]
transmission rate of outgoing link	R	[bit/sec]
fixed packets length	L	[bit]
average bits arrival rate	La	[bit/sec]
traffic intensity	La/R	

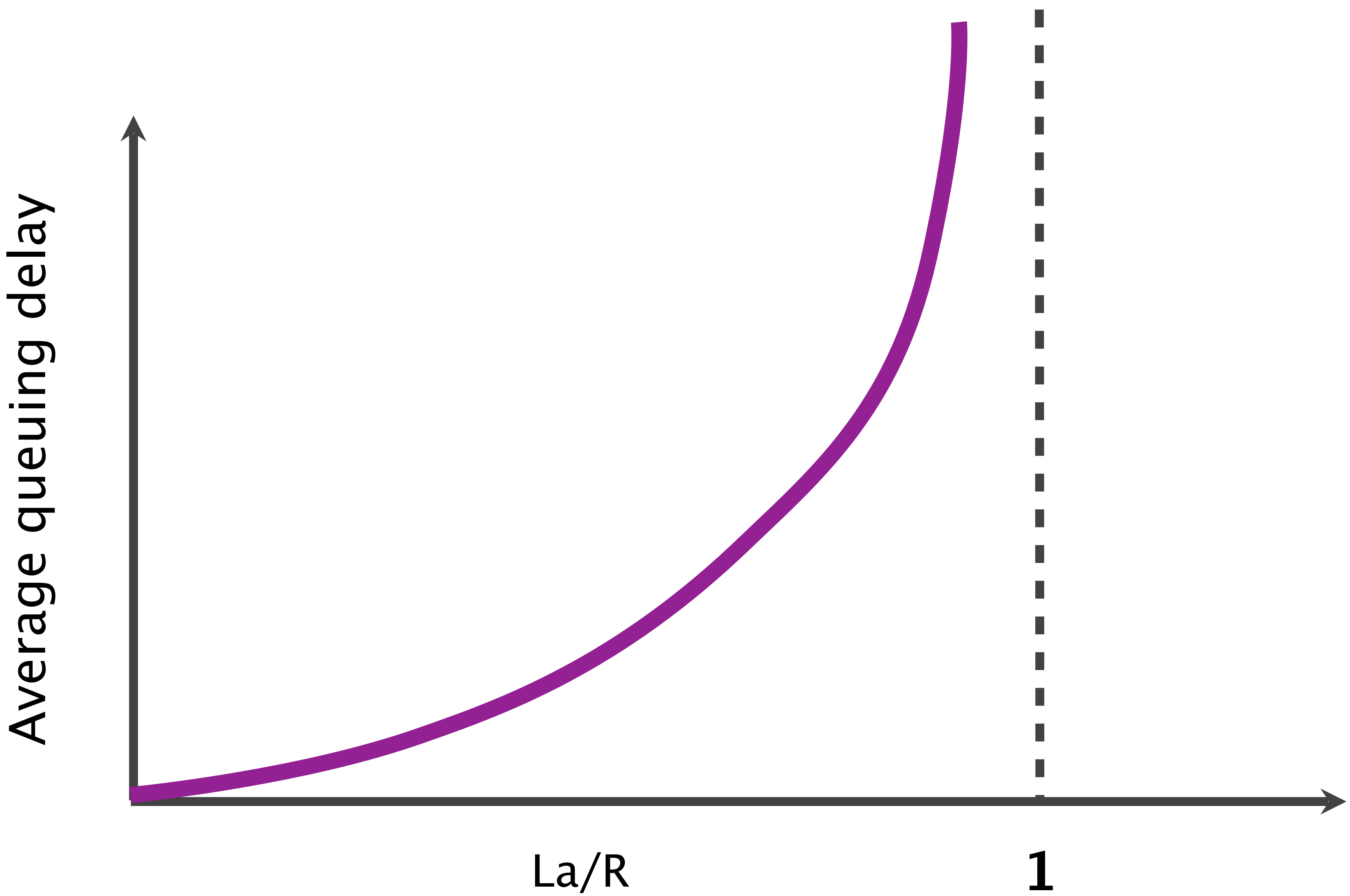
When the traffic intensity is >1, the queue will increase without bound, and so does the queuing delay

Golden rule

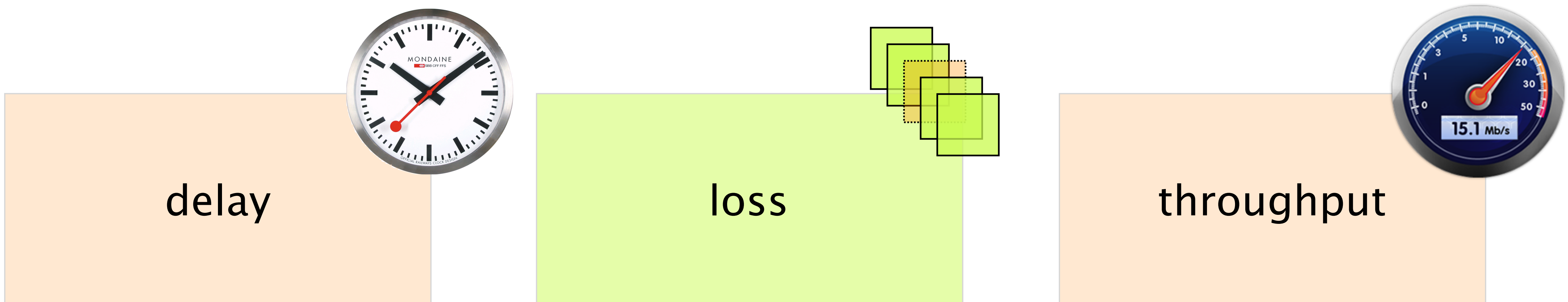
Design your queuing system,

so that it operates far from that point

When the traffic intensity is <=1,
queueing delay depends on the burst size

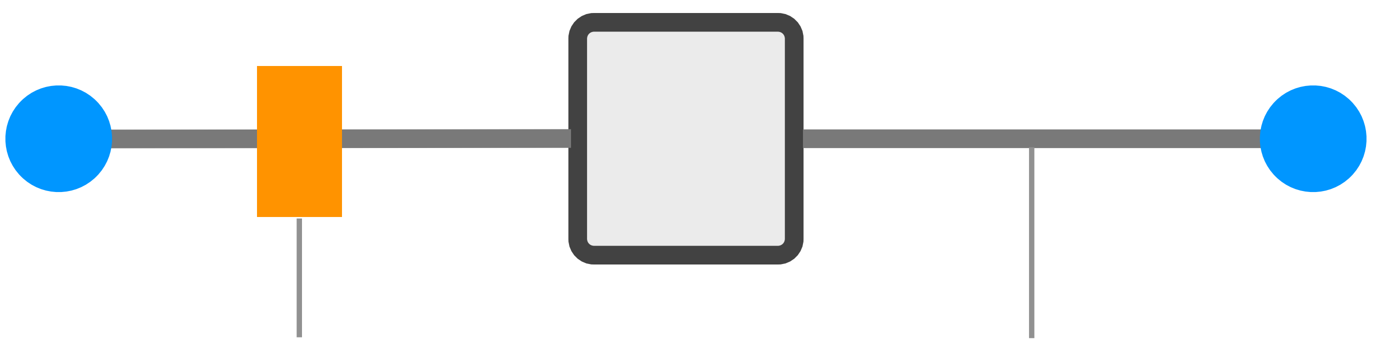


A network *connection* is characterized by its delay, loss rate and throughput



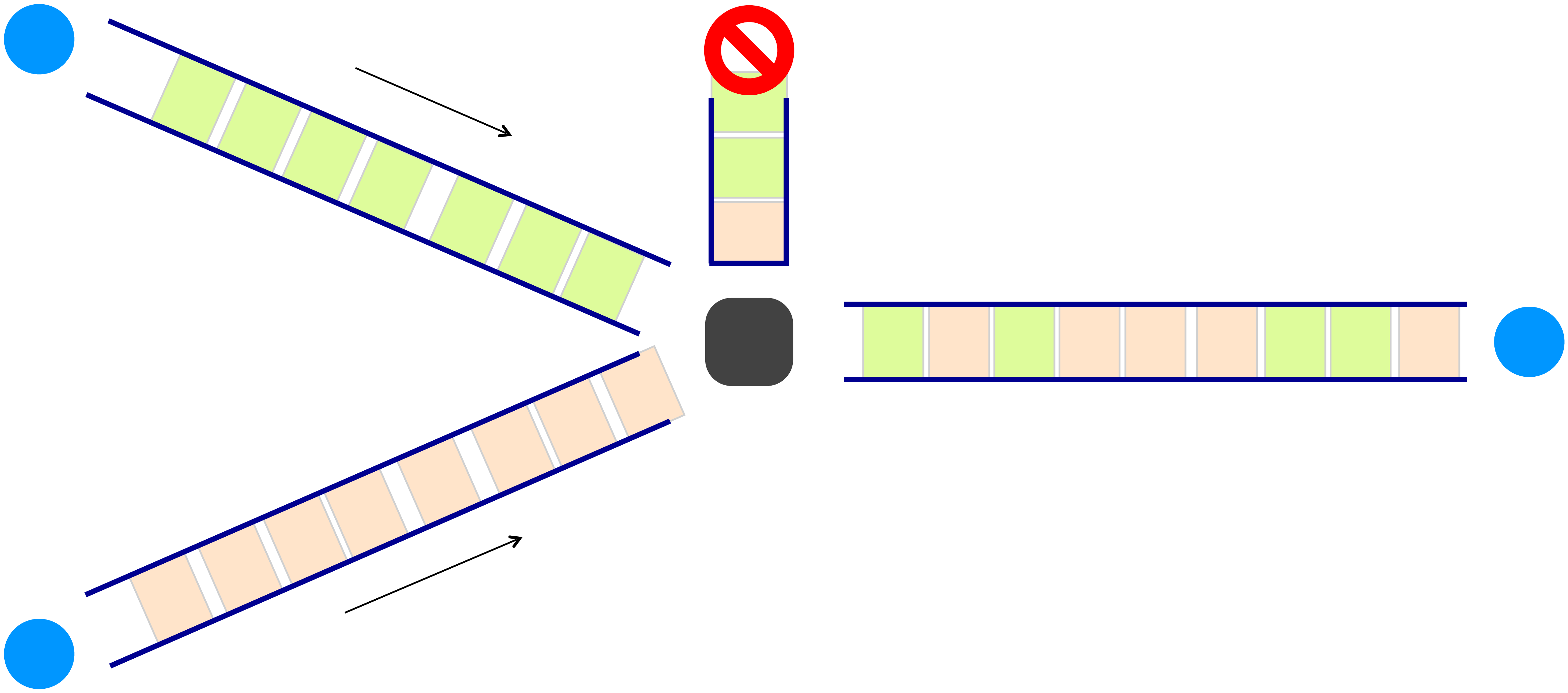
In practice, queues are not infinite.

There is an upper bound on queuing delay.

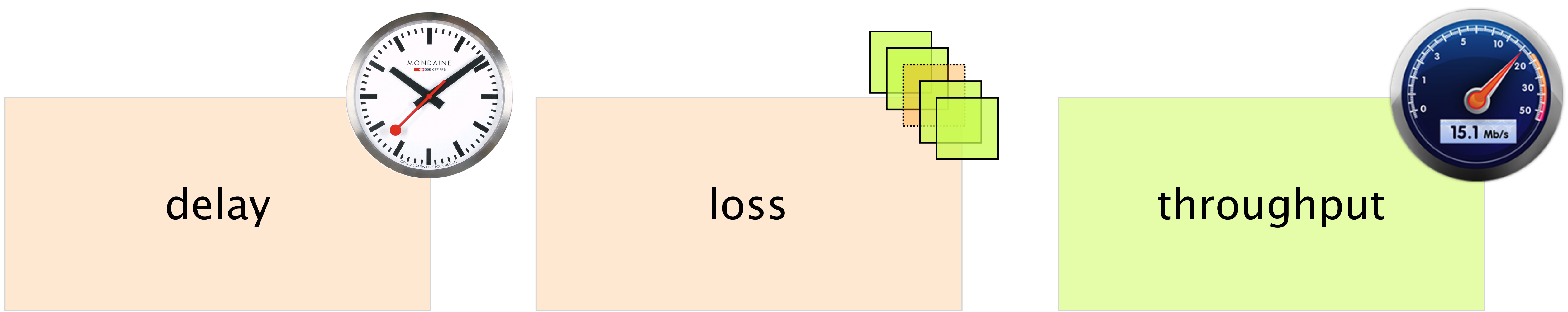


queuing delay upper bound: N*L/R

If the queue is persistently overloaded,
it will eventually drop packets (loss)



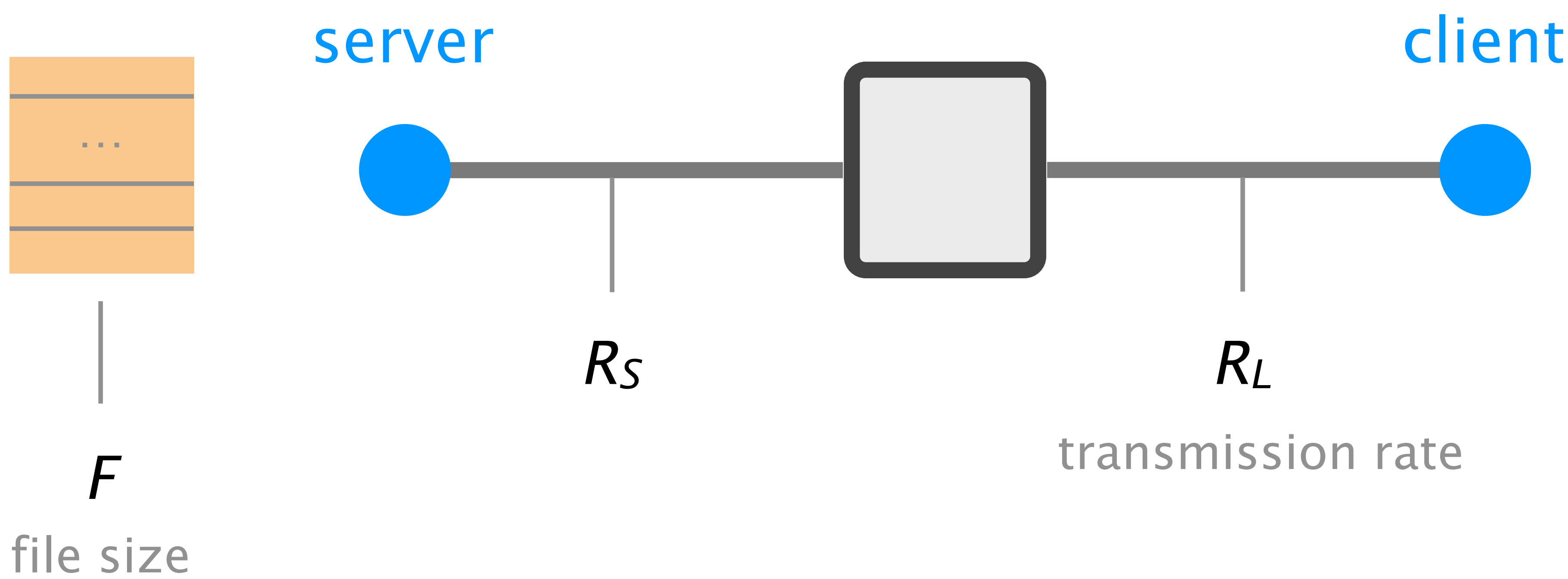
A network *connection* is characterized by its delay, loss rate and throughput



The throughput is the instantaneous rate at which a host receives data

$$\underset{[\#bits/sec]}{\text{Average throughput}} = \frac{\text{data size} \quad [\#bits]}{\text{transfer time} \quad [sec]}$$

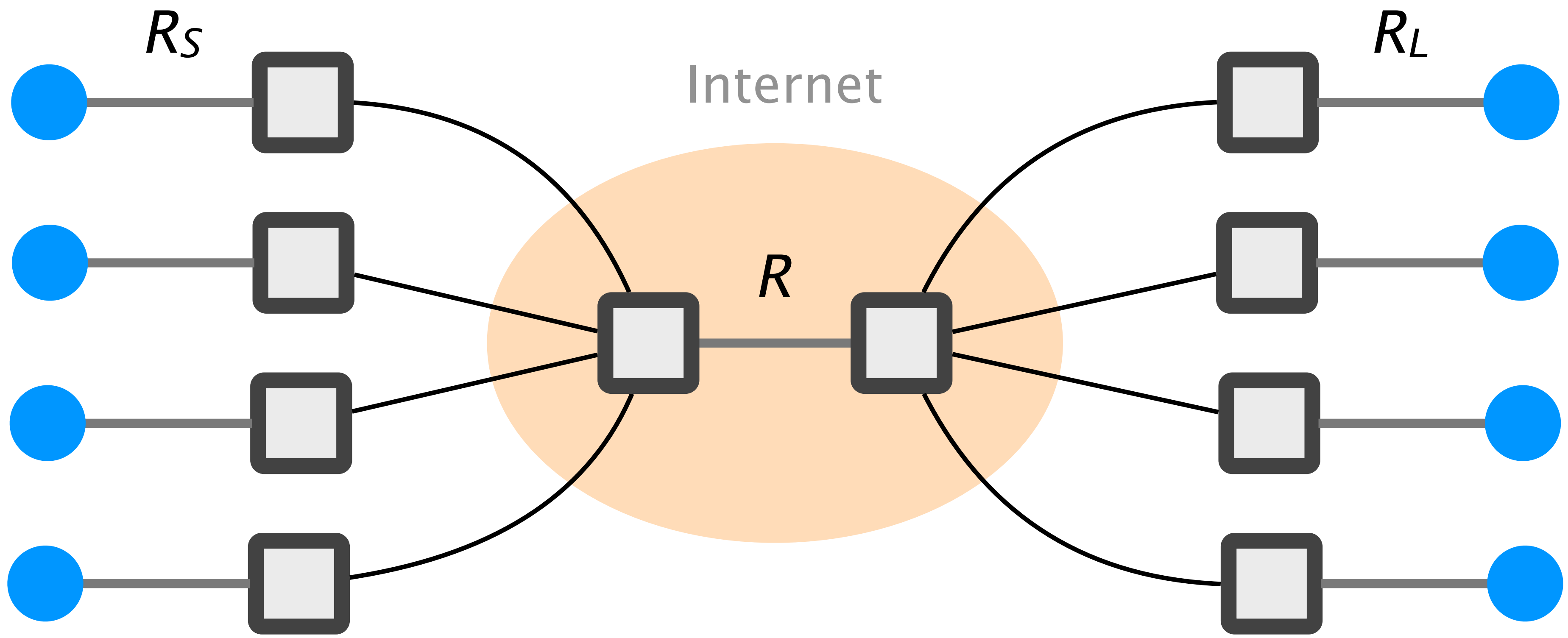
To compute throughput, one has to consider the bottleneck link



Average throughput $\min(R_S, R_L)$

= transmission rate of the bottleneck link

To compute throughput, one has to consider the bottleneck link… **and the intervening traffic**

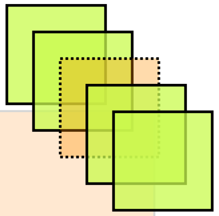


if $4*\min(R_S, R_L) > R$

the bottleneck is now in the core, providing each download R/4 of throughput

A network *connection* is characterized by its delay, loss rate and throughput

delay

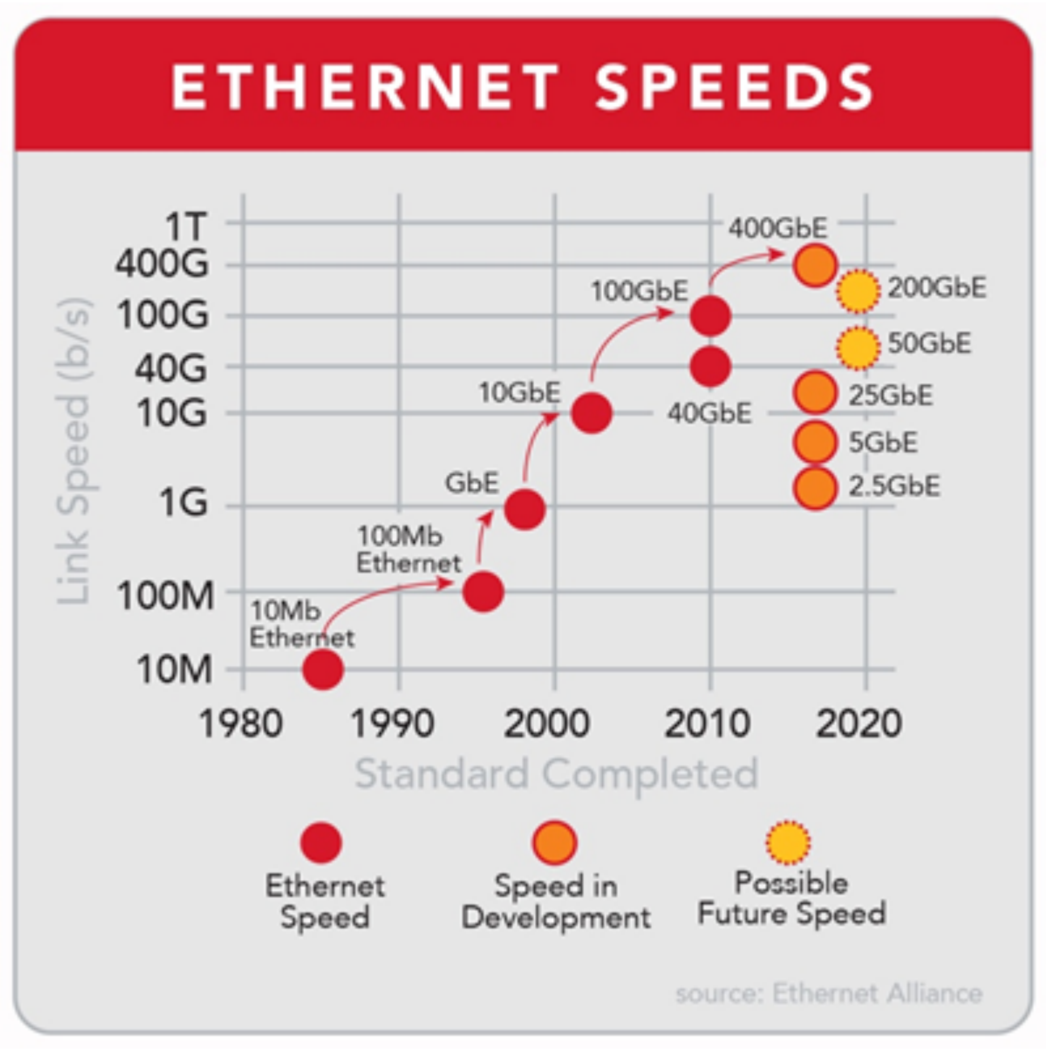


loss

throughput

As technology improves, throughput increase & delays are getting lower except for propagation

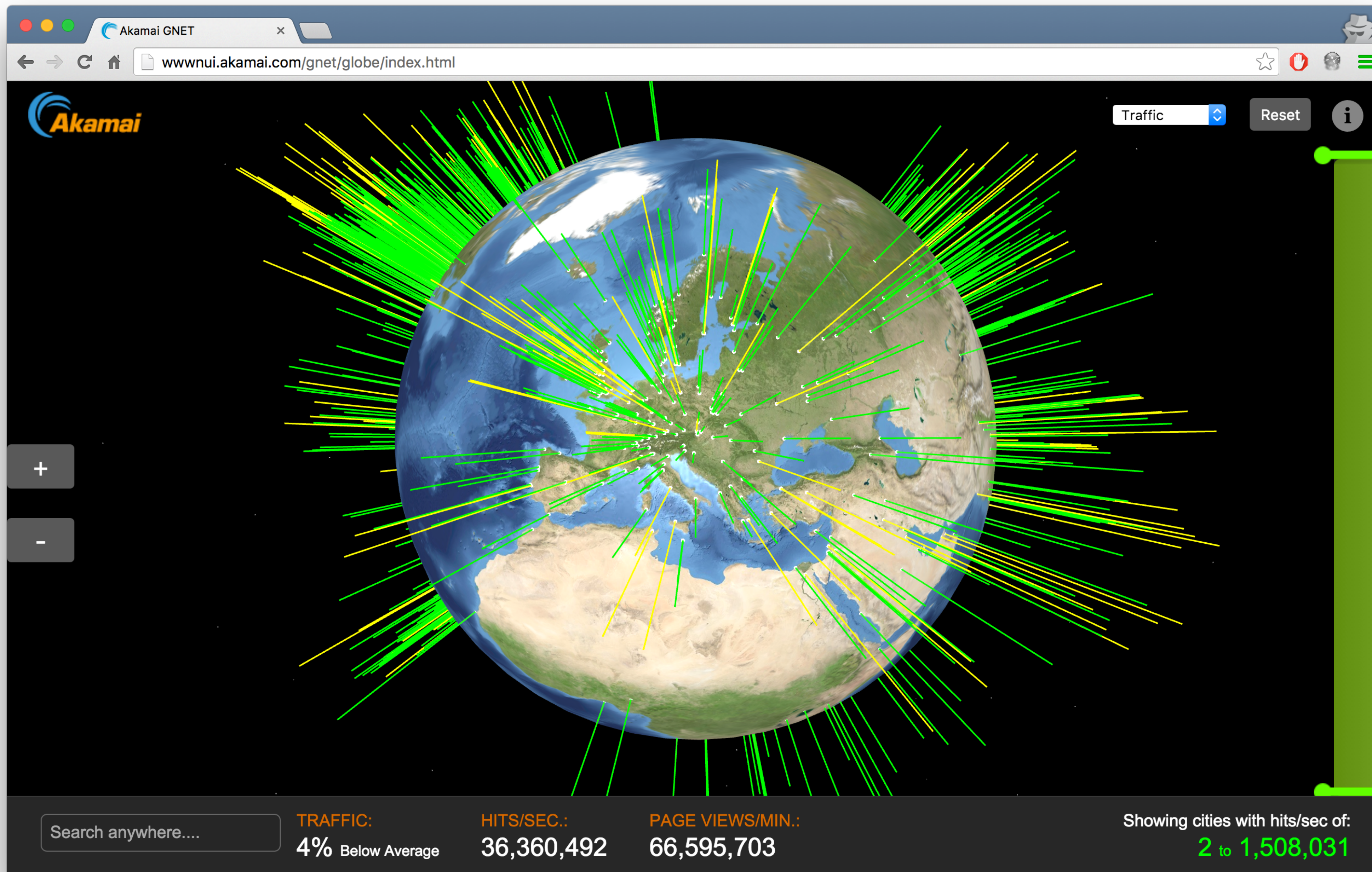
(speed of light)



source: ciena.com

Because of propagation delays,

Content Delivery Networks move content closer to you



http://wwwnui.akamai.com/gnet/globe/index.html

Communication Networks

Part 1: General overview

What is a network made of?

How is it shared?

How is it organized?

How does communication happen?

How do we characterize it?

Communication Networks

Part 2: Concepts

routing

reliable delivery

Communication Networks

Part 2: Concepts

routing

How do you guide IP packets from a source to destination?

reliable delivery

How do you ensure reliable transport on top of best-effort delivery?

This week

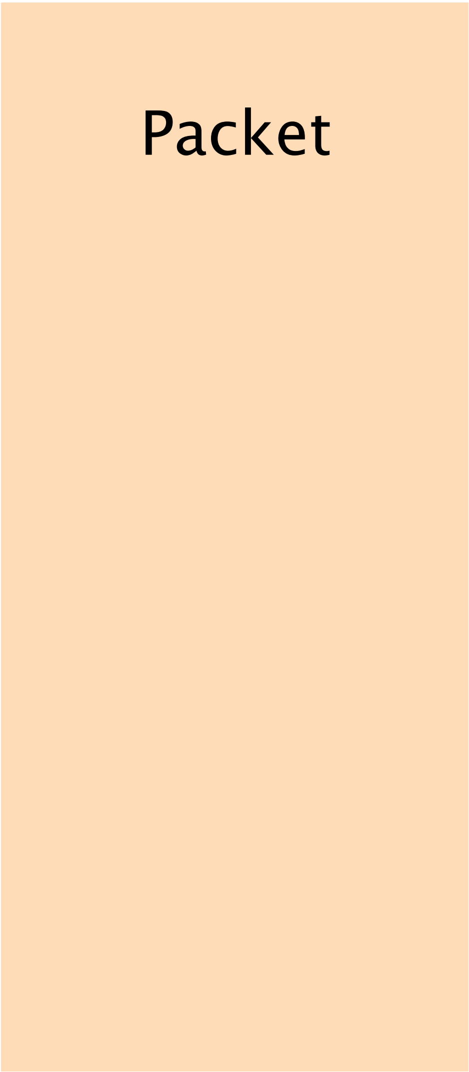
routing

Next week

reliable delivery

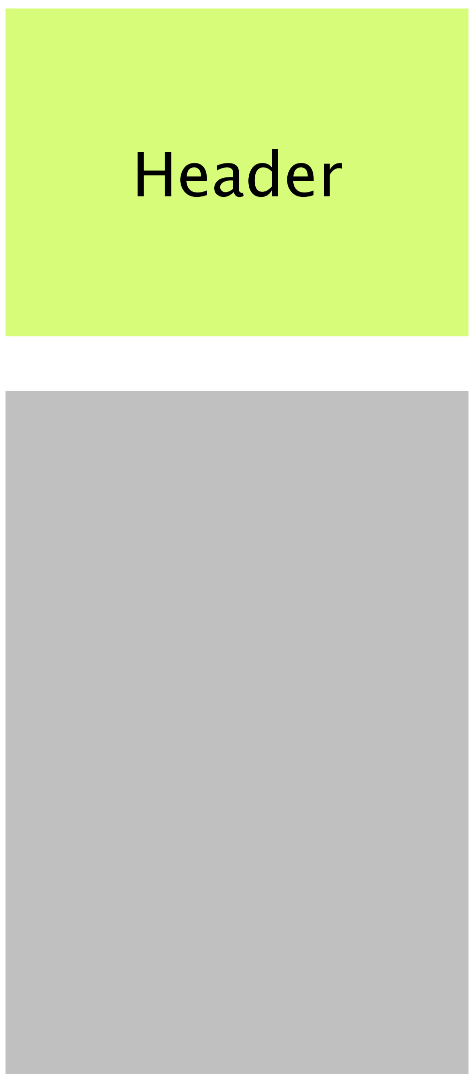
How do you guide IP packets from a source to destination?

Think of IP packets as envelopes



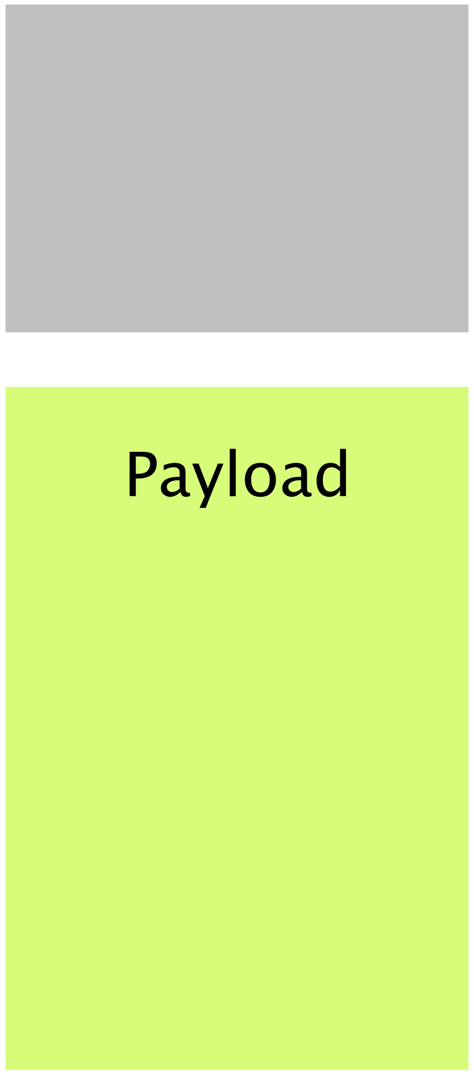
Like an envelope,

packets have a header

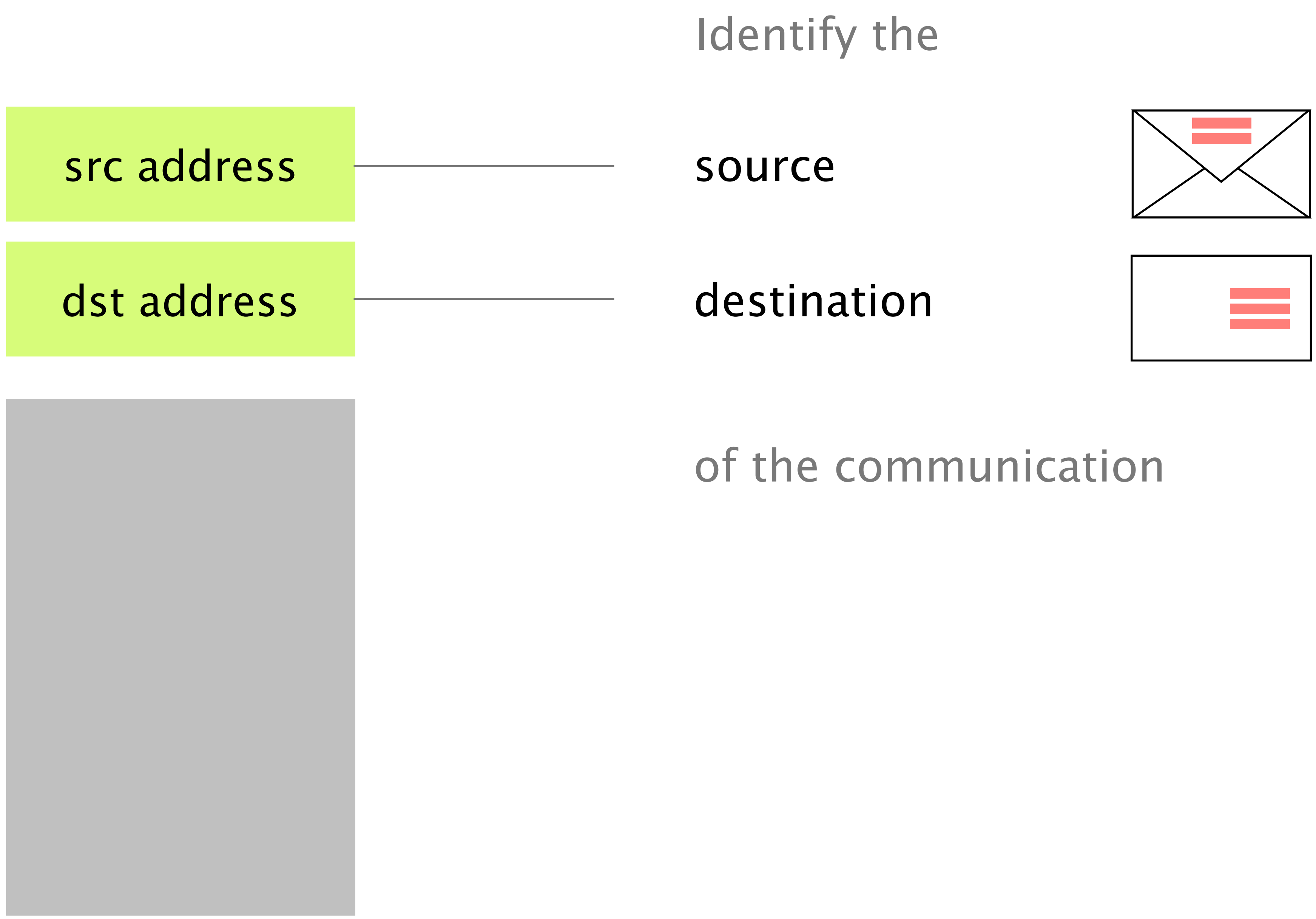


Like an envelope,

packets have a payload

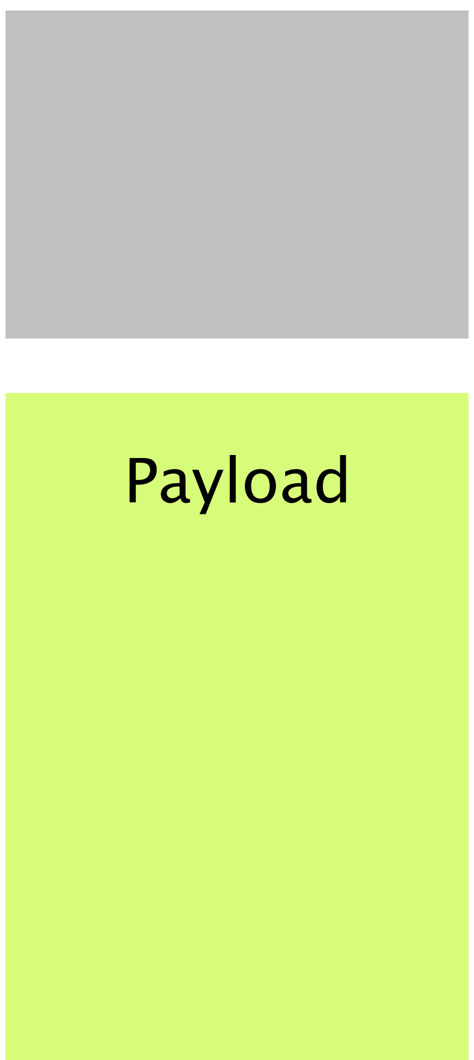


The **header** contains the **metadata** needed **to forward** the packet

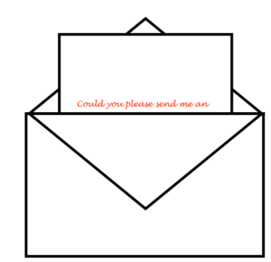


The **payload** contains

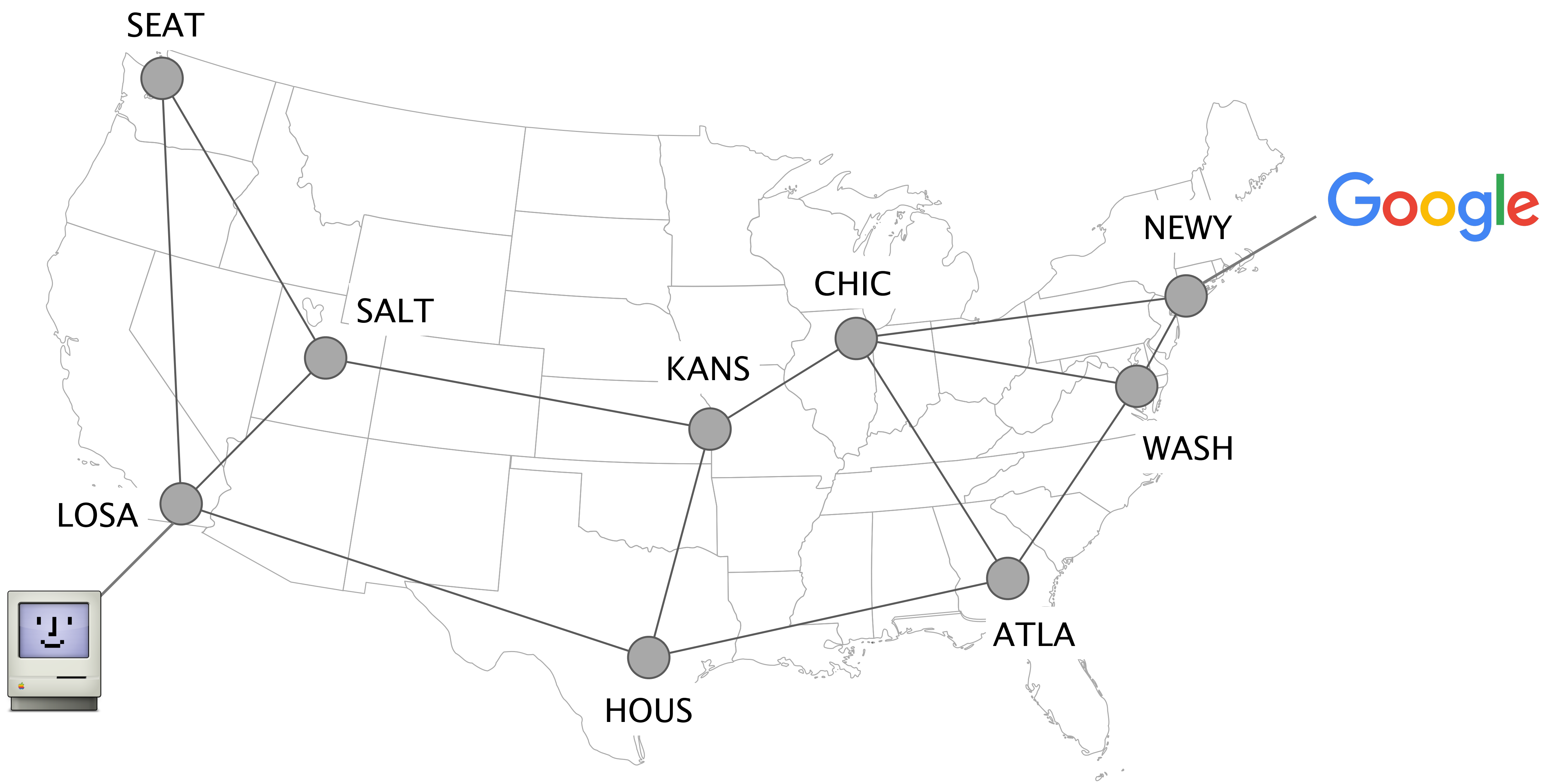
the **data to be delivered**

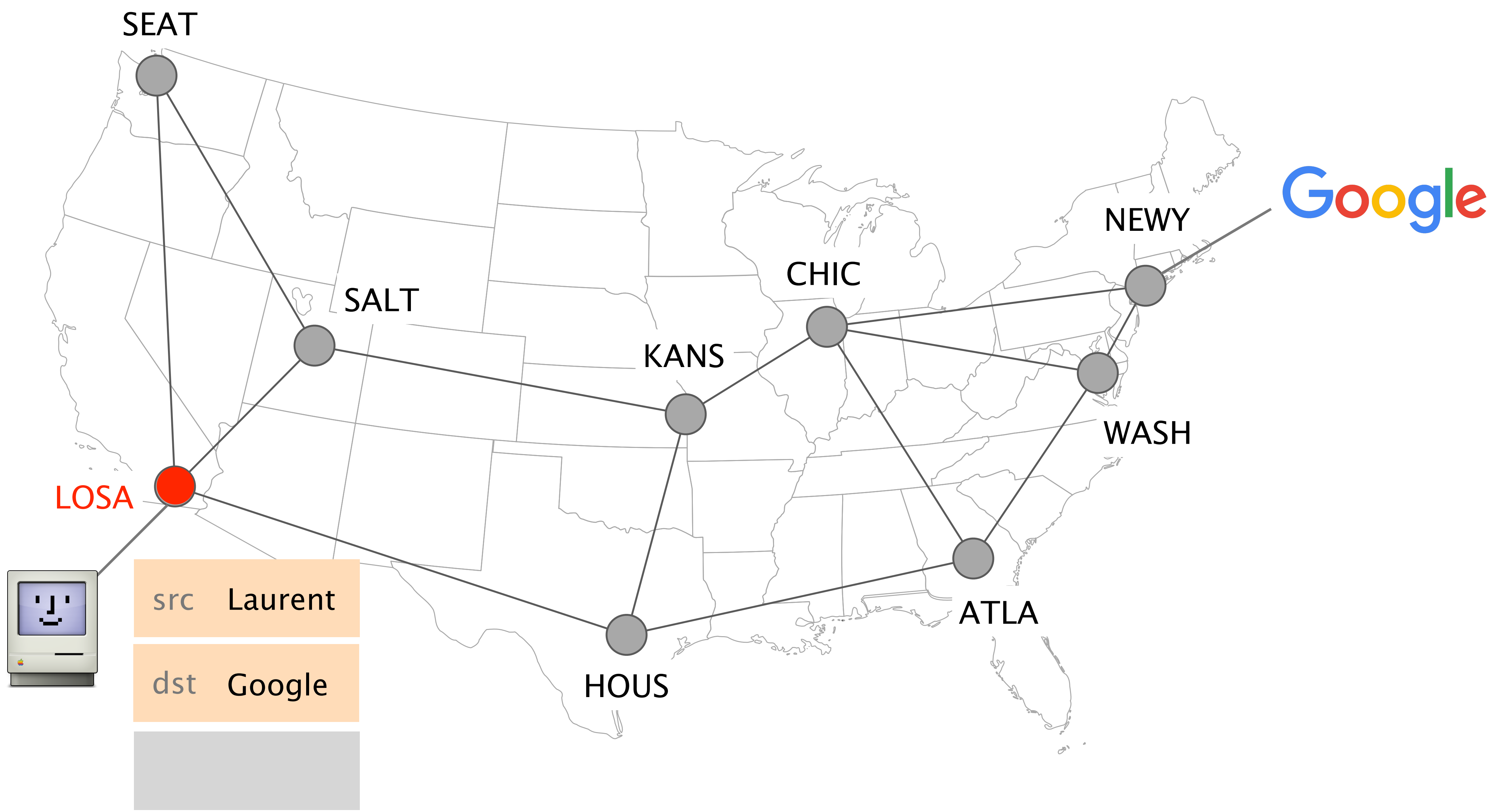



```
<html><head>
<meta http-equiv="content-type" content="text/html; charset=UTF-8">
<title>Google</title>
</head><body>
  <img alt="Google" height=110 src="images/logo.gif" width=276>
  <form action="/search" name=f>
   <input name=hl type=hidden value=en>
   <input name=q size=55 title="Google Search" value="">
   <input name=btnG type=submit value="Google Search">
   <input name=btnI type=submit value="I'm Feeling Lucky">
 </form>
</body></html>
```

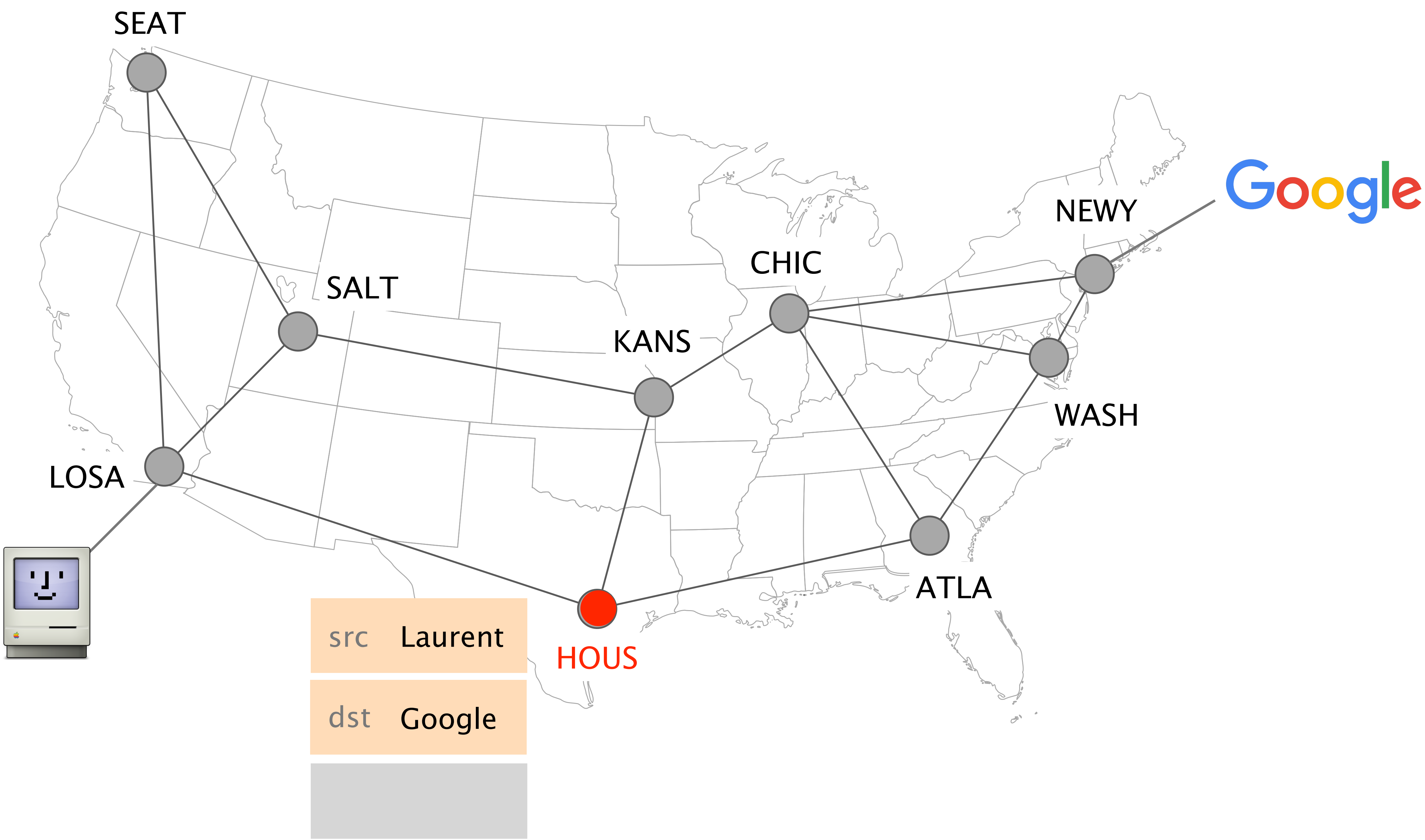


Routers forward IP packets **hop-by-hop** towards their destination





SEAT
SALT
KANS
CHIC
NEWY
Google
WASH
LOSA
ATLA
HOUS
src Laurent
dst Google


SEAT
SALT
LOSA
KANS
CHIC
NEWY
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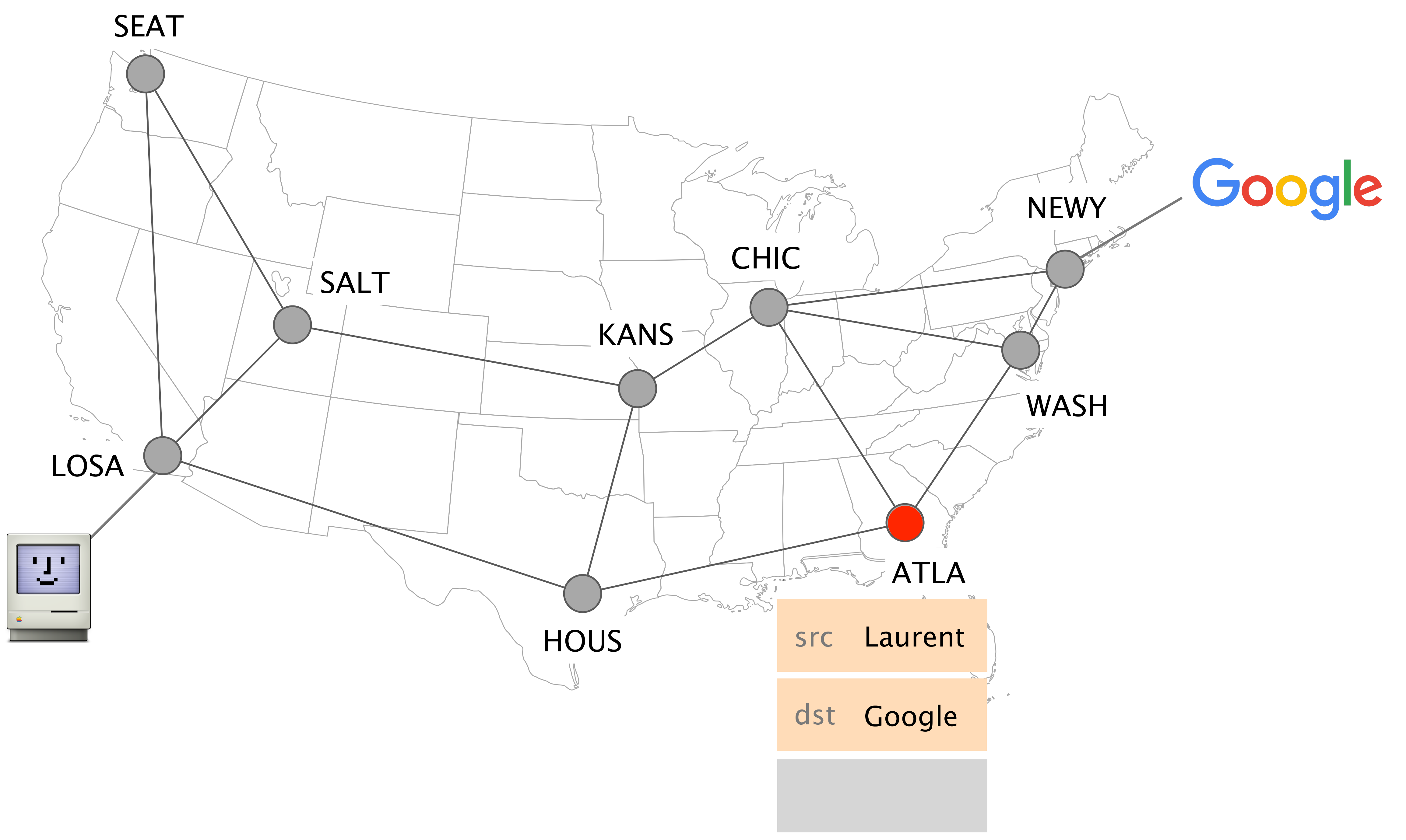
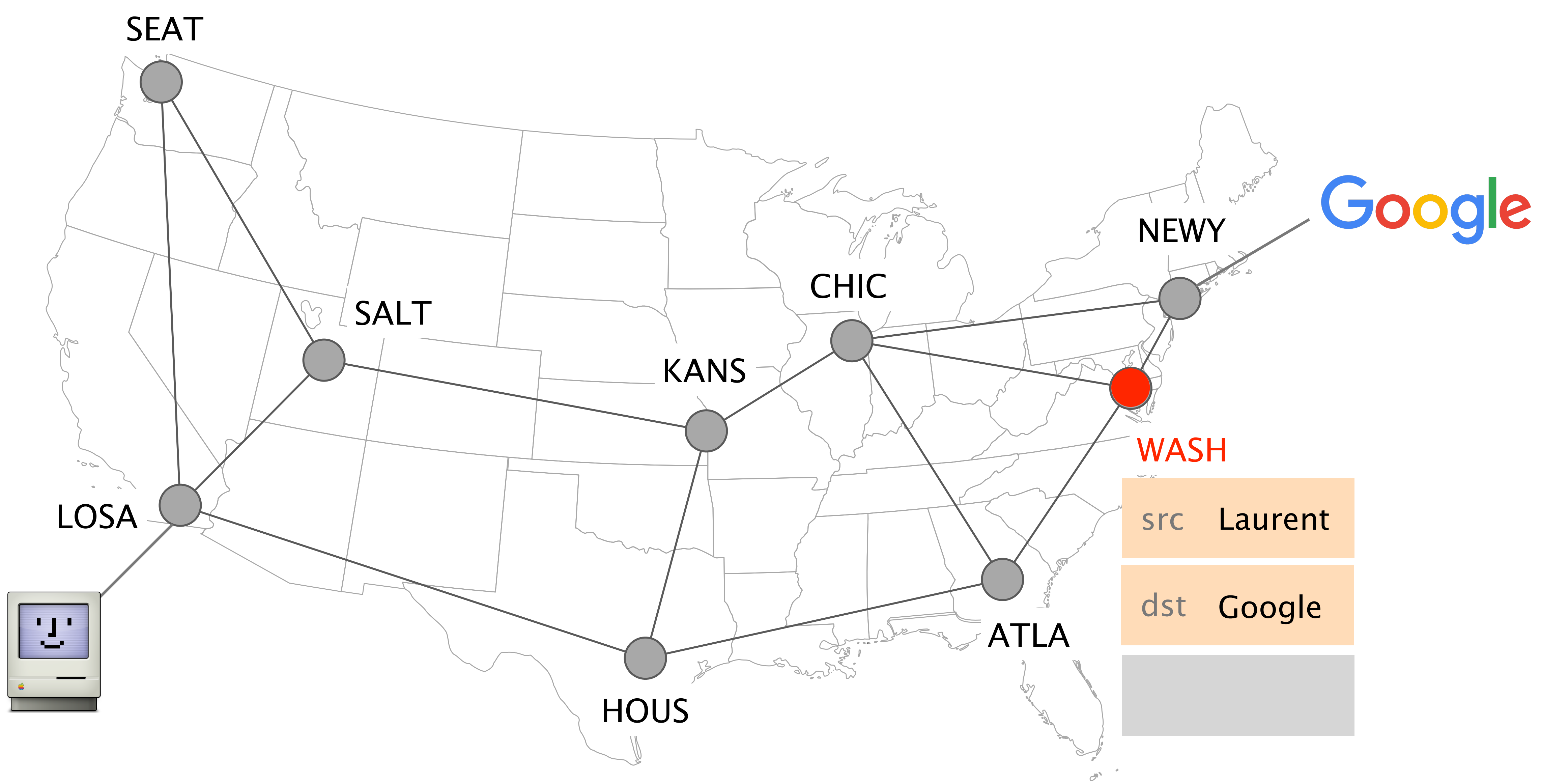
HOUS
src Laurent
dst Google


SEAT
SALT
KANS
CHIC
NEWY
WASH
LOSA
HOUS
ATLA
Google
src Laurent
dst Google


SEAT
SALT
KANS
CHIC
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src Laurent
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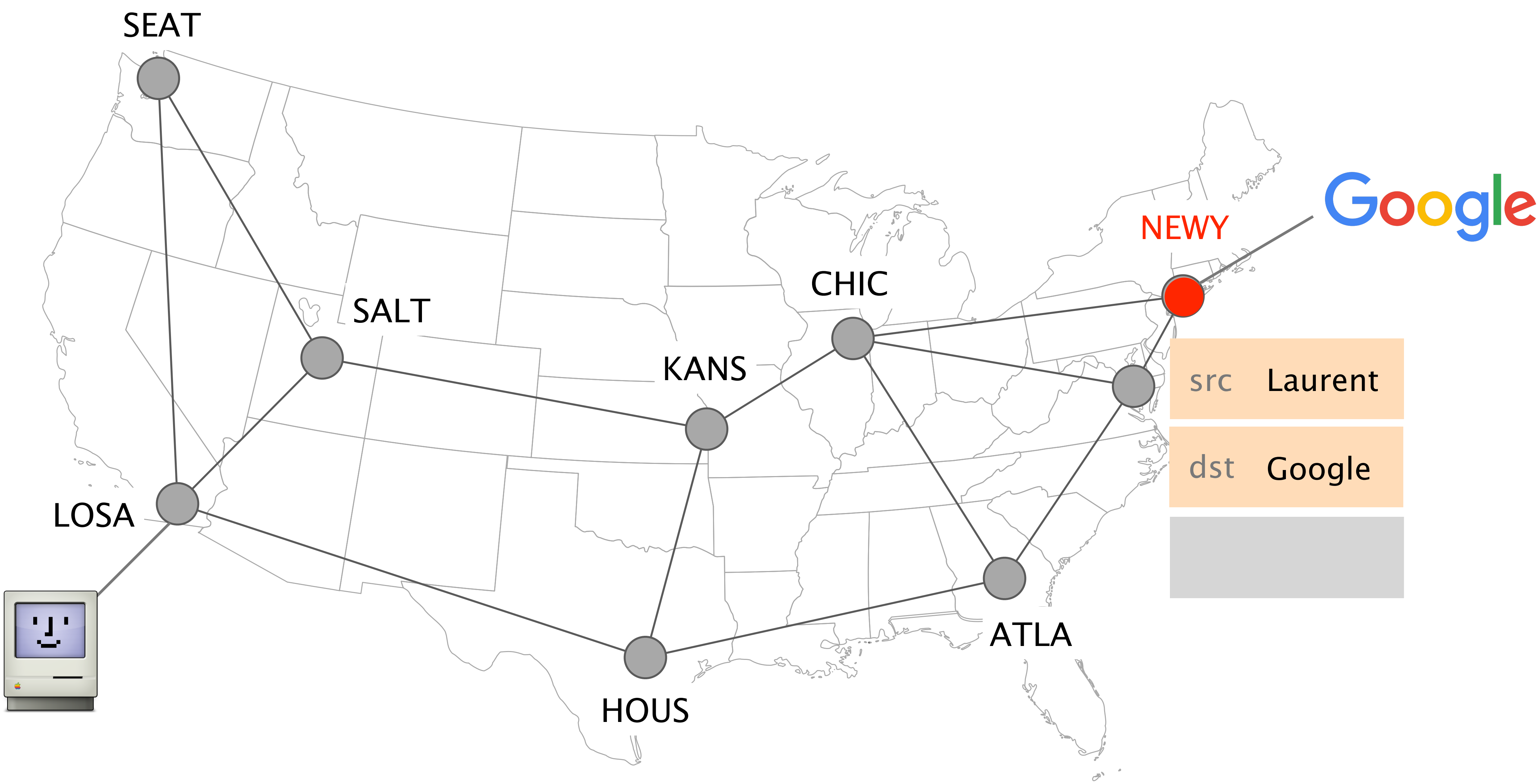
SEAT
SALT
LOSA
KANS
CHIC
NEWY
Google
HOUS
ATLA
src Laurent
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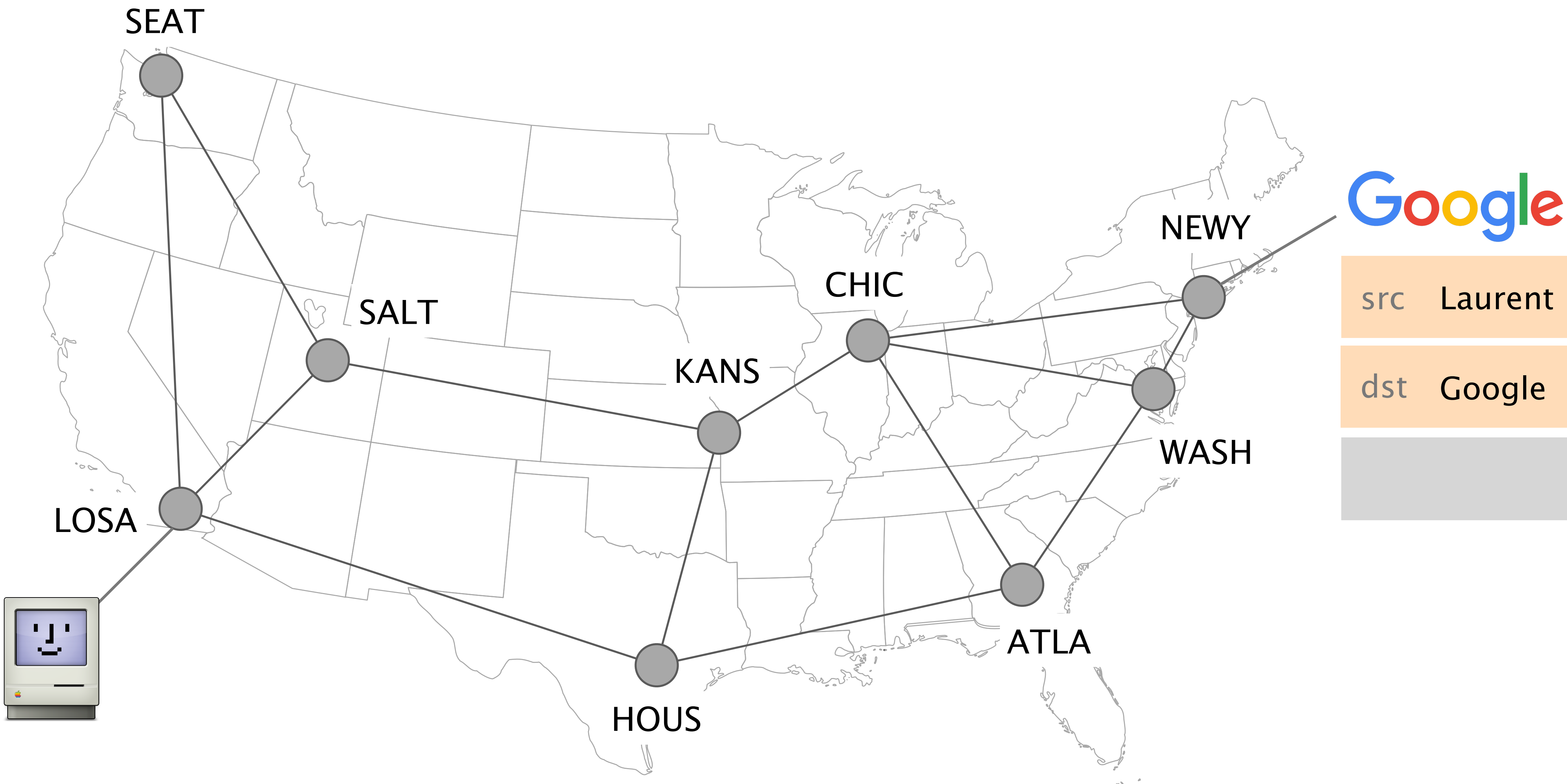
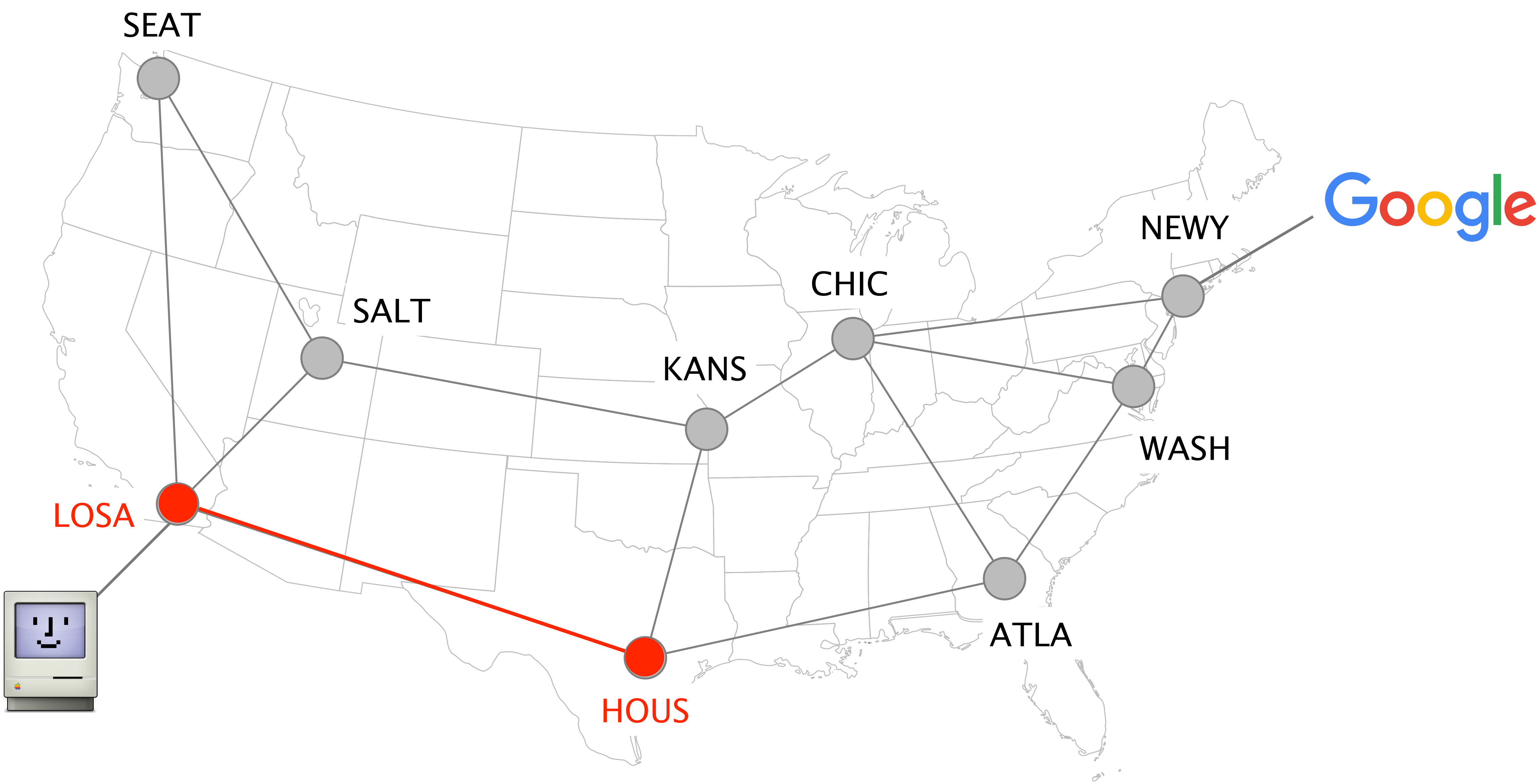
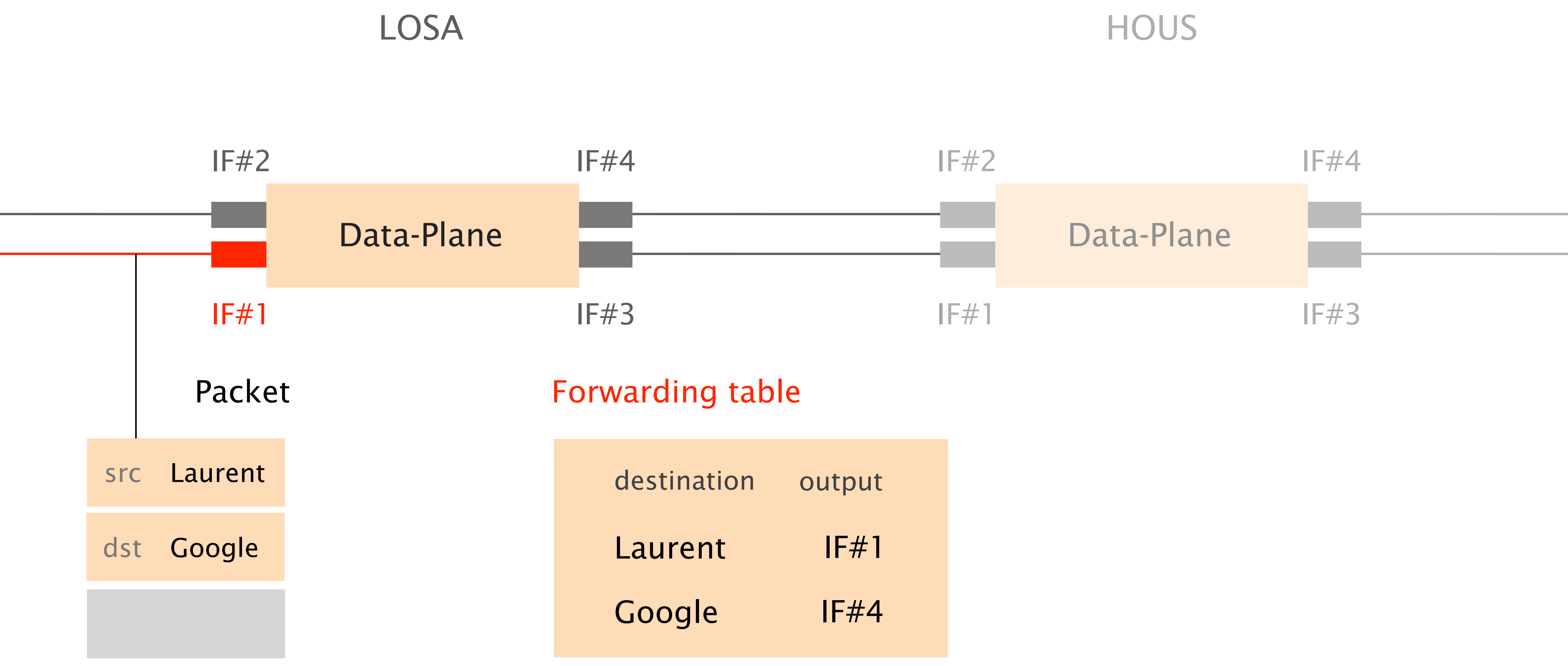

SEAT
SALT
LOSA
KANS
CHIC
NEWY
WASH
ATLA
HOUS
Google
src Laurent
dst Google

Let's zoom in on what is going on between **two adjacent routers**

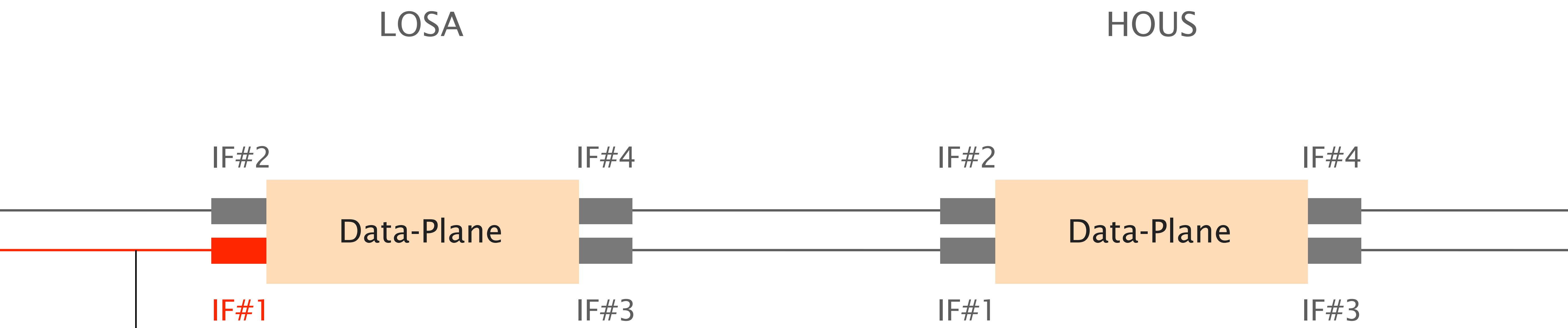



LOSA
HOUS
IF#2
IF#4
Data-Plane
IF#1
IF#3
IF#2
IF#4
Data-Plane
IF#1
IF#3

Upon packet reception, routers locally look up their forwarding table to know where to send it next



Here, the packet should be directed to IF#4

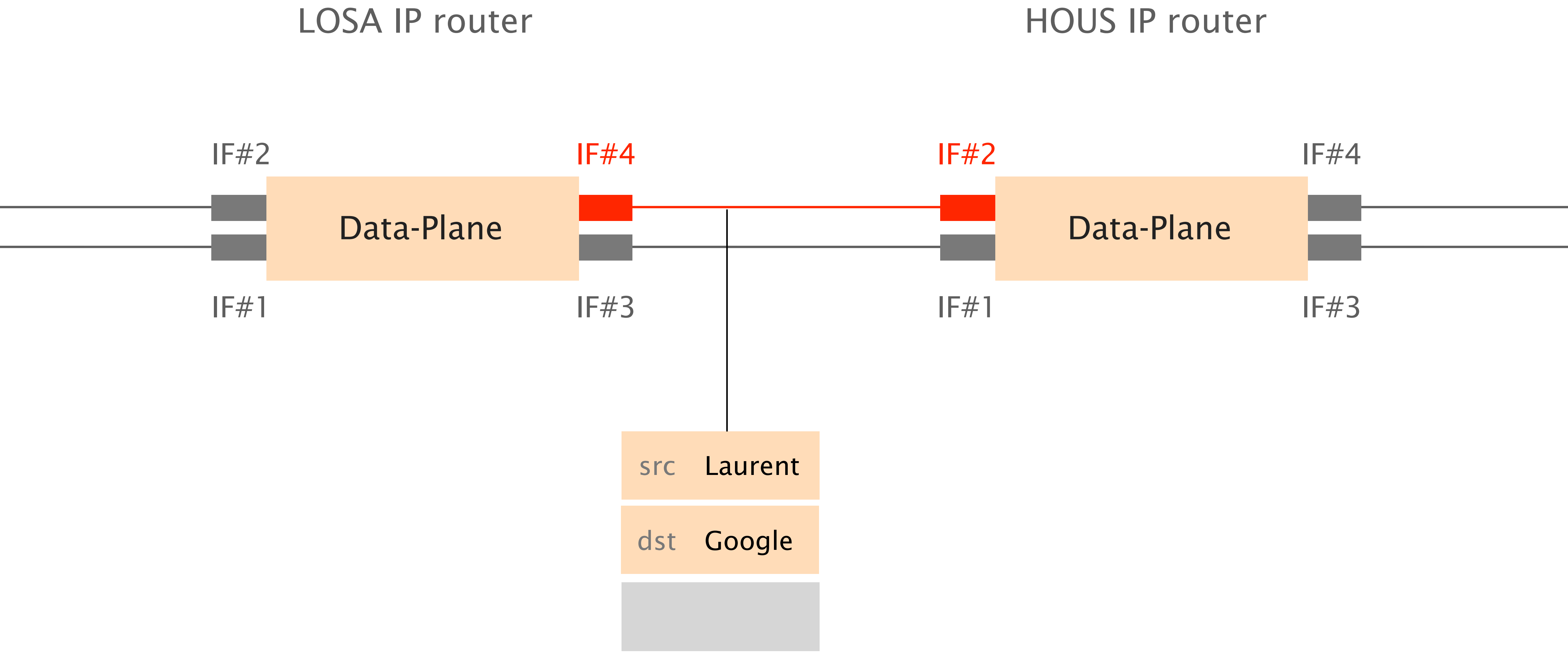


Packet

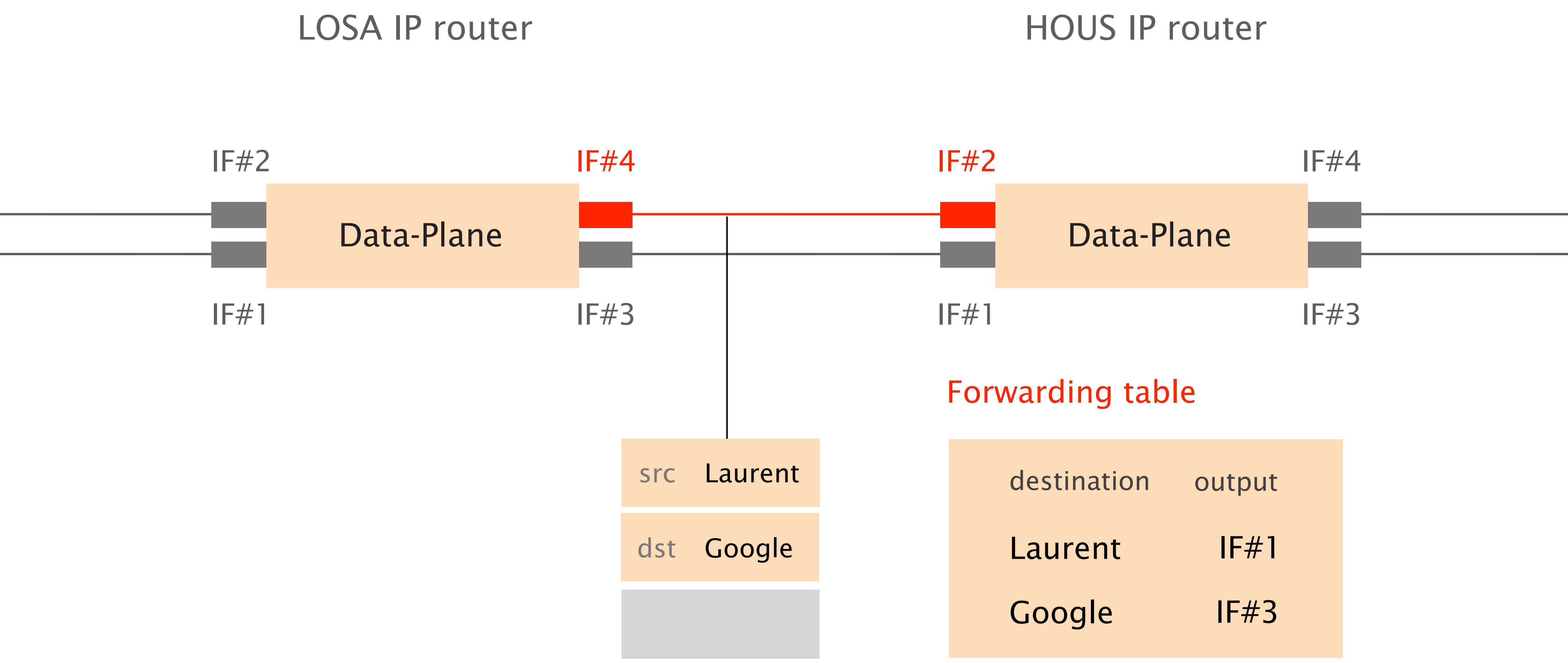
src	Laurent
dst	Google

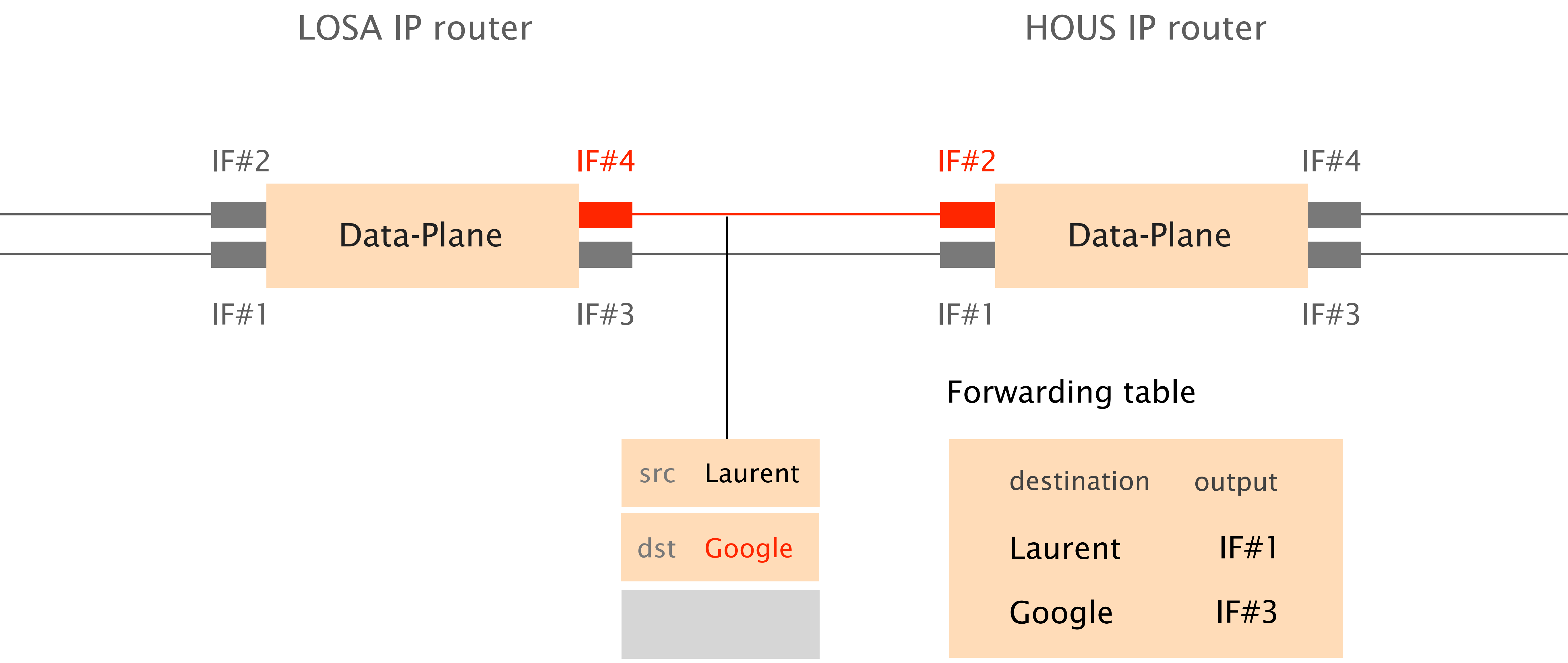
Forwarding table

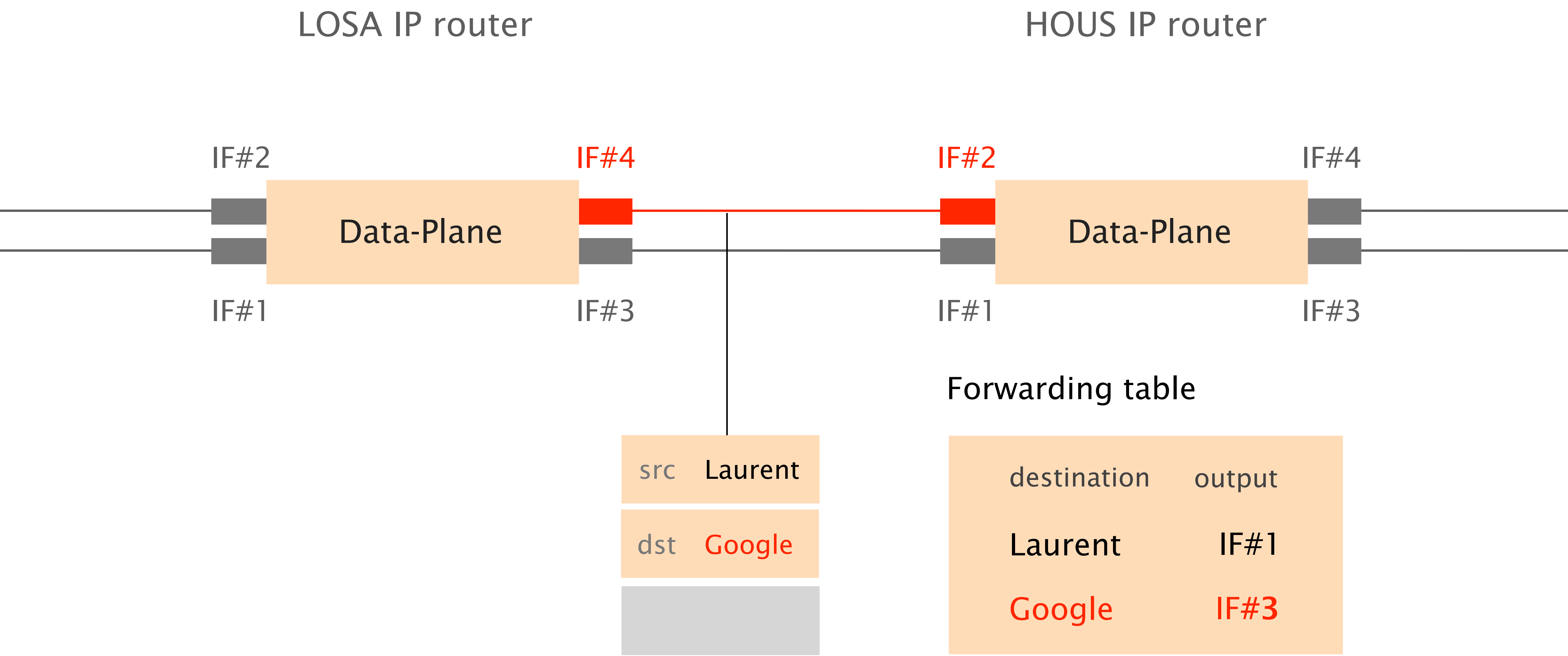
destination	output
Laurent	IF#1
Google	IF#4

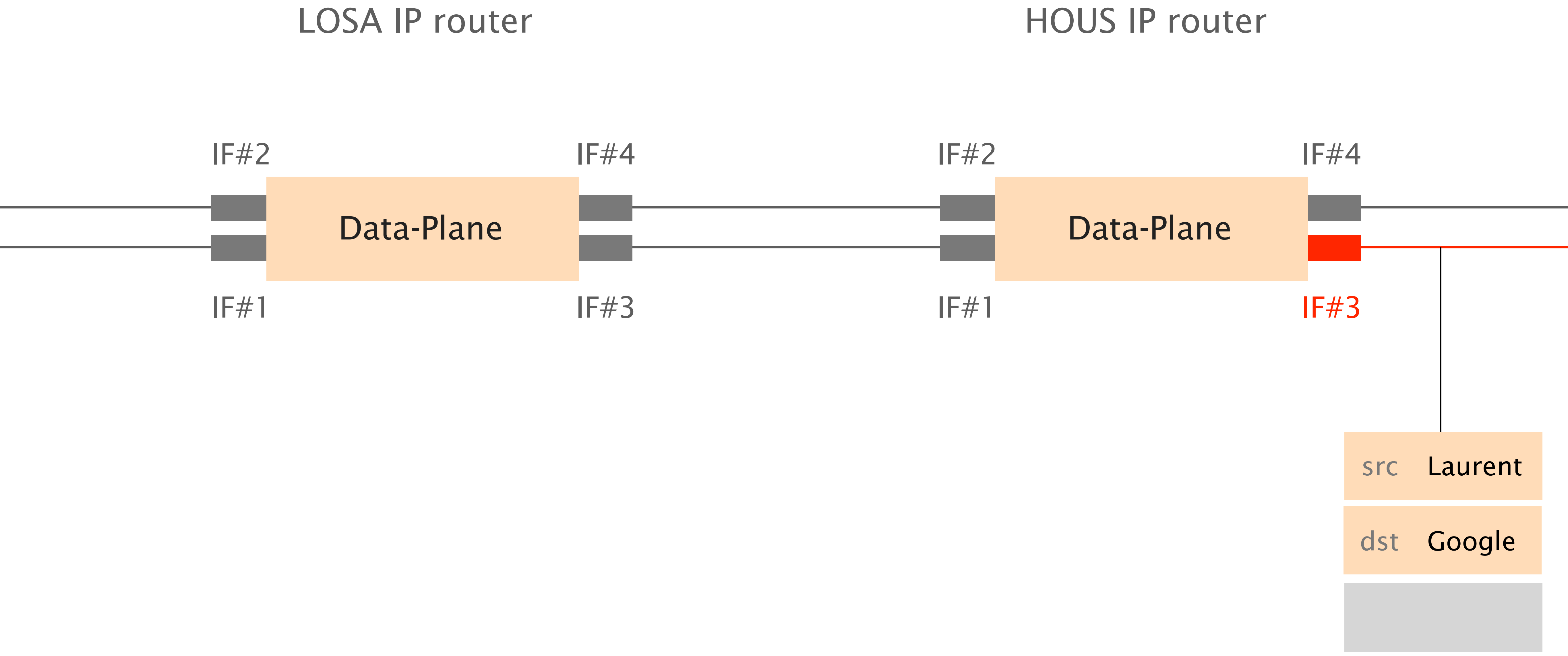

LOSA IP router
HOUS IP router
IF#2
IF#4
IF#2
IF#4
Data-Plane
Data-Plane
IF#1
IF#3
IF#1
IF#3
src Laurent
dst Google

Forwarding is repeated at each router, until the destination is reached




LOSA IP router
HOUS IP router
IF#2
IF#4
IF#2
IF#4
Data-Plane
Data-Plane
IF#1
IF#3
IF#1
IF#3
Forwarding table
src Laurent
dst Google
destination output
Laurent IF#1
Google IF#3


LOSA IP router
HOUS IP router
IF#2
IF#4
IF#2
IF#4
Data-Plane
Data-Plane
IF#1
IF#3
IF#1
IF#3
Forwarding table
src Laurent
dst Google
destination output
Laurent IF#1
Google IF#3


LOSA IP router
HOUS IP router
IF#2
IF#4
Data-Plane
IF#1
IF#3
IF#2
IF#4
Data-Plane
IF#1
IF#3
src Laurent
dst Google

Forwarding decisions necessarily depend on the destination, but can also depend on other criteria

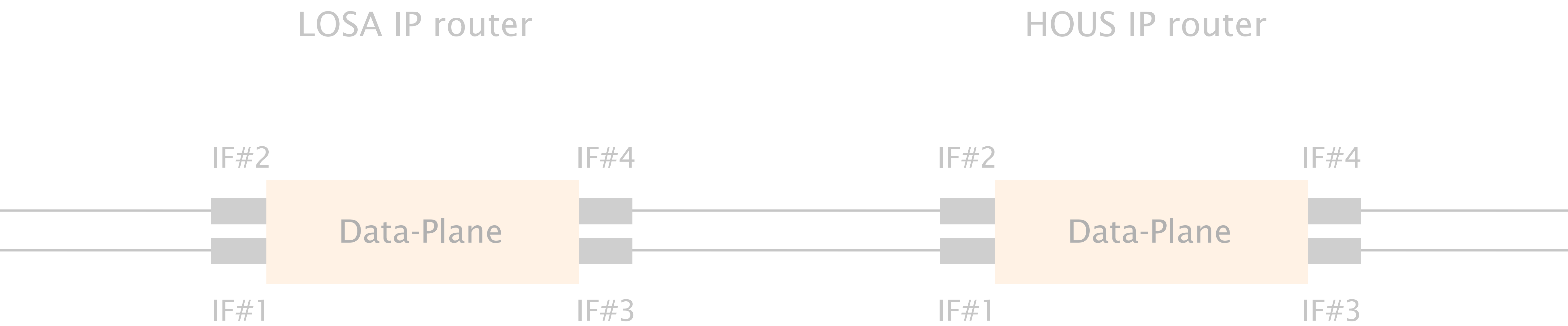
criteria		
	destination	mandatory (why?)
	source	requires n^2 state
	input port	traffic engineering
	+any other header	

In the rest of the lecture,

we'll consider destination-based routing

the default in the Internet

Where are these forwarding tables coming from?



Forwarding table

destination	output
Laurent	IF#1
Google	IF#4

Forwarding table

destination	output
Laurent	IF#1
Google	IF#3


Data-Plane
Data-Plane

In addition to a data-plane,

routers are also equipped with a control-plane

Think of the control-plane as the router's brain

Roles

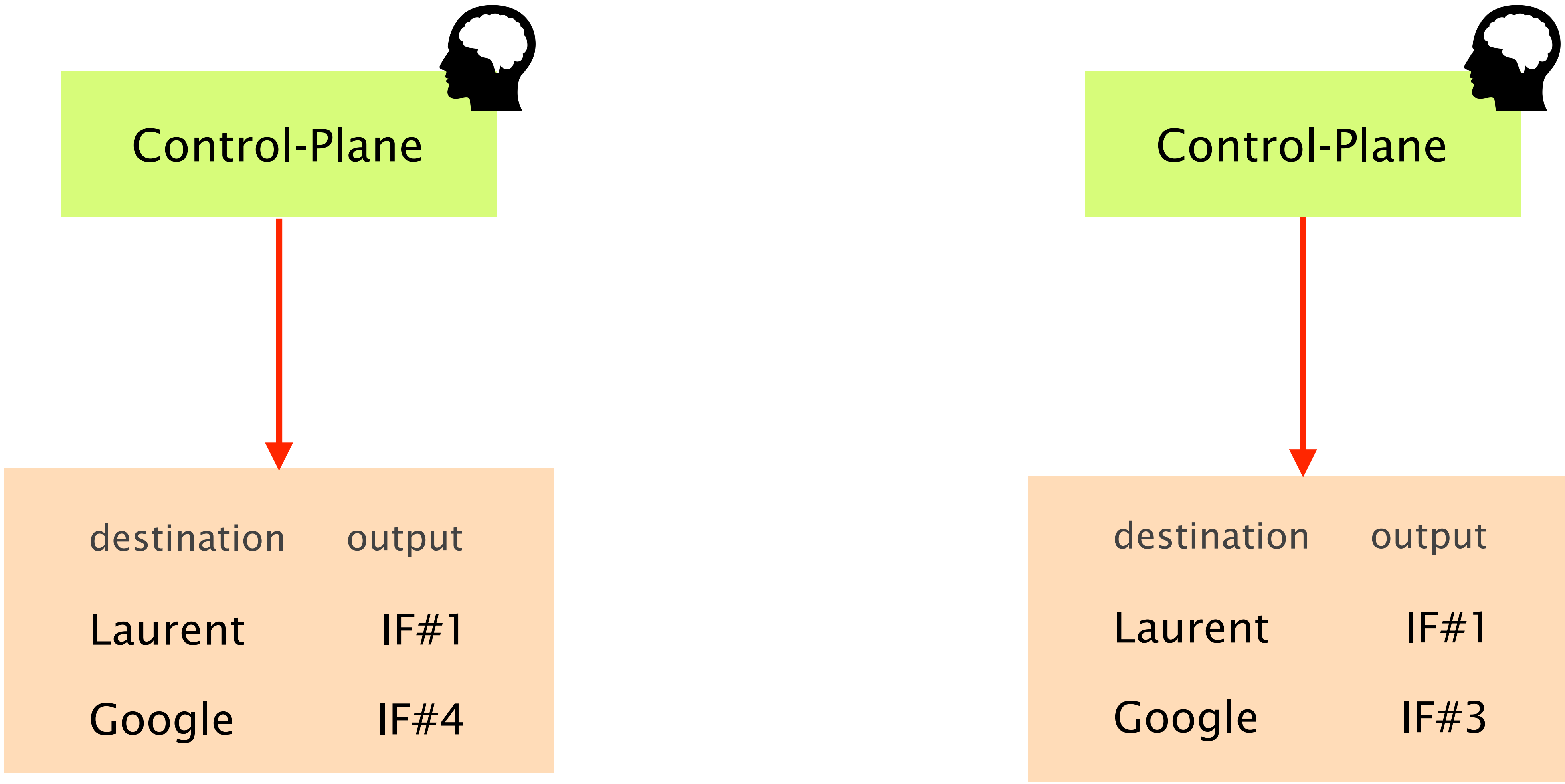
Routing

Configuration

Statistics

…

Routing is the control-plane process that **computes** and **populates** the forwarding tables



While forwarding is a *local* process,
routing is inherently a *global* process

How can a router know
where to direct packets
if it does not know what
the network looks like?

Forwarding *vs* Routing

summary

	forwarding	routing
goal	directing packet to an outgoing link	computing the paths packets will follow
scope	local	network-wide
implem.	hardware usually	software usually
timescale	nanoseconds	milliseconds (hopefully)

The goal of routing is to compute
valid global forwarding state

Definition

a global forwarding state is valid if

it always delivers packets
to the correct destination

sufficient and necessary condition

Theorem

a global forwarding state is valid if and only if

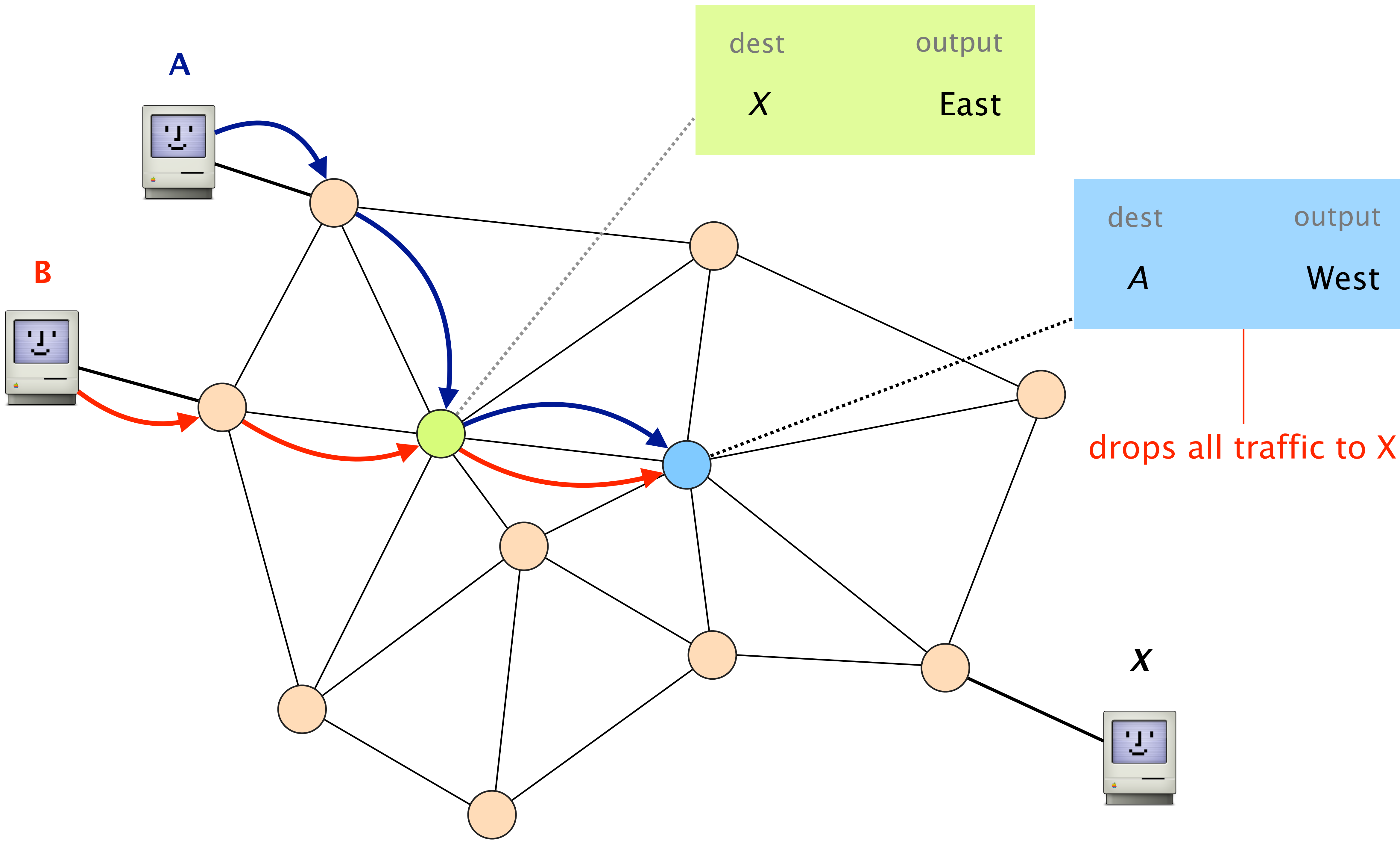
- there are no dead ends

  no outgoing port defined in the table

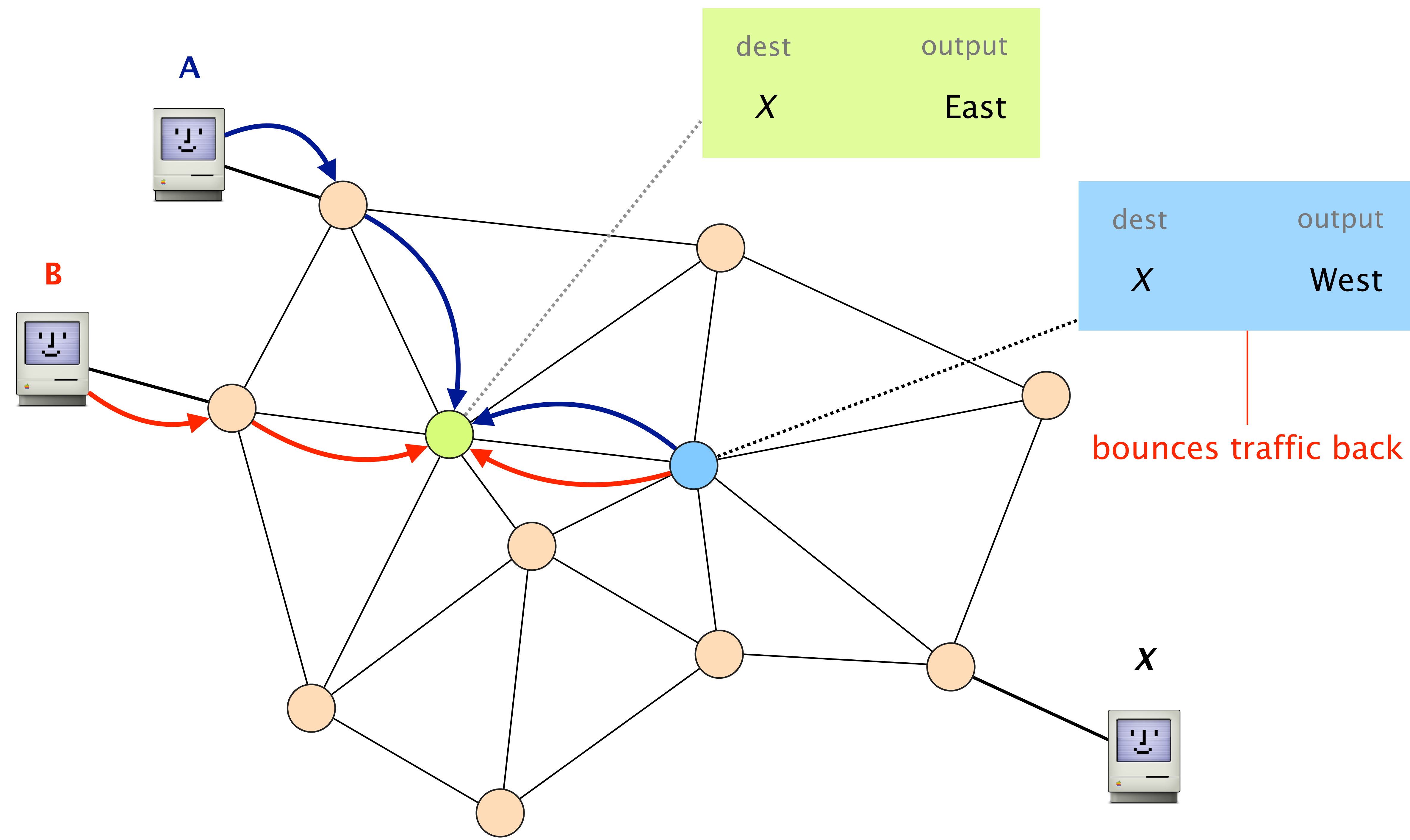
- there are no loops

  packets going around the same set of nodes

A global forwarding state is valid if and only if there are no dead ends



A global forwarding state is valid if and only if there are no forwarding loops



Proving the **necessary** condition is **easy**

Theorem

If a routing state is valid

then there are no loops or dead-end

Proof

If you run into a dead-end or a loop

you'll never reach the destination

so the state cannot be correct (contradiction)

Proving the **sufficient** condition is **more subtle**

Theorem

If a routing state has no dead end and no loop
then it is valid

Proof

There is only a finite number of ports to visit

A packet can never enter a switch via the same port, otherwise it is a loop (which does not exist by assumption)

As such, the packet must eventually reach the destination question 1 How do we verify that a forwarding state is valid?

question 2 How do we compute valid forwarding state?

question 1 How do we verify that a forwarding state is valid?

How do we compute valid forwarding state?

Verifying that a routing state is valid is easy

simple algorithm

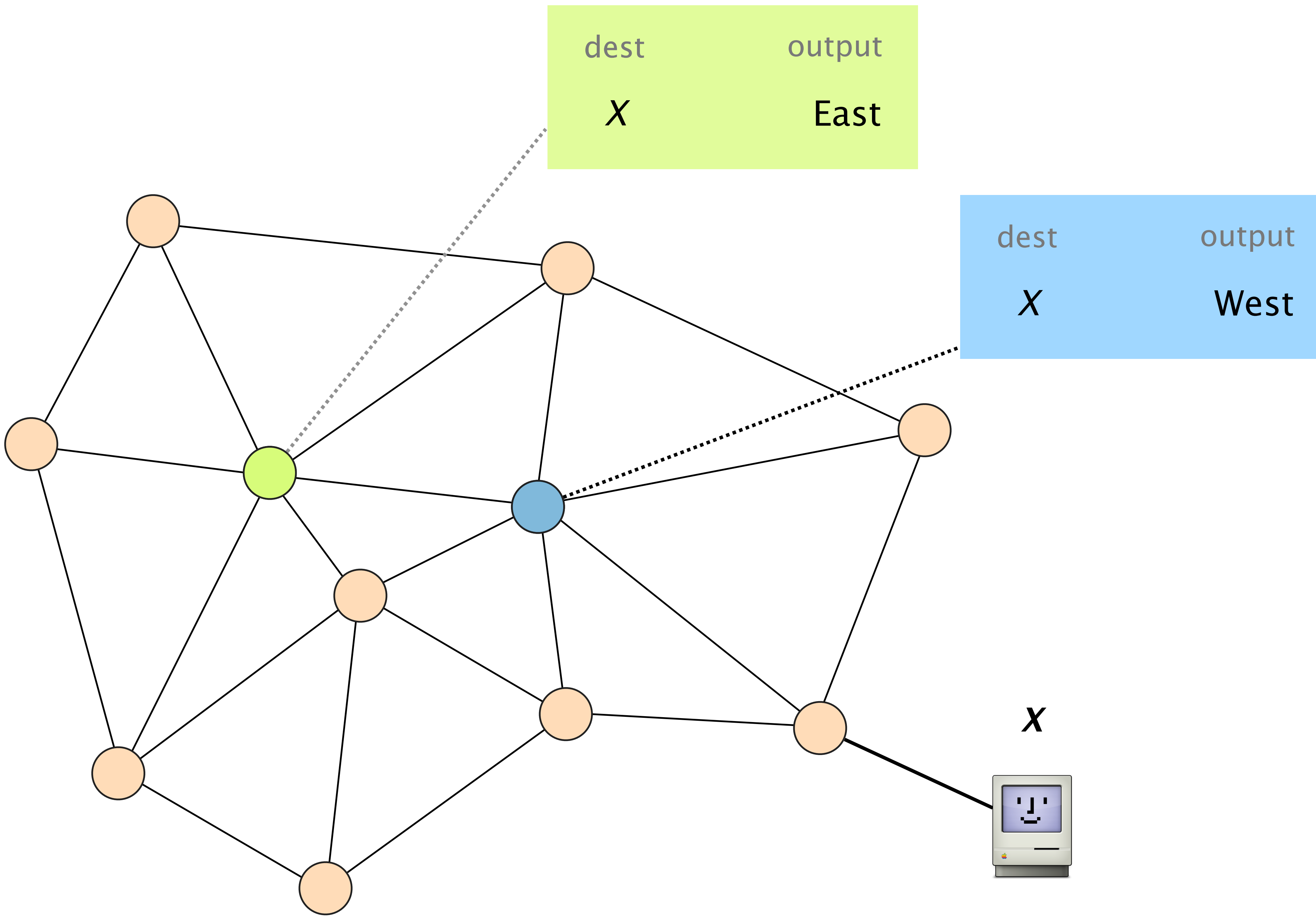
for one destination

Mark all outgoing ports with an arrow

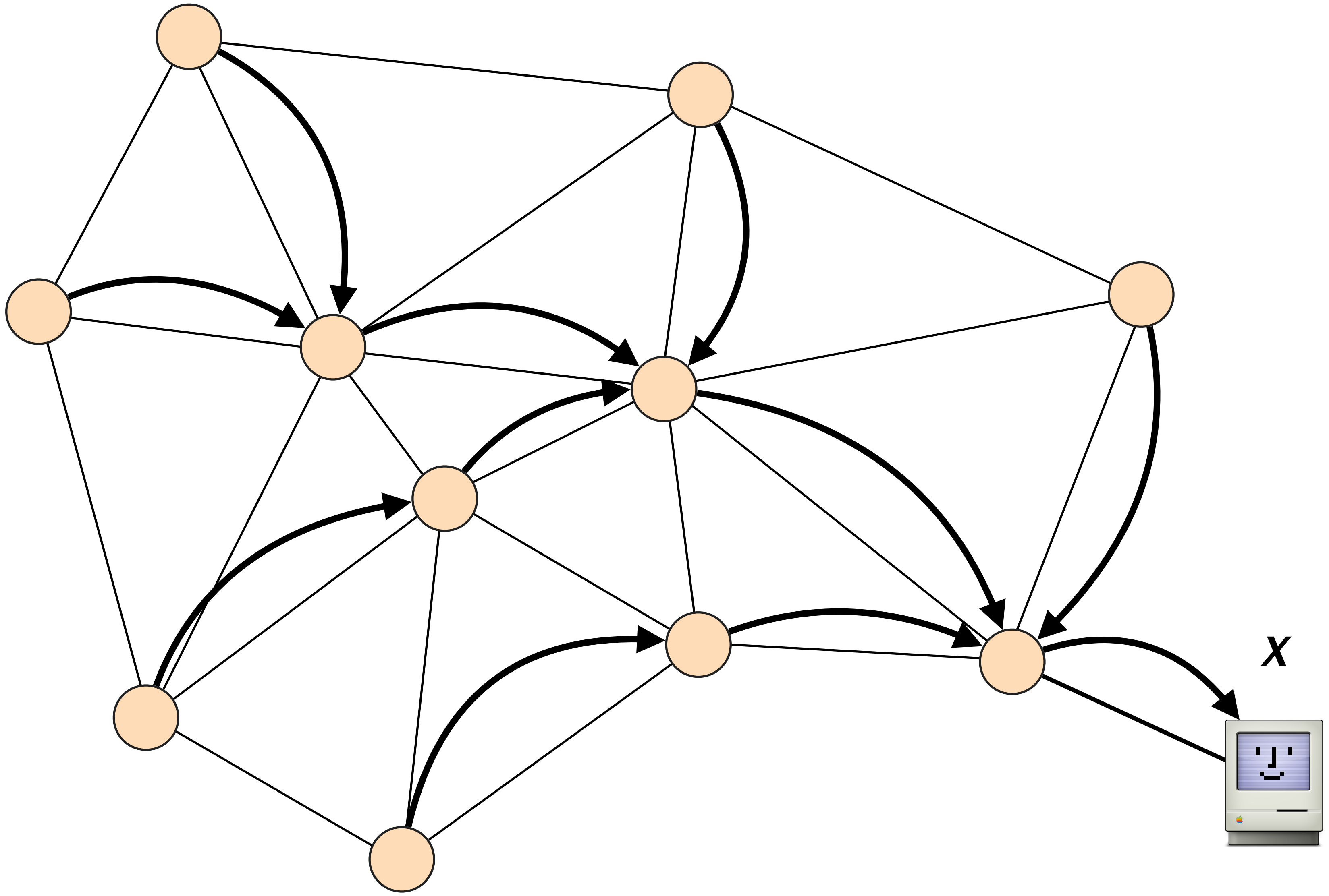
Eliminate all links with no arrow

State is valid *iff* the remaining graph is a spanning-tree

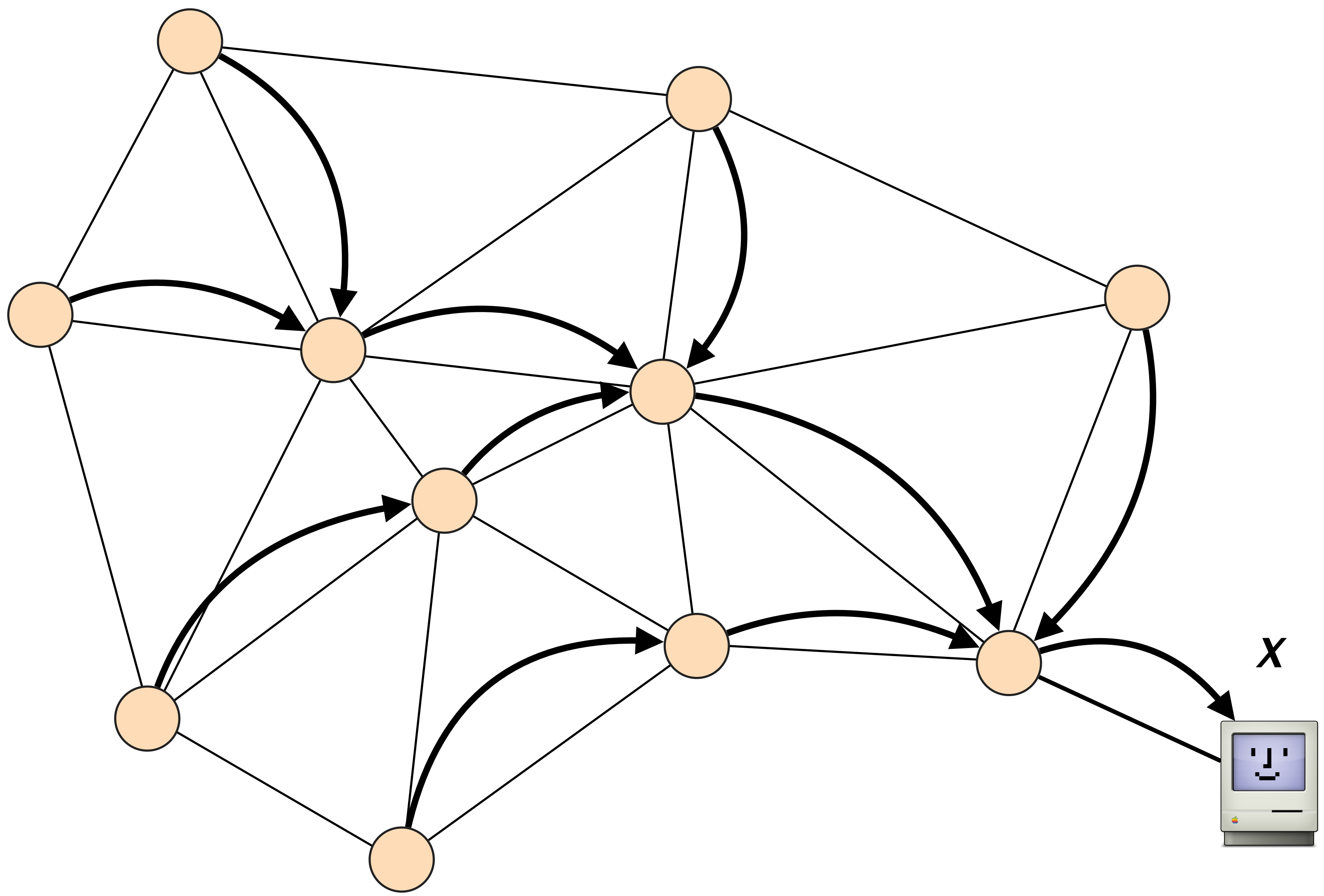
Given a graph with the corresponding forwarding state

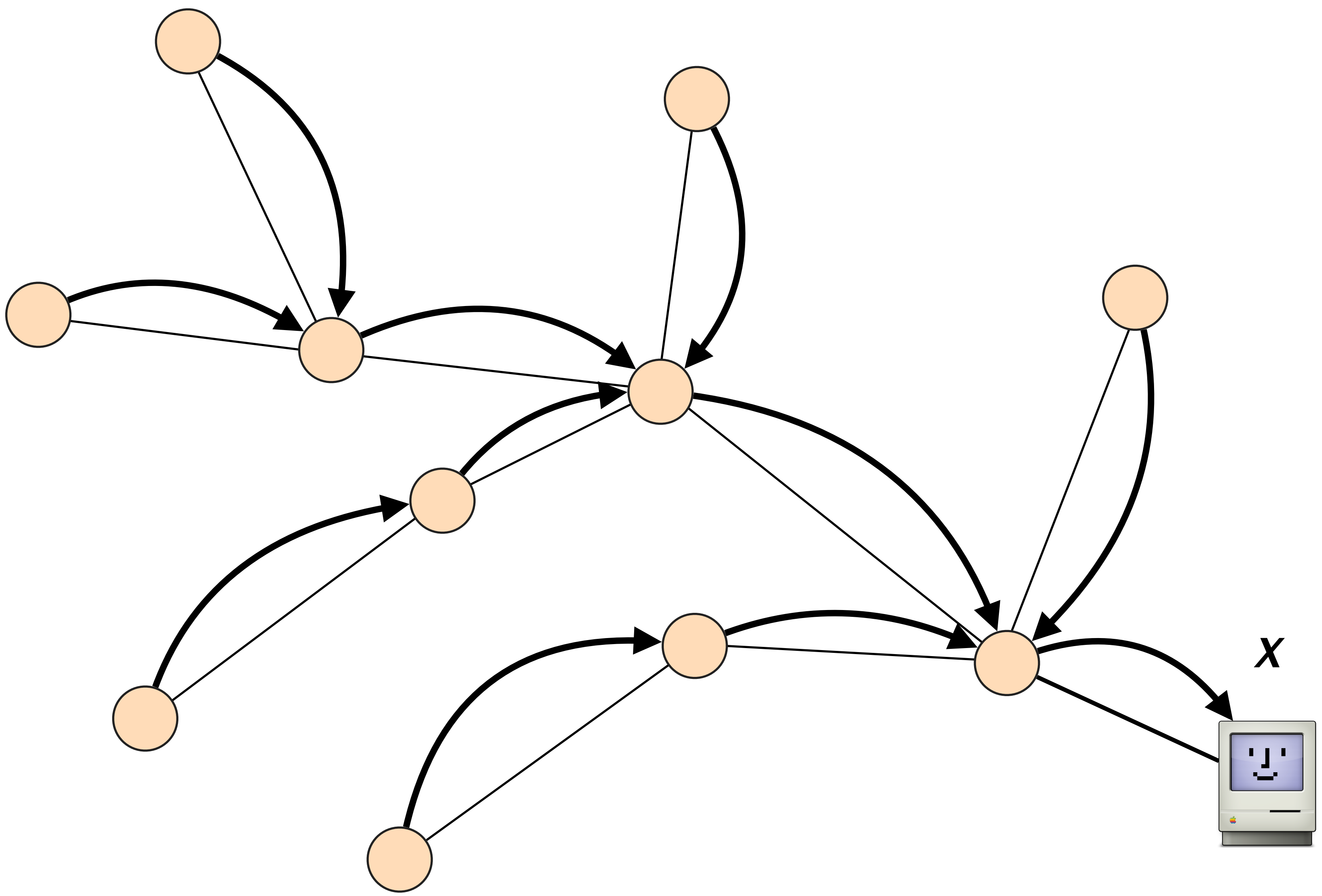


Mark all outgoing ports with an arrow



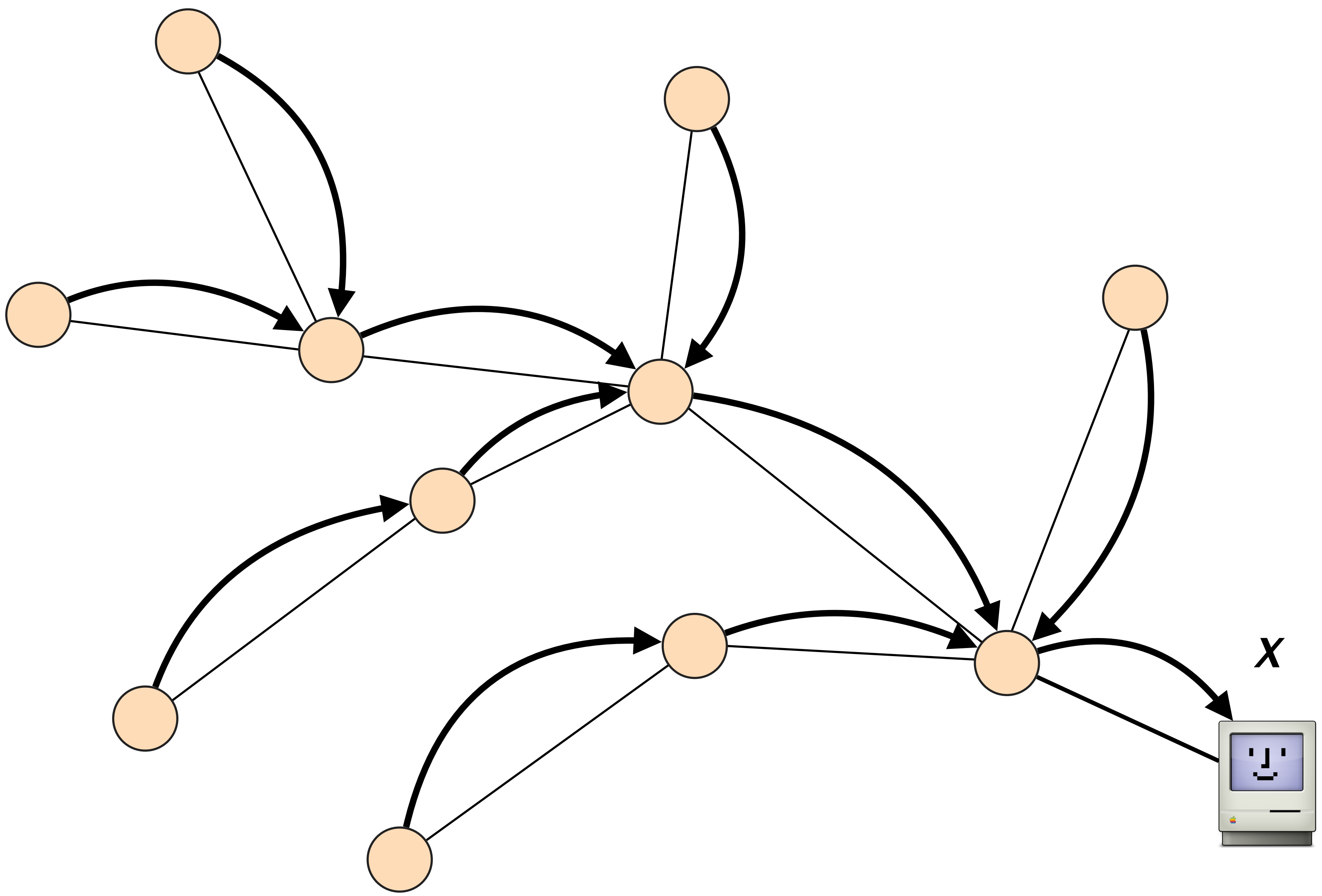
Eliminate all links with no arrow



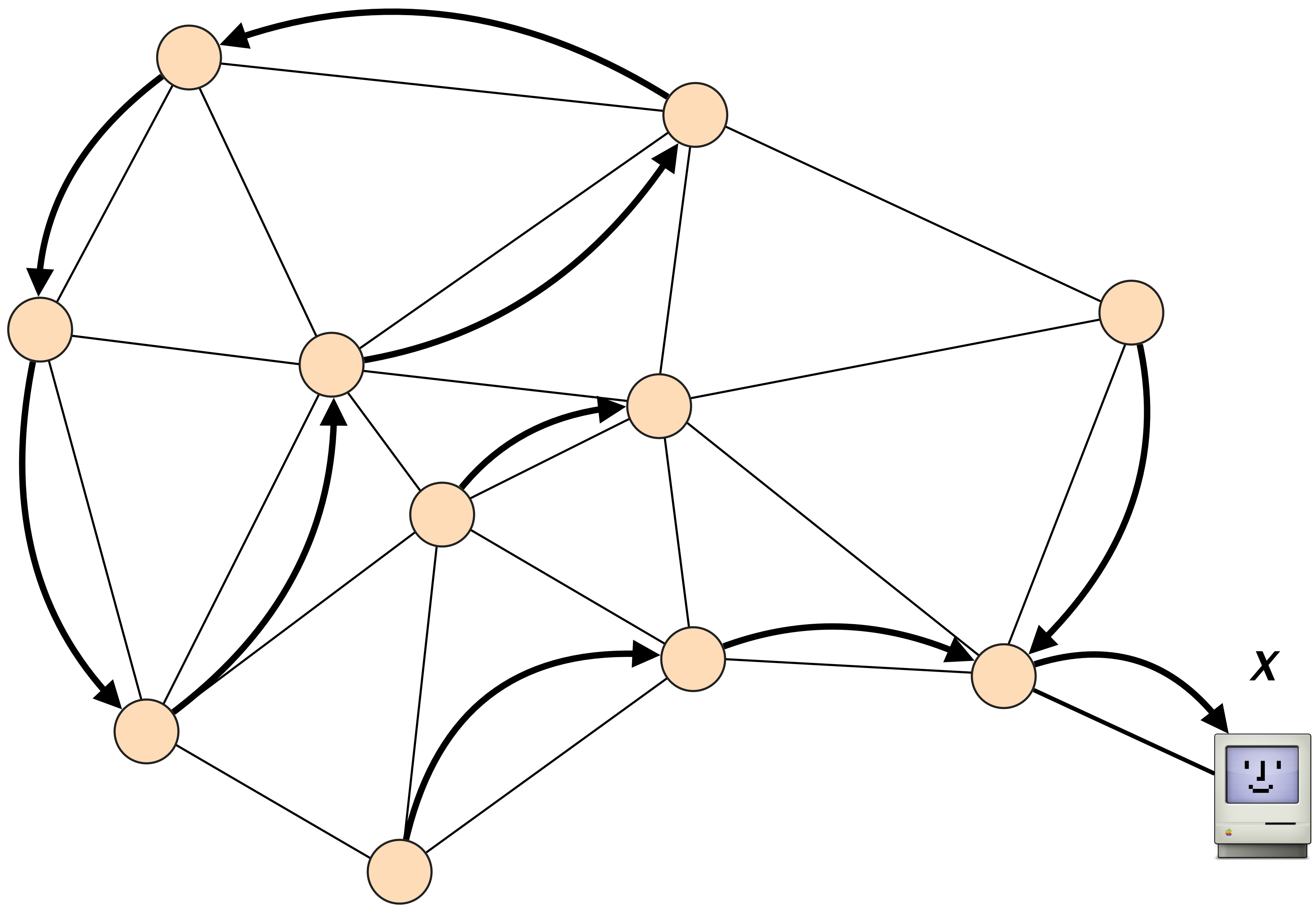

X

The result is a spanning tree.

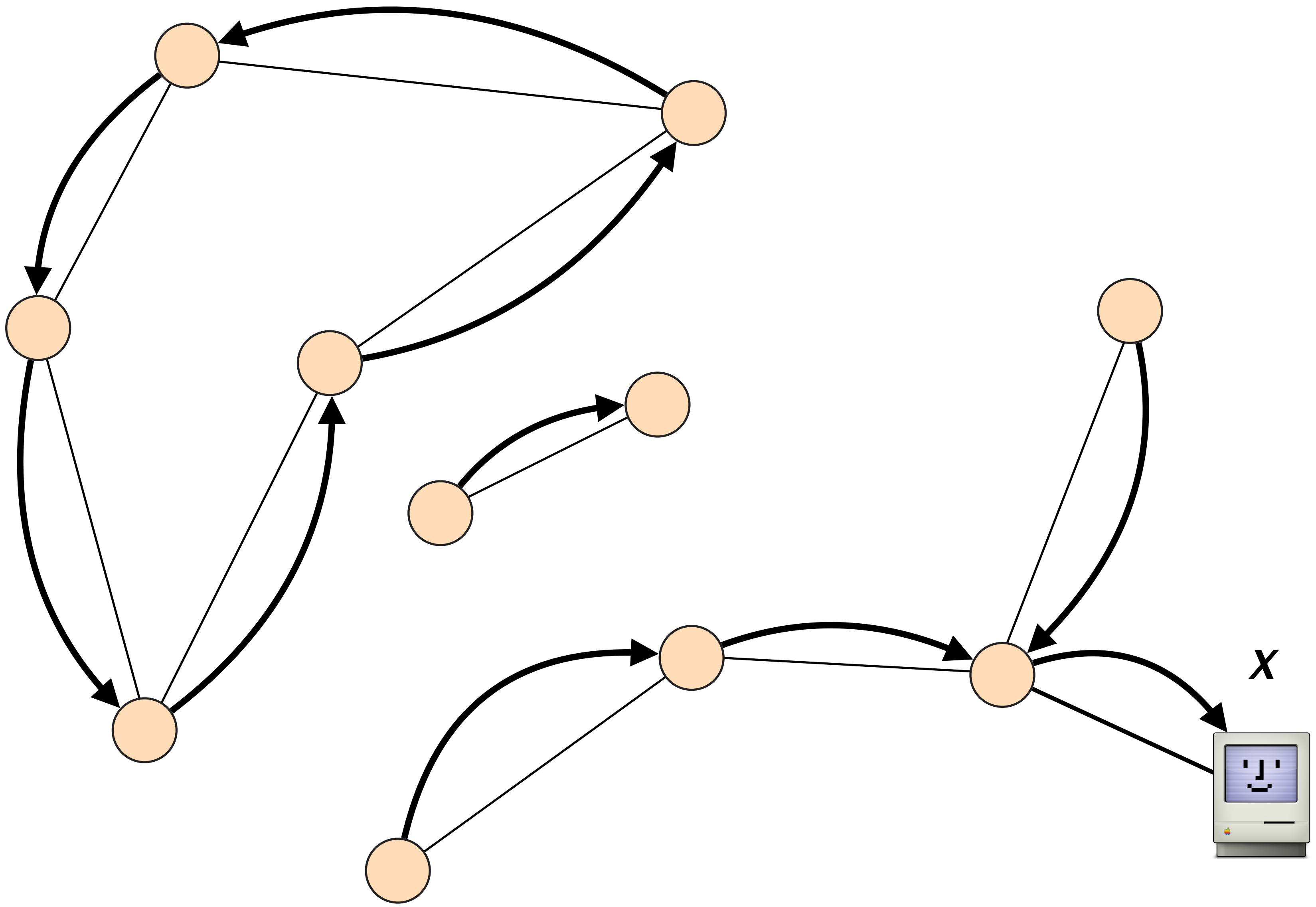
This is a valid routing state



Mark all outgoing ports with an arrow

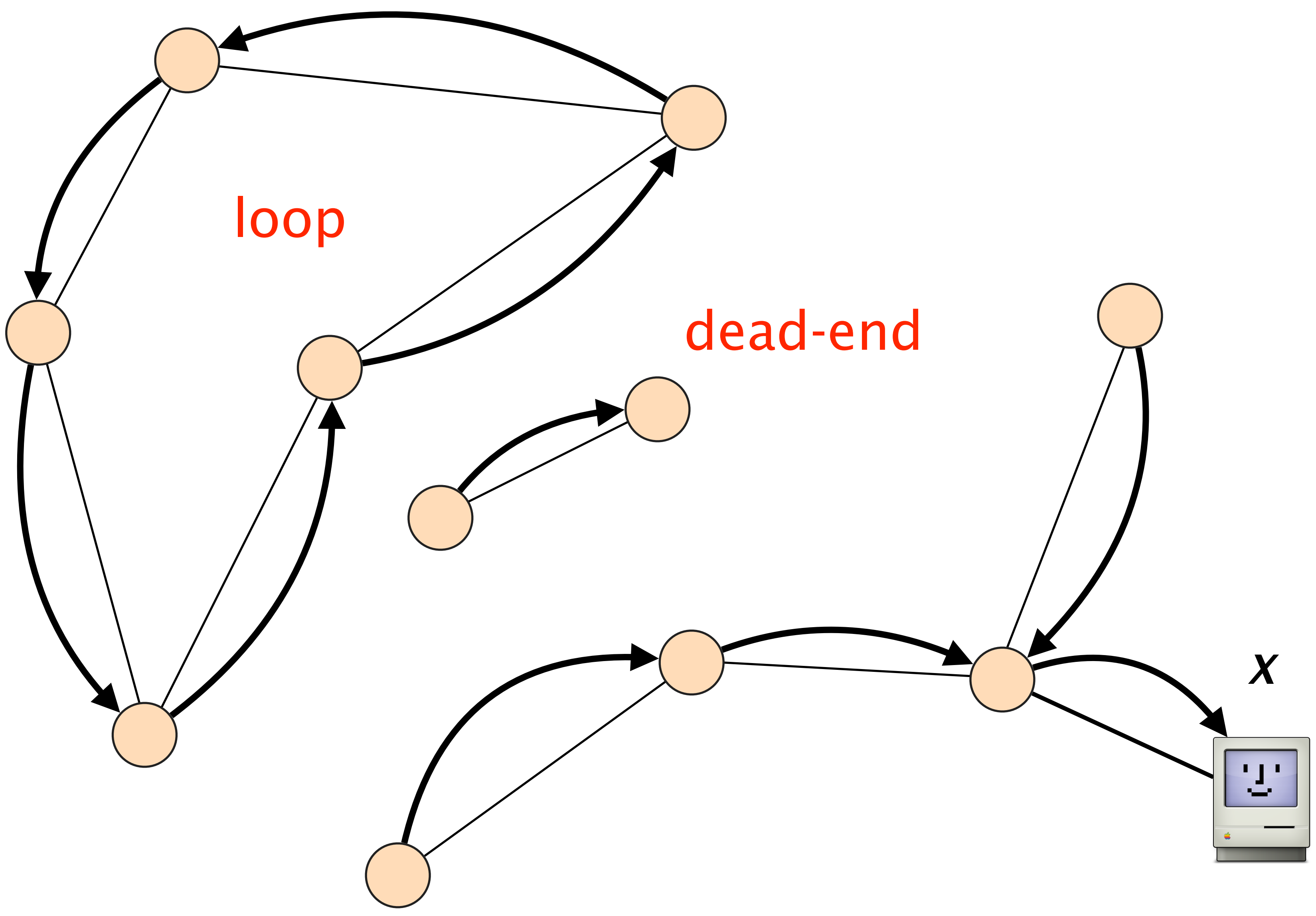


Eliminate all links with no arrow



The result is not a spanning-tree.

The routing state is not valid



How do we verify that a forwarding state is valid?

question 2

How do we compute valid forwarding state?

Producing valid routing state is harder

prevent dead ends

easy

prevent loops

hard

Producing valid routing state is harder
but doable

prevent dead ends

easy

prevent loops

hard

This is the question
you should focus on

Existing routing protocols differ in

how they avoid loops

prevent loops

hard

Essentially,

there are **three ways to compute** valid routing state

	Intuition	Example
#1	Use tree-like topologies	Spanning-tree
#2	Rely on a global network view	Link-State SDN
#3	Rely on distributed computation	Distance-Vector BGP

Communication Networks

Spring 2020

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February 24 2020

Materials inspired from Scott Shenker & Jennifer Rexford